

Efficient Location Based Services for Groups of Mobile Users



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14th IEEE intl. conf. on mobile data management
Milan, 3-6 June 2013

Observation

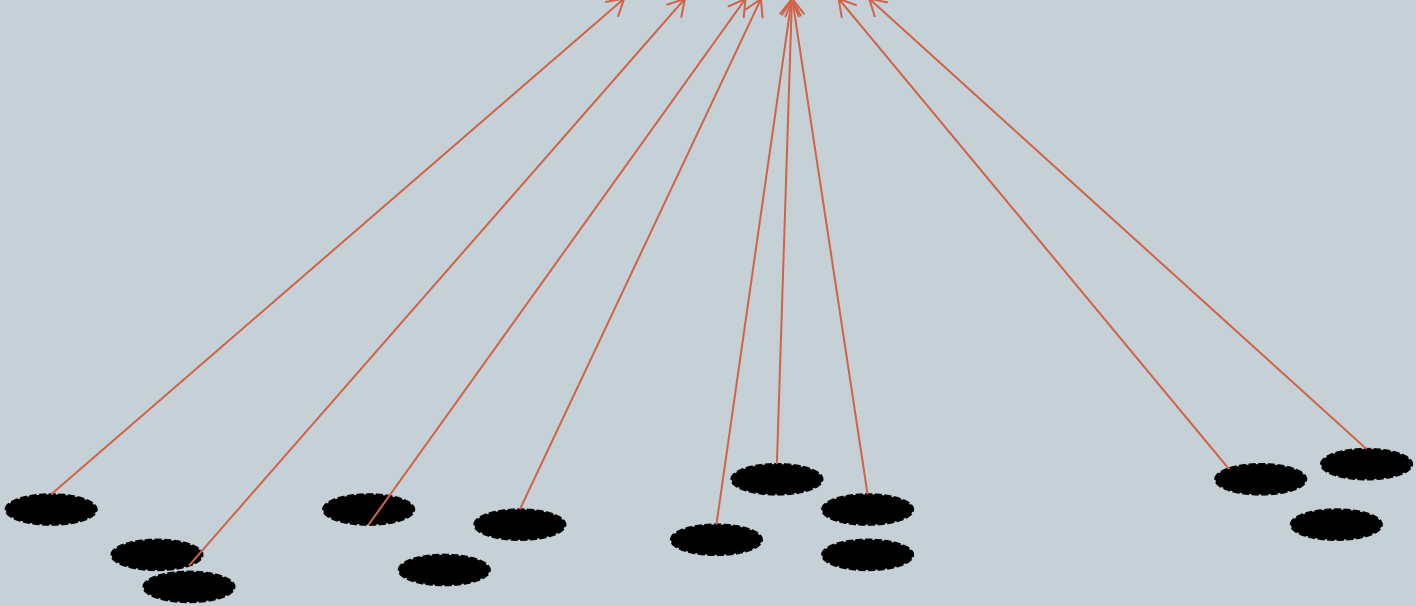


- Mobile users **send** location updates to LBS server (back-end system; BES)
- The BES **pushes content** (possibly, personalized) to a mobile user **if** she enters/crosses a **region**

Observation



LBS server

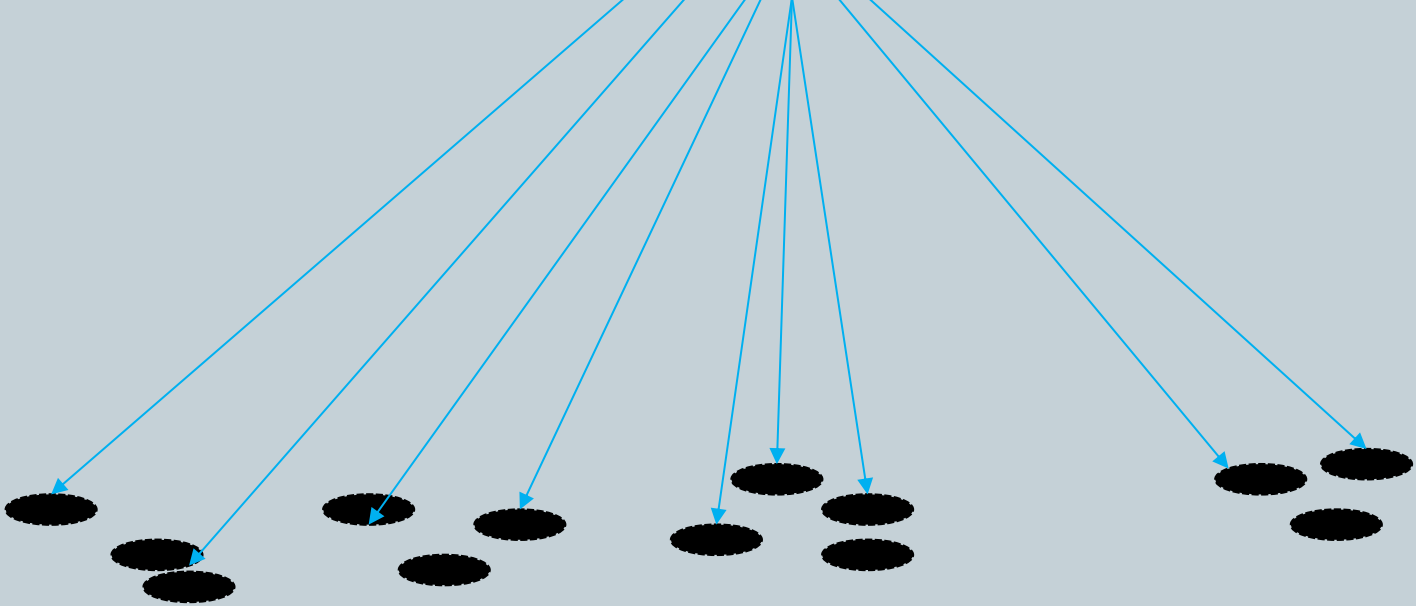


- Location update report
- Mobile node

Observation



LBS server



- Location-dependent content
- Mobile node

Observation



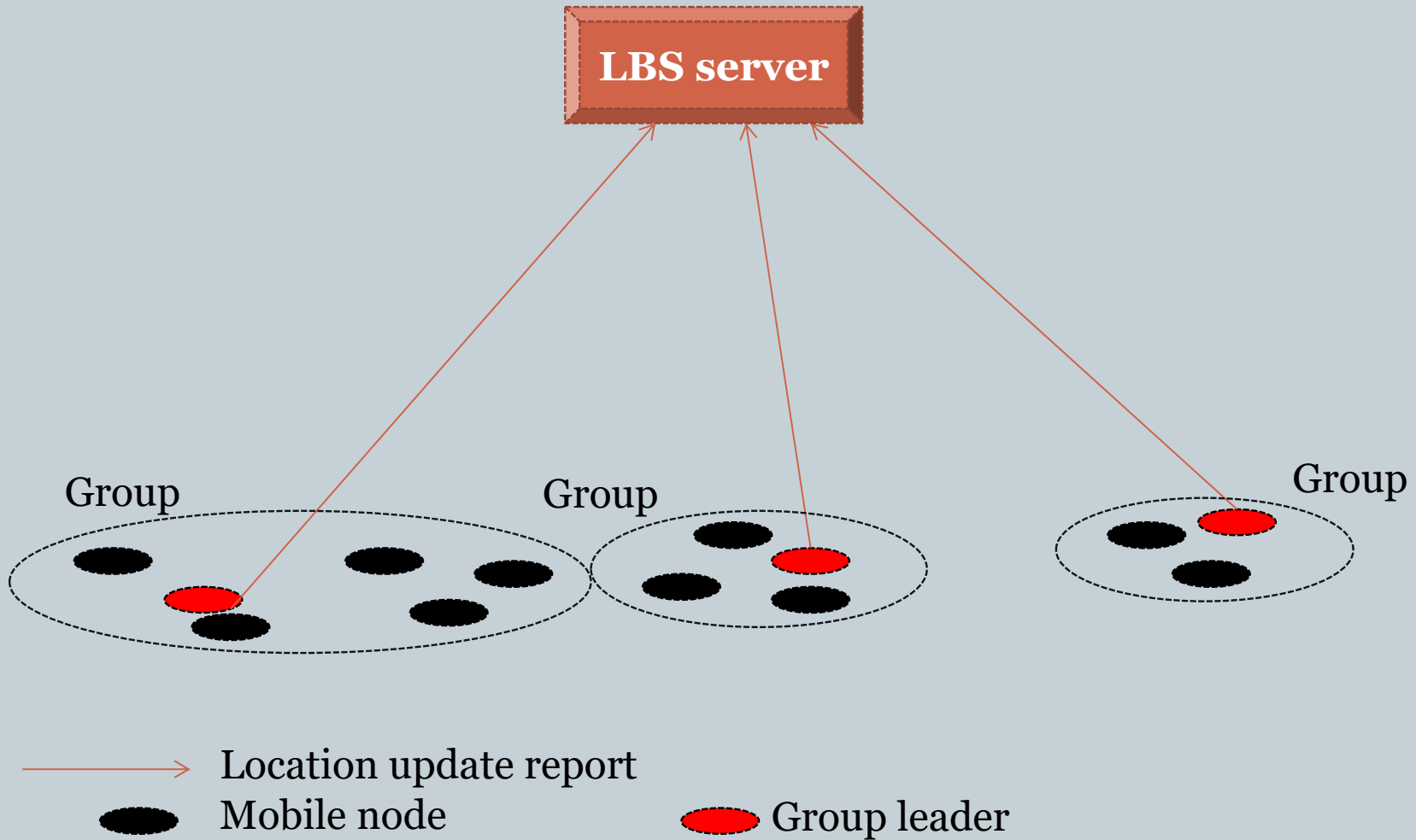
- As the **number** of mobile users increases, and
- As the users **continuously change** their position **then**
- The **communication overhead** for location reporting *and* location-based content delivery becomes quite significant

Group Identification

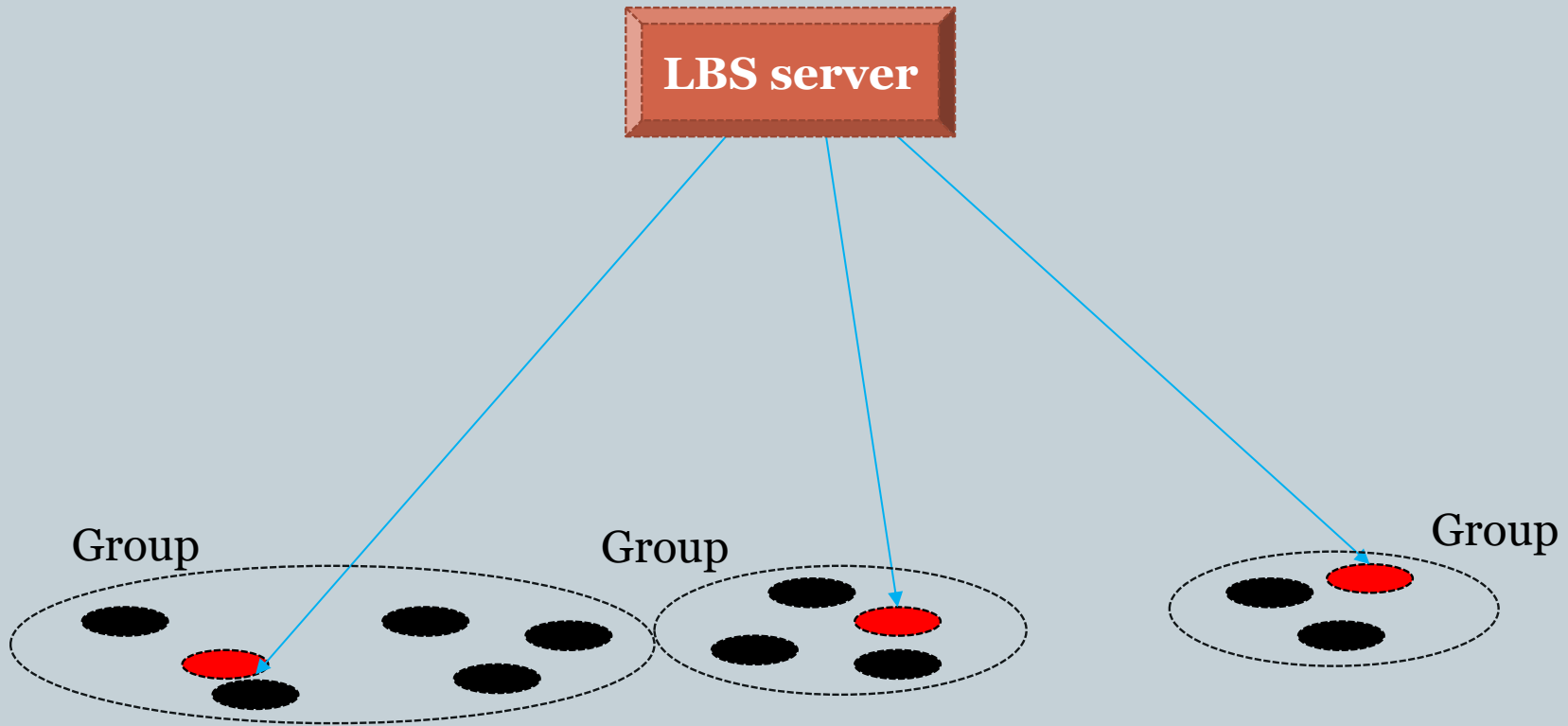


- Mobile users are not only **close to each other** but also likely to **move together** for a certain time horizon, thus, forming a **moving group**.
- A group has a unique **group leader** (GL).
- GL ...is the **representative** of the group, i.e.,
 - **sends** location updates to BES
 - **receives** possible location-dependent content and **disseminates** it to its members

Moving Object Groups Monitoring



Moving Object Groups Monitoring

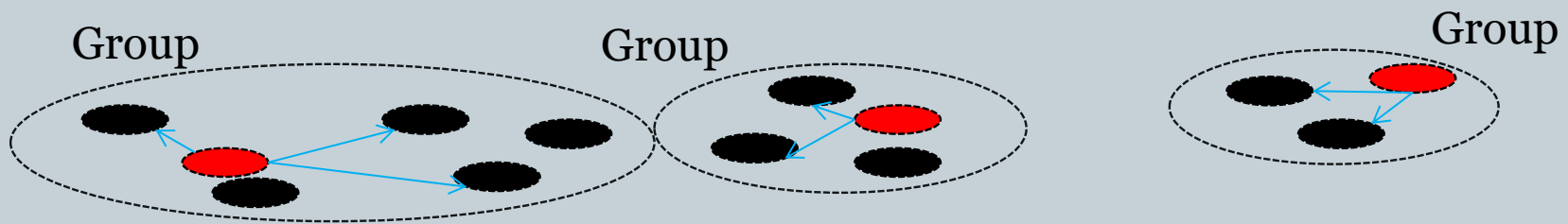


- ← Location-dependent content
- Mobile node
- Group leader

Moving Object Groups Monitoring



LBS server



- ← Location-dependent content
- Mobile node
- Group leader

Problem



- **How** to *evaluate* that a formed group Q at time t , will remain the **same** group at time $t+k$, $k>1$
- ...‘same’ group means:
 - **coherence**: group consists of the same mobile users as **initially** identified, and
 - **compactness**: all group members are **within** the communication range of the GL during a *finite time horizon*

Some Definitions



Group Q is defined as: $Q = \langle l, V_l \rangle$

$$l = \arg \min_{u \in Q} \left\| \mathbf{p}_u - \mathbf{c}_Q \right\| \quad \text{i.e., the GL of } Q$$

$$\mathbf{c}_Q = \frac{1}{|Q|} \sum_{u \in Q} \mathbf{p}_u \quad \text{i.e., the center point of } Q$$

$$V_l = \left\{ u \mid \left\| \mathbf{p}_u - \mathbf{p}_l \right\| \leq R \right\} \quad \text{i.e., the group members}$$

Some Definitions



Group validity at (discrete) time t :

$$x_t = \left(\frac{1}{|Q|} \sum_{u \in Q} \left\| \mathbf{p}_u^t - \mathbf{c}_Q^t \right\| \right) \cdot \mathbf{I}_{[\exists u \in V_l^t : \left\| \mathbf{p}_u^t - \mathbf{p}_l^t \right\| > R]} \left(V_l^t \right)$$

i.e., the average distance between members, *within the range of GL* , and group centroid

Group Validity and Validation System



- Consider a *clustering algorithm*, which is invoked periodically **every** N time instances and produces a set of groups $\mathbf{Q} = \{Q_1, Q_2, \dots, Q_{|\mathbf{Q}|}\}$;
- Let $Q \in \mathbf{Q}$ be formed at t_0 . Q is **valid** at $t > t_0$
if $x_t \leq \theta$
- The lower the x_t is, the more compact Q is at t

Group Validity and Validation System



Group Validation System (GVS)

- [1] *monitors the behavior of Q in $[t_0, k]$, i.e., checks whether Q maintains its **initial** structure.*
- [2] *considers Q as a **persistent** group in $(k, t_0 + N]$*

At time t in $[t_0, k]$

- (1) **each** member reports its location to BES; $f_t(|Q|)$
- (2) GVS **evaluates** the group validity x_t ; $g_t(|Q|)$

(3) Total **cost** c_t up to t is
$$c_t = \sum_{l=t_0}^t (f_l(|Q|) + g_l(|Q|))$$

Group Validity and Validation System



If GVS validates Q at k then for t in $(k, t_0 + N]$

- (1) BES communicates **only** with GL;
- (2) **only** the GL **reports** its location to BES;
- (3) (possible) content is delivered to the members **through** the GL.

Formation, Validation & Persistence Phase



validation decision



t_0

k^*

$t_0 + N$



group formation
phase (clustering)

Group Validity and Validation System



The problem is to find (an optimal) time k^* , $1 < k \leq N$:

$$\inf_{1 < k \leq N} E[X_k + c_k] = \inf_{1 < k \leq N} E[X_k + kc]$$

Optimal Stopping Theory



- Choose the **best** time instance to take a decision of performing a certain action.
- **Observe** the current state of a system and decide whether to:
 - **continue** the process or
 - **stop** the process, and incur a certain **cost**.
- ...the *odds* algorithm, the *secretary* problem, the *parking* problem, the *asset-selling* problem, etc.

Application to Group Validity Process



- Adoption of Optimal Stopping Theory for evaluating an criterion for k^* .
- The more **validity** values the GVS *observes*, the more *certain* is on concluding on a ‘group persistence’
- The GVS **observes** X_t at each t . The **decision** is:
 - **stop** observing X_t and classify Q as *persistent*
 - **continue** observing X_{t+1} with *additional cost* C
- ‘Finite’ Horizon GVS (FVS), i.e., $1 < k \leq N$
- ‘Infinite’-like Horizon GVS (IVS), i.e., $1 < k$

'Finite' Horizon GVS (FVS)

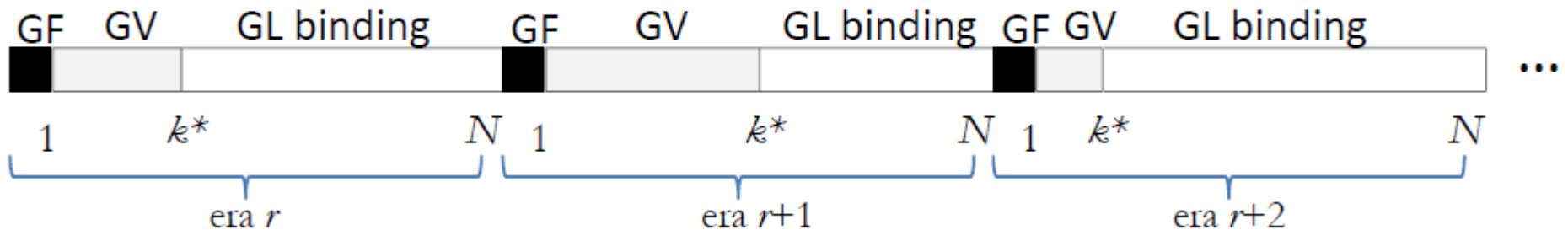


- FVS should *validate Q up to N*.
- The criterion is the sequence (a_1, a_2, \dots, a_N) such that FVS stops at t **iff** $x_t < a_t$.

$$a_t = a_{t+1} (P_X(X \leq x) - P_X(X \leq a_{t+1})) + \int_0^{a_{t+1}} x dP_X(x) + c$$

- The {c} 'princ' $a_N = E[X] + c$ in the 'ction'.

(FVS)



'Infinite'-like Horizon GVS (IVS)



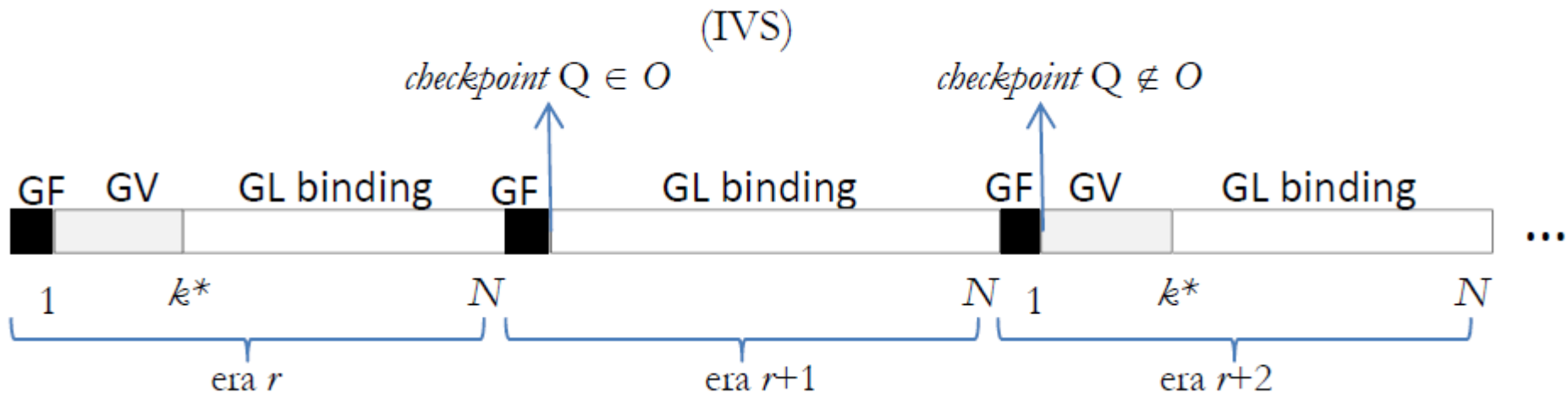
- IVS *validates* Q independent of N .
- At certain **checkpoints**, it checks whether Q maintains its initial structure

$$\int_0^{a^*} (a^* - x) dP_X(x) = c$$

- If Q is ...
the v...

ting

- The criterion for optimal stopping is: stop at t **iff** $x_t < a^*$



Performance Evaluation

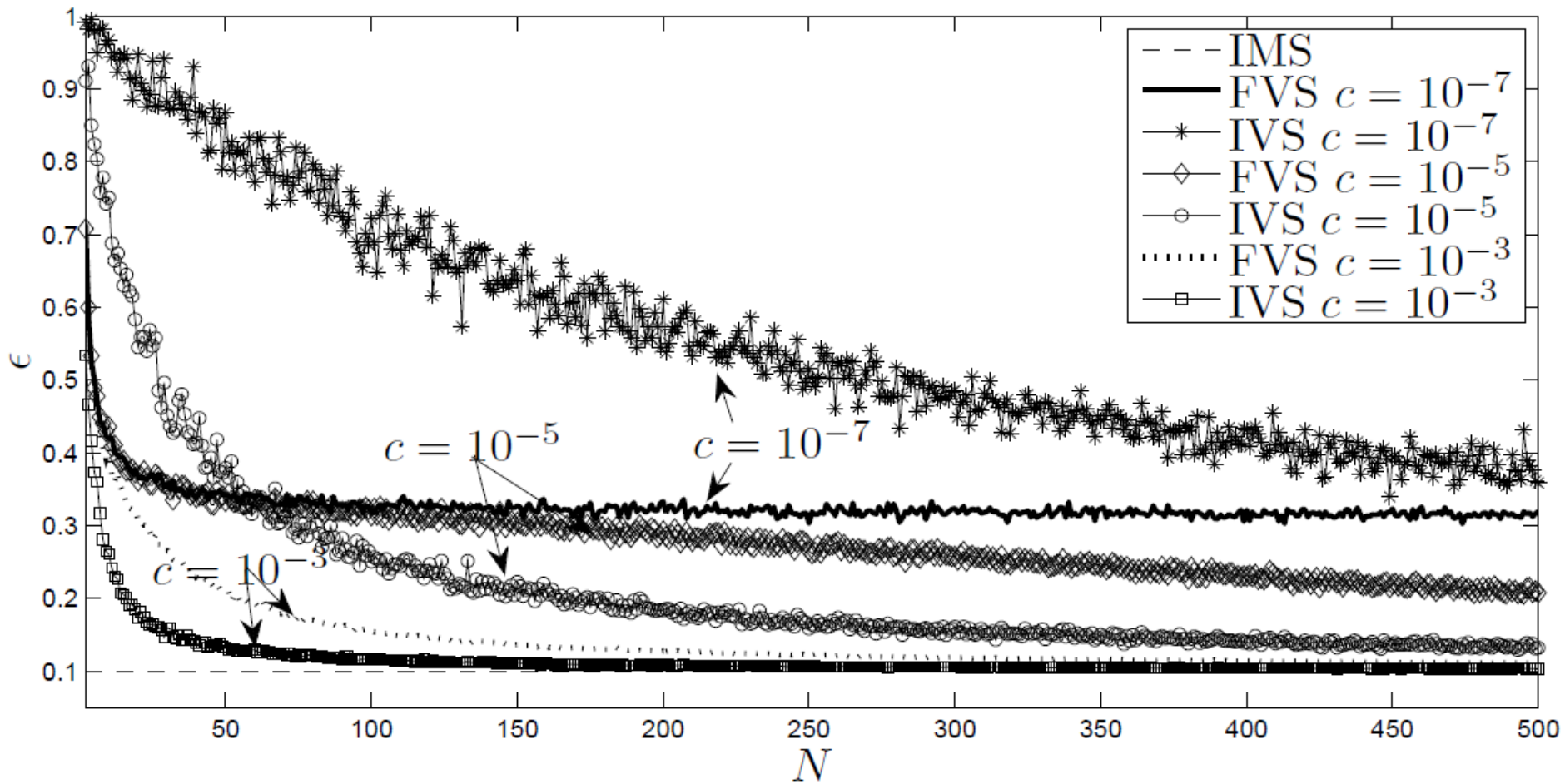


Metric ϵ :communication overhead during N (total number of messages) **w.r.t.** a continuous monitoring system.

- For a *continuous* monitoring system: $|Q|N$
- For FVS is: $|Q|k^* + (N - k^*)$
- For IVS is: $|Q|k^* + (Nm - k^*)$; m^{-1} = re-validation rate
- For IMS (immediate validation system) is: N

IMS: **periodically** performs re-clustering with freq. $1/N$

The load savings



Performance Evaluation



Required communication load:

$I_t = 1$, if Q is *valid* at t ;

$I_t \in (1, |Q|]$, **otherwise**;

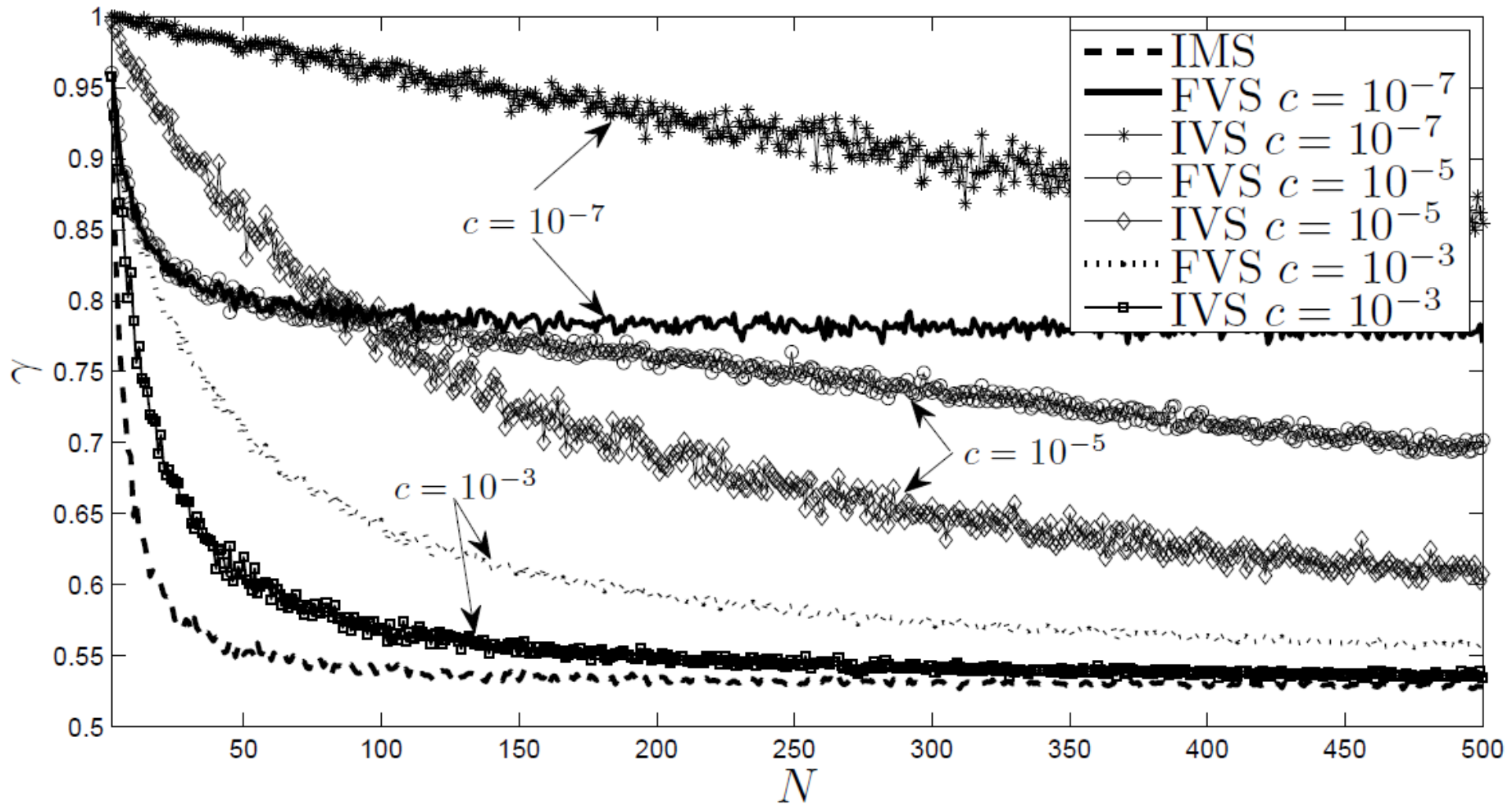
i.e., the case in which some members are *not* within the communication range of the GL.

Metric γ : efficiency is defined as:

$$\gamma = \frac{|Q|k^* + (N - k^*)}{|Q|k^* + \sum_{t=k^*}^N I_t} \quad \gamma_{\max} \Leftrightarrow (N - k^*) = \sum_{t=k^*}^N I_t$$

Low γ indicates that the GVS **improperly** validated Q

The efficiency



Conclusions



There is **a trade-off** between:

- *validating* a group rapidly, thus, achieving low communication load (low ε) **and**
- *delaying* the validation decision for being certain on concluding on group persistence (high γ)



Thank you!