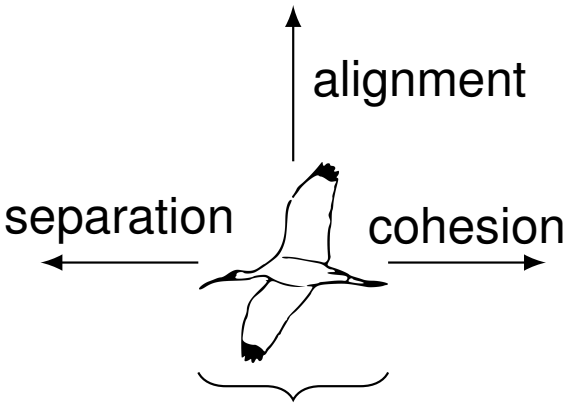
Imperial College London

From Communicating Machines to Graphical Choreographies

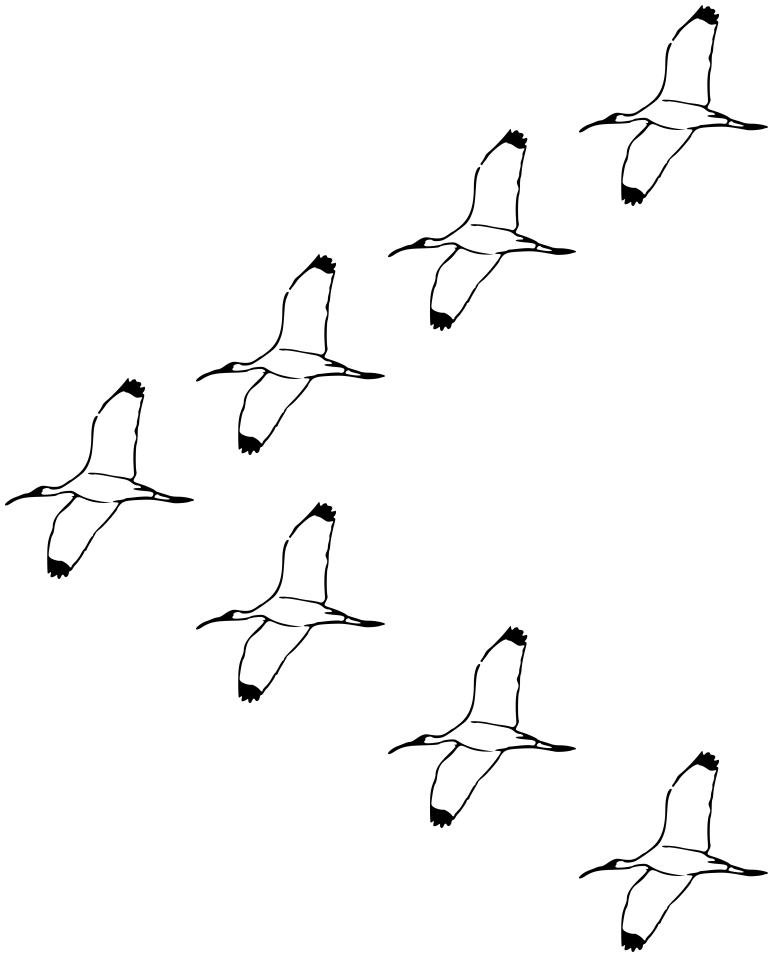
Julien Lange

joint work with Emilio Tuosto *and* Nobuko Yoshida

April 2015



Single bird $\simeq$ local behaviour $\simeq$ CFSM

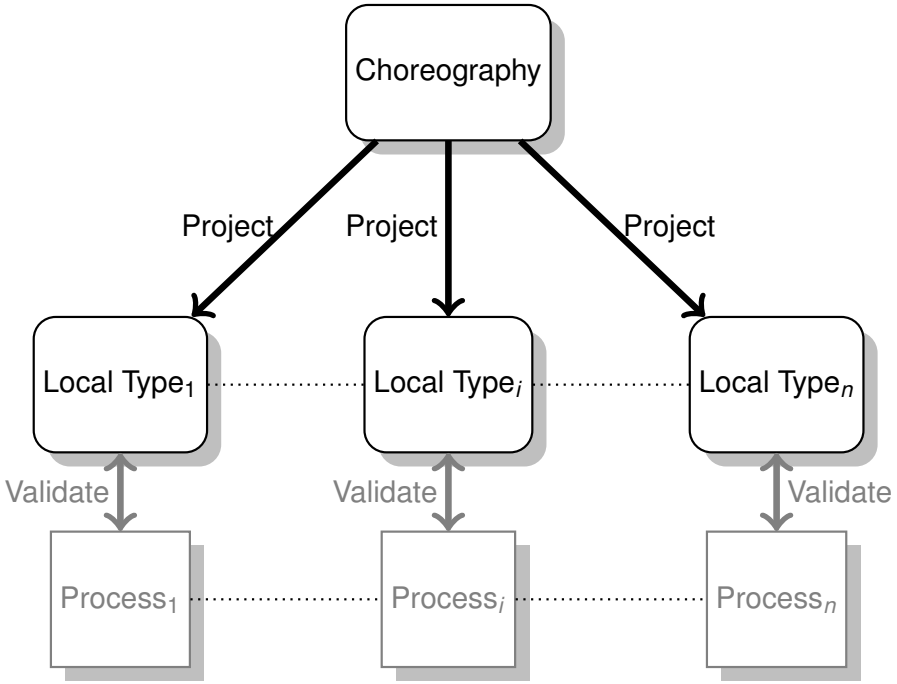


Flock $\simeq$ global behaviour $\simeq$ choreography

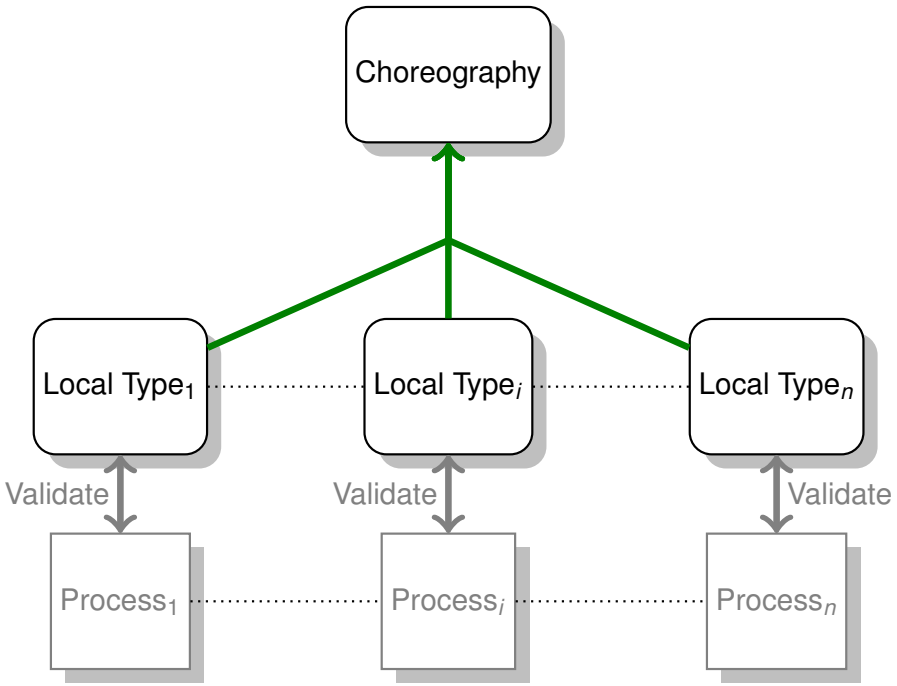
Introduction

- Parts of distributed systems change/evolve, not always in a coordinated way,
- these changes are often *not* documented.
- Service oriented systems are sometimes composed dynamically,
- it is often unclear how complex the overall system has become.
- Cognizant's Zero Deviation Lifecyle Business Unit.

A *global* point of view of a distributed system is *essential* for top-level management.

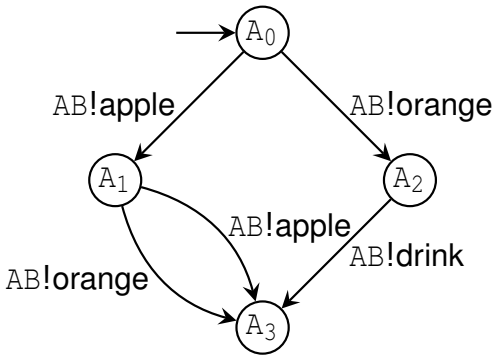


- Choreography-driven development, cf. Multiparty Session Types top-down approach (POPL'08 & ESOP'12)
- Not applicable without *a priori knowledge* of a choreography



- Choreography-driven development, cf. Multiparty Session Types top-down approach (POPL'08 & ESOP'12)
- Not applicable without *a priori knowledge* of a choreography
- Our goal: from *Communicating Finite-State Machines* to *Global Graphs*

Background: CFSMs



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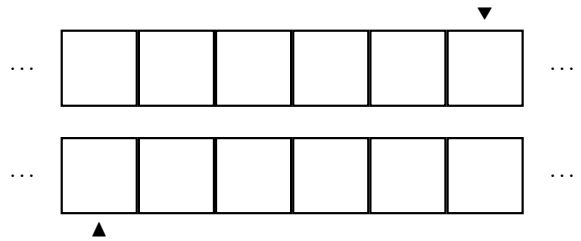
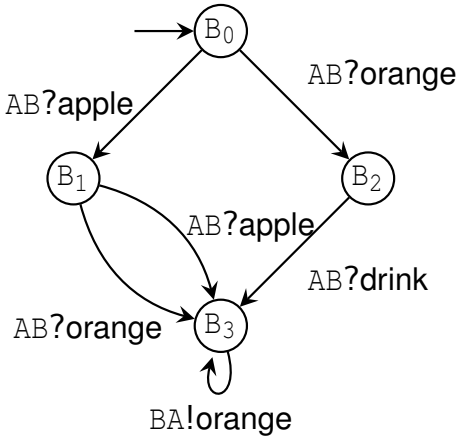
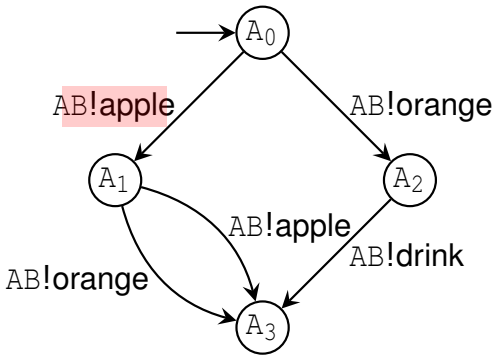




B

"On Communicating Finite-State Machines", Brand & Zafiropulo ('83)

Background: CFSMs



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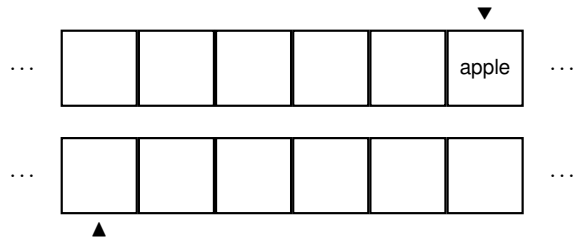

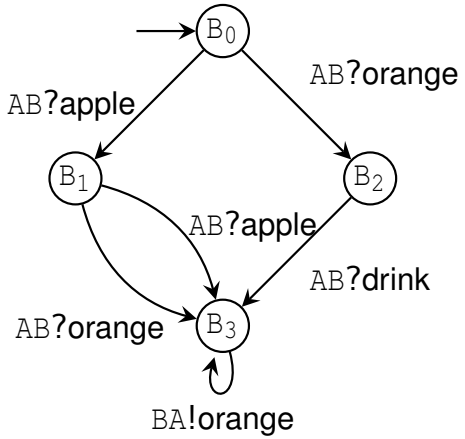

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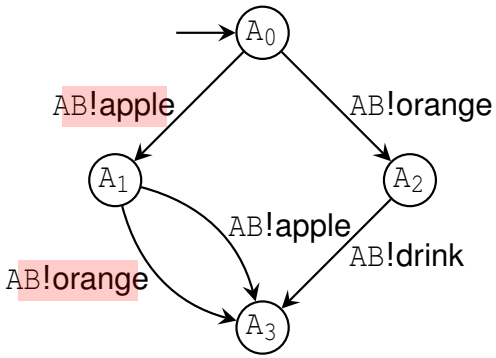

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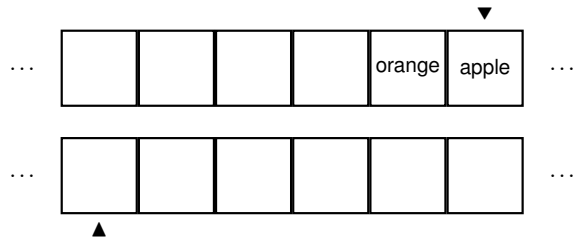

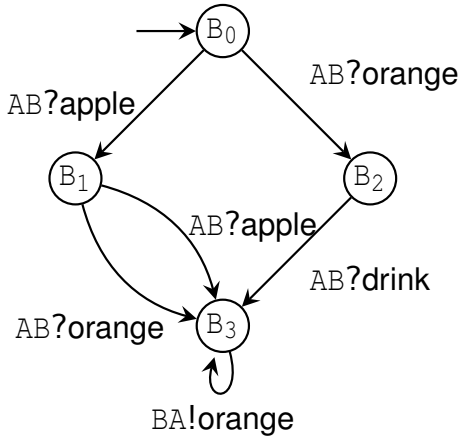

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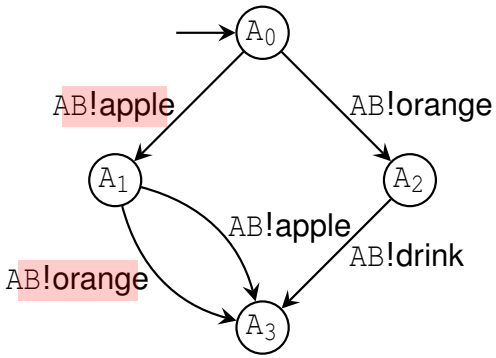

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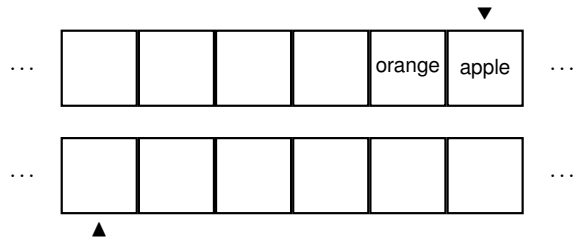

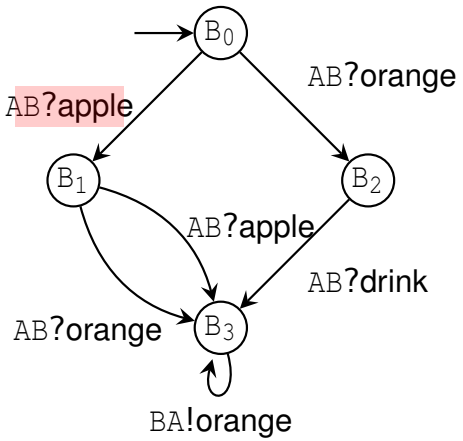

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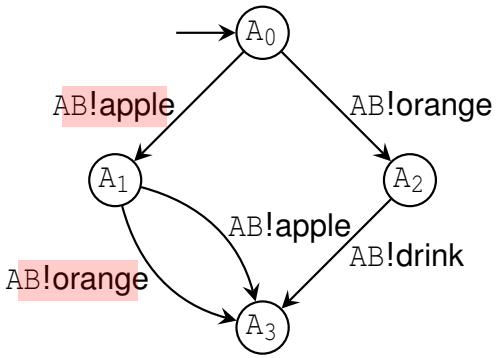

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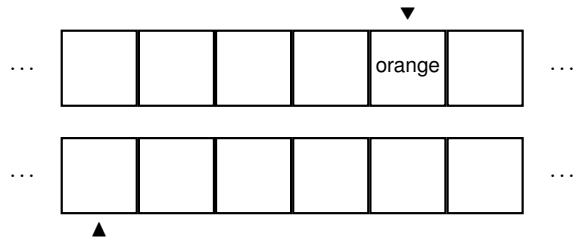

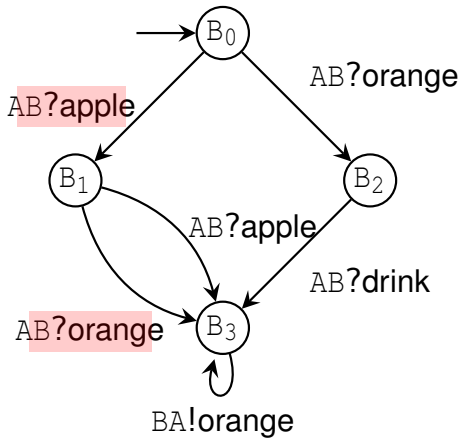

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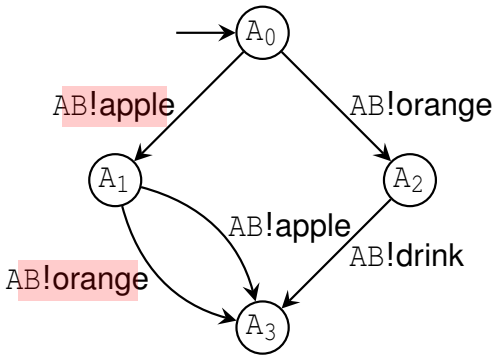

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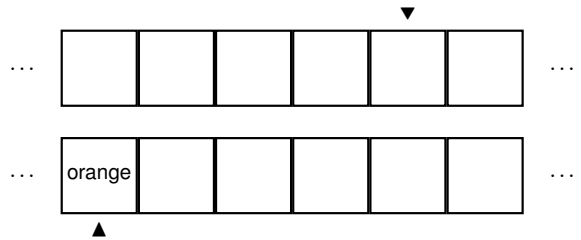

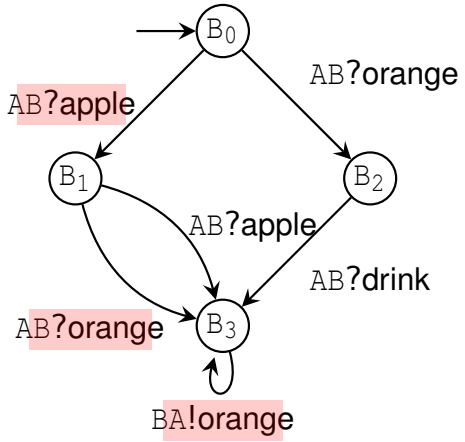

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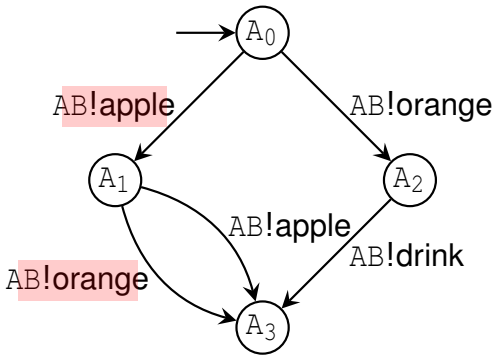

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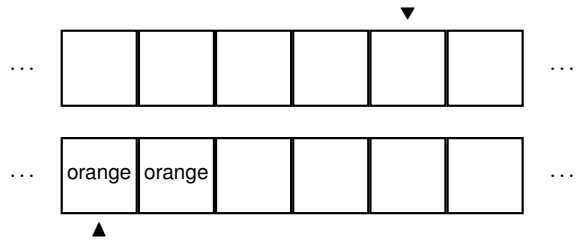

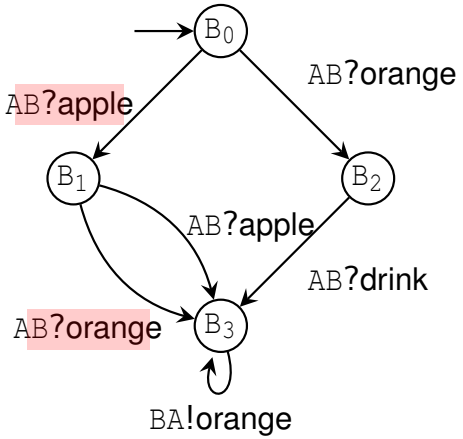

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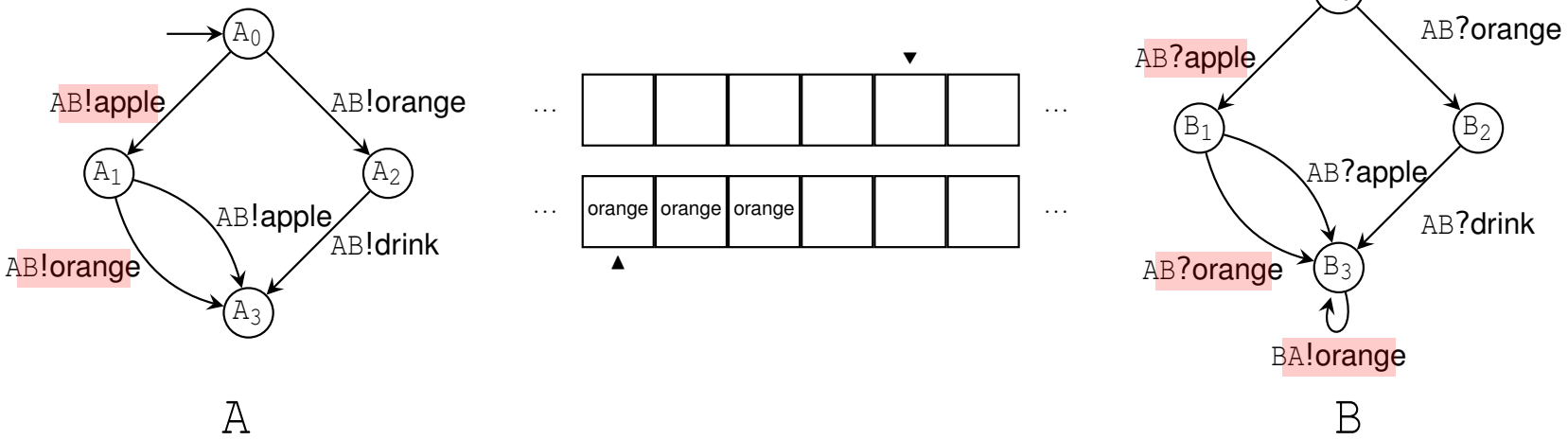

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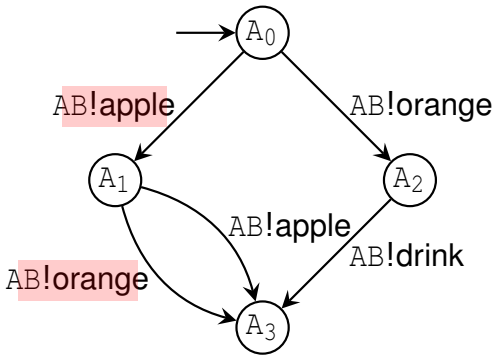

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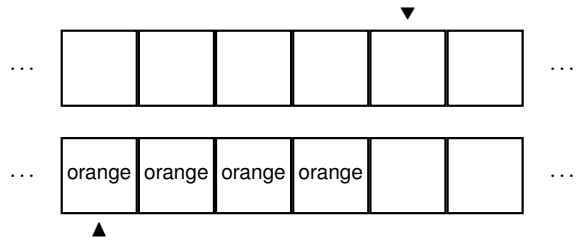

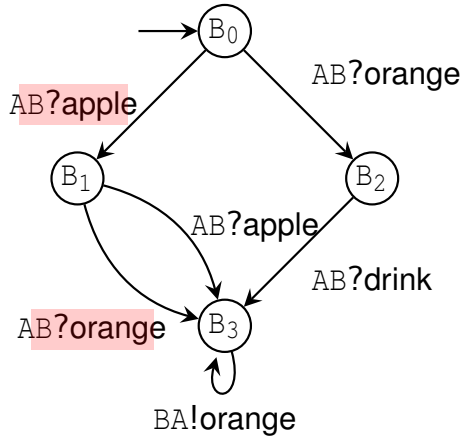

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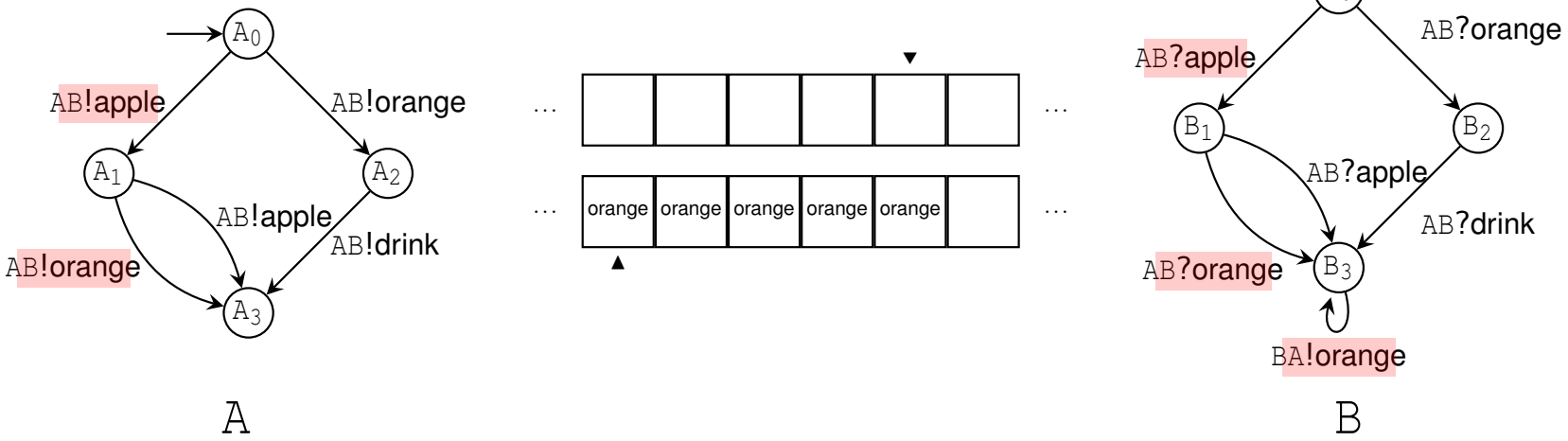

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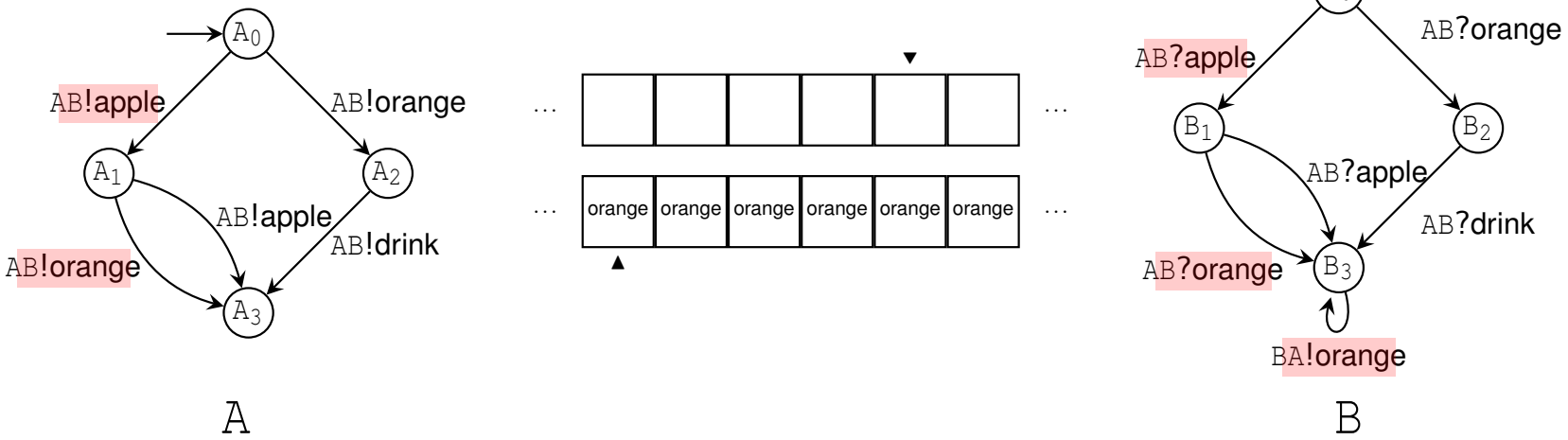

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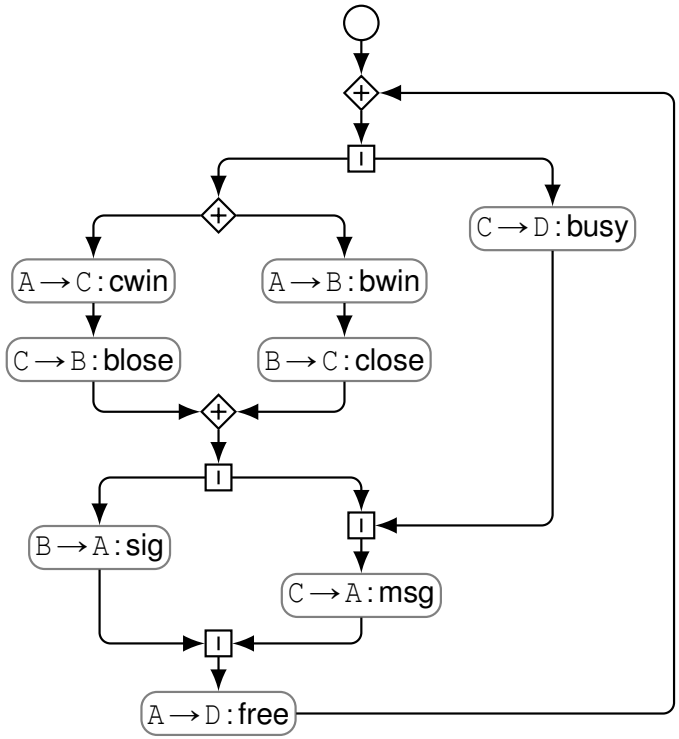
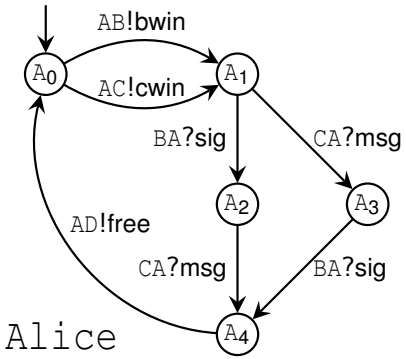
Global Graphs

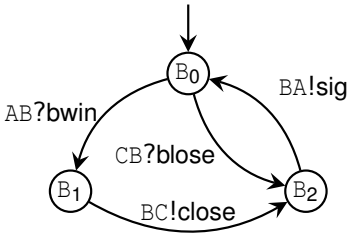


Four Player Game

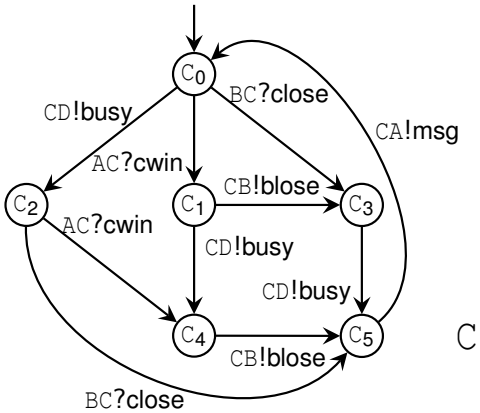
Global Graphs



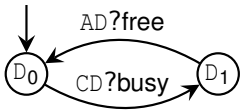
Alice



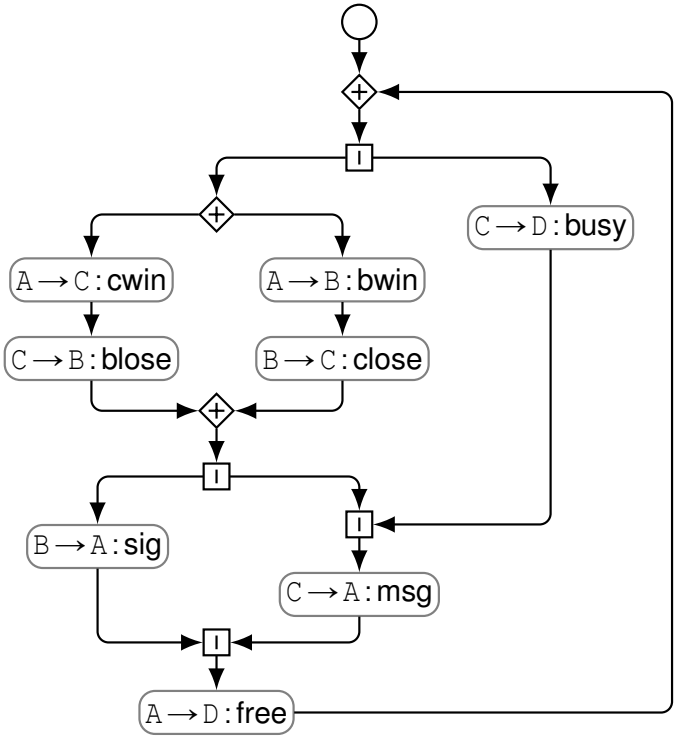
Bob



Carol



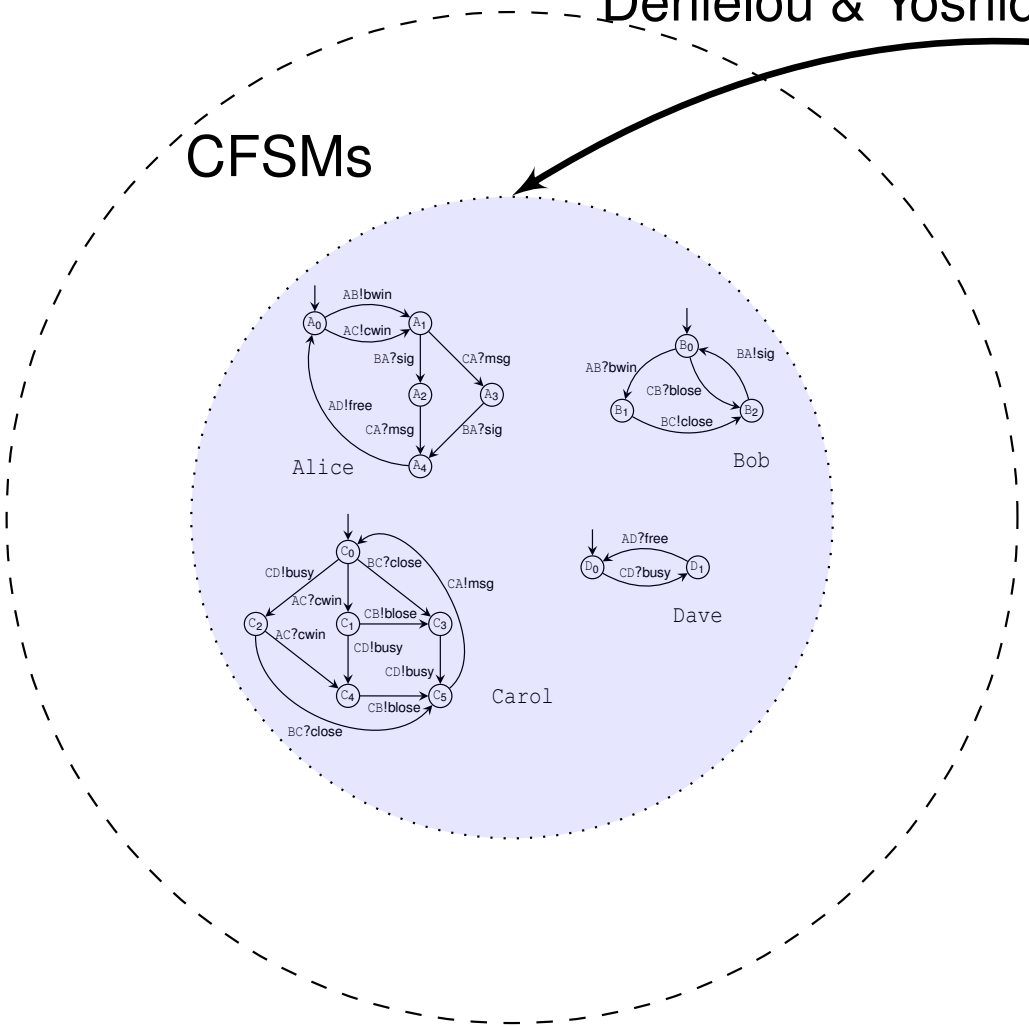
Dave

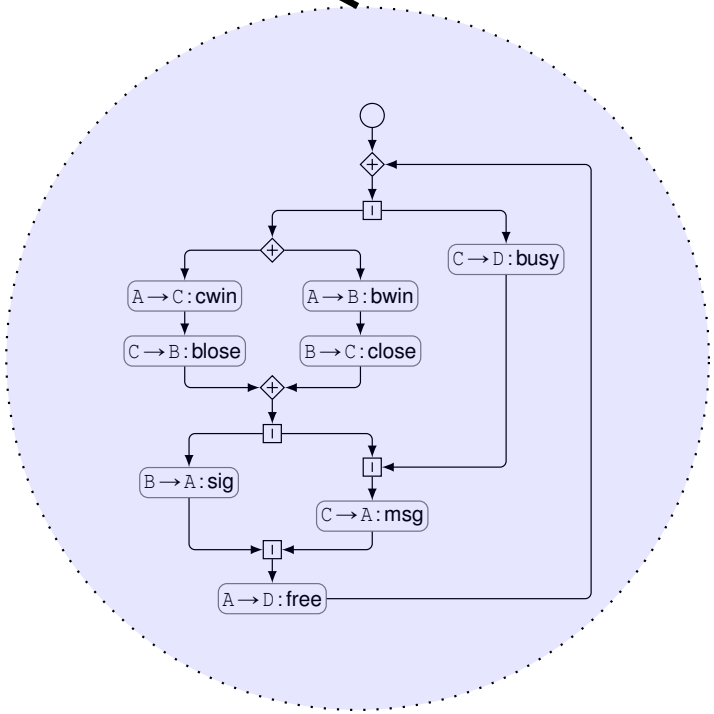


Four Player Game

Deniélou & Yoshida – ESOP'12

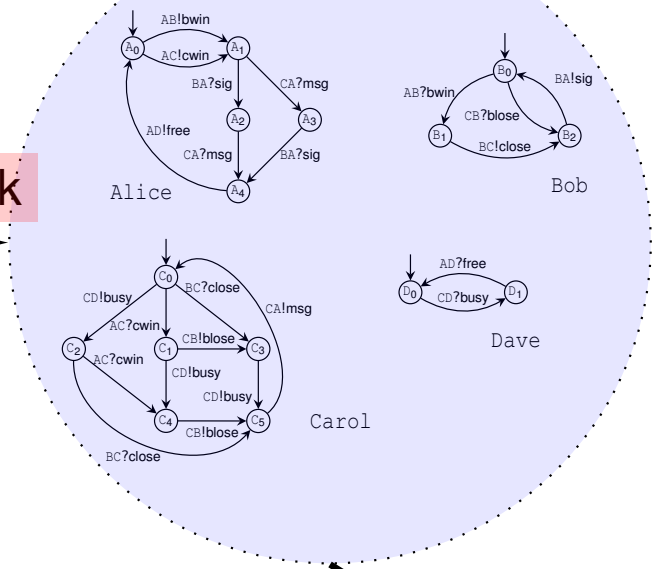
CFSMs

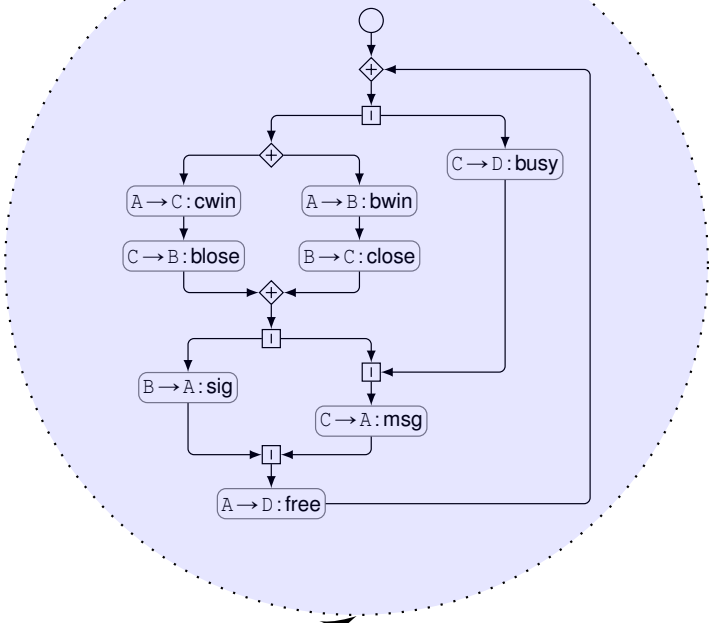




Deniélou & Yoshida – ESOP'12

CFSMs





This work

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Objectives

Two main objectives:

- **Sound Condition for Safety:** generalised multiparty compatibility

  If $S = (M_1, \ldots, M_k)$ is *compatible* then S is "safe", i.e., every sent message is eventually received and no deadlock.

- **Construction of a Global Graph:**

  If G is the global graph constructed from S, then

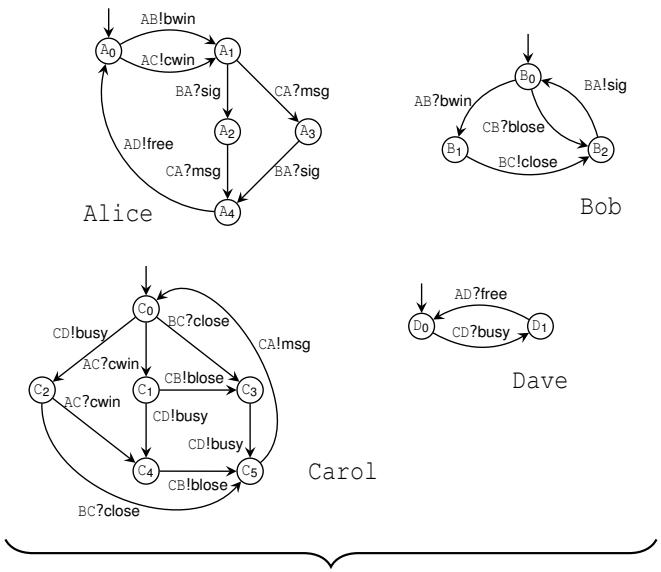
$$S = (M_1, \ldots, M_k) \equiv (G\!\upharpoonright_1, \ldots, G\!\upharpoonright_k)$$

The Plan

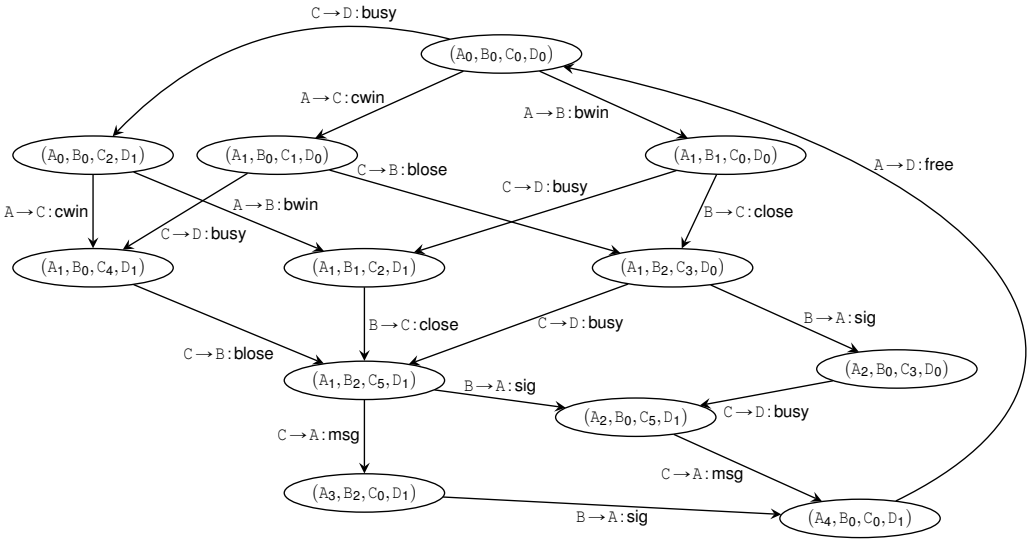
1. Build TS(S), the transition system of all *synchronous* executions
2. Check for safety on TS(S) to
   - ensure equivalence between original system and the projections of the choreography,
   - guarantee safety (no deadlock, no orphan message)
3. Build a choreography (global graph) from TS(S), relying on
   - the theory of regions, and
   - safe Petri nets.

1. Synchronous Transition System of CFSMs

Synchronous Transition System ($TS(S)$)



Four Player Game (CFSMs)



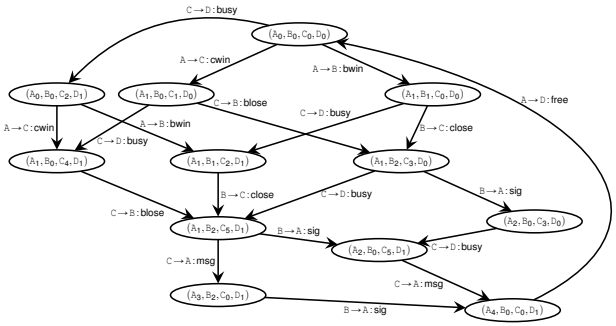
Synchronous Transition System ($TS(S)$)

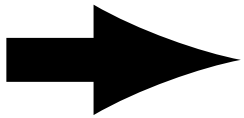
2. **Check for Safety**: *Generalised Multiparty Compatibility* (GMC)

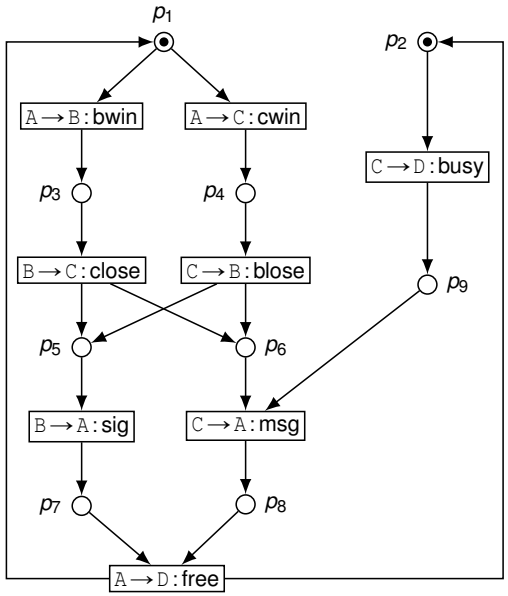
1. Representability: ensures that TS contains enough information
   - Essentially: the projection of TS onto a participant must be time branching bissimilar to the original machine.
2. Branching Property: each branching in TS is “well-formed”
   - it either corresponds to concurrent interactions, or
   - each participant eventually knows which branch was chosen.

3. Build a global graph

1: From TS to Petri Net

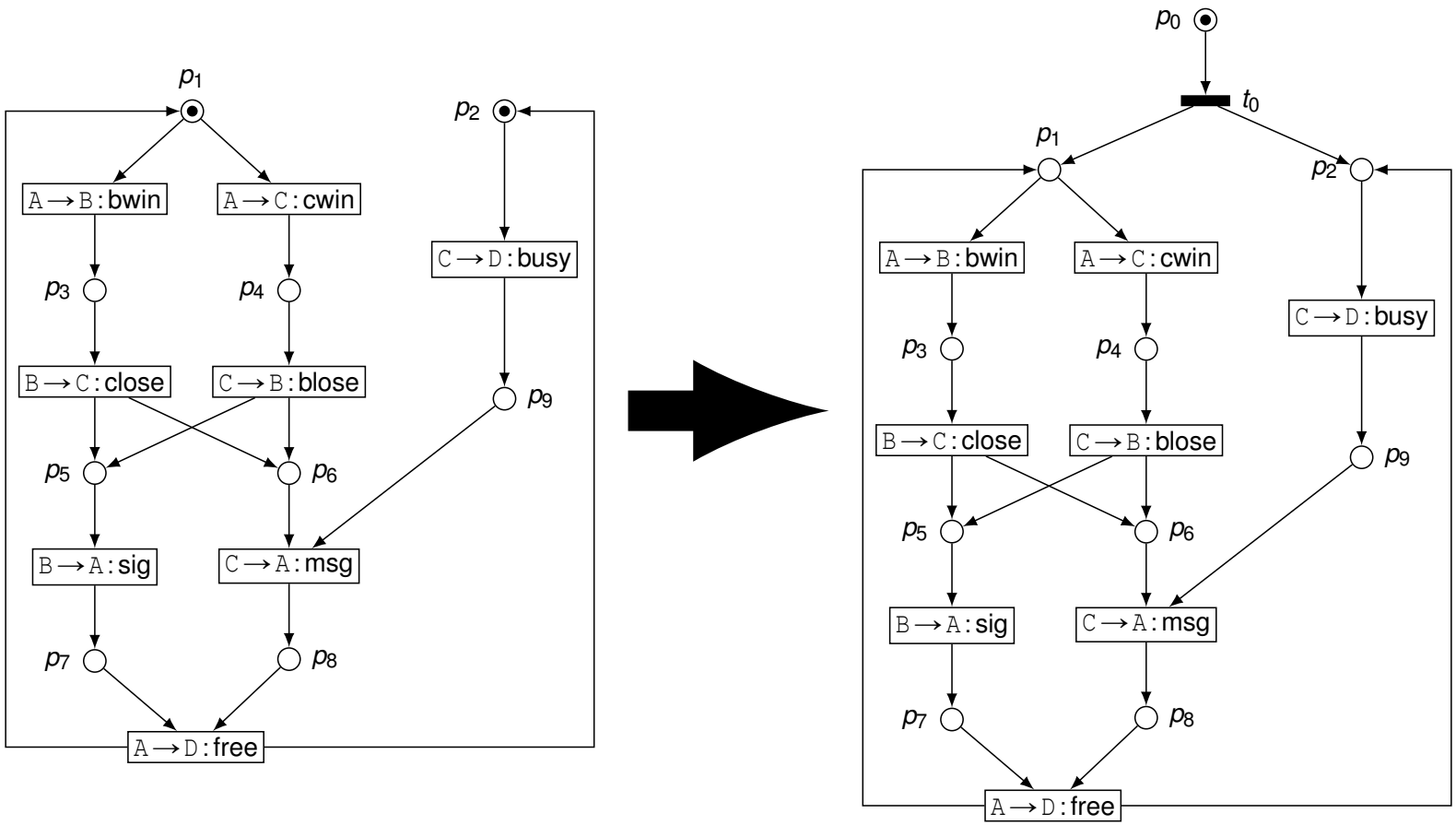




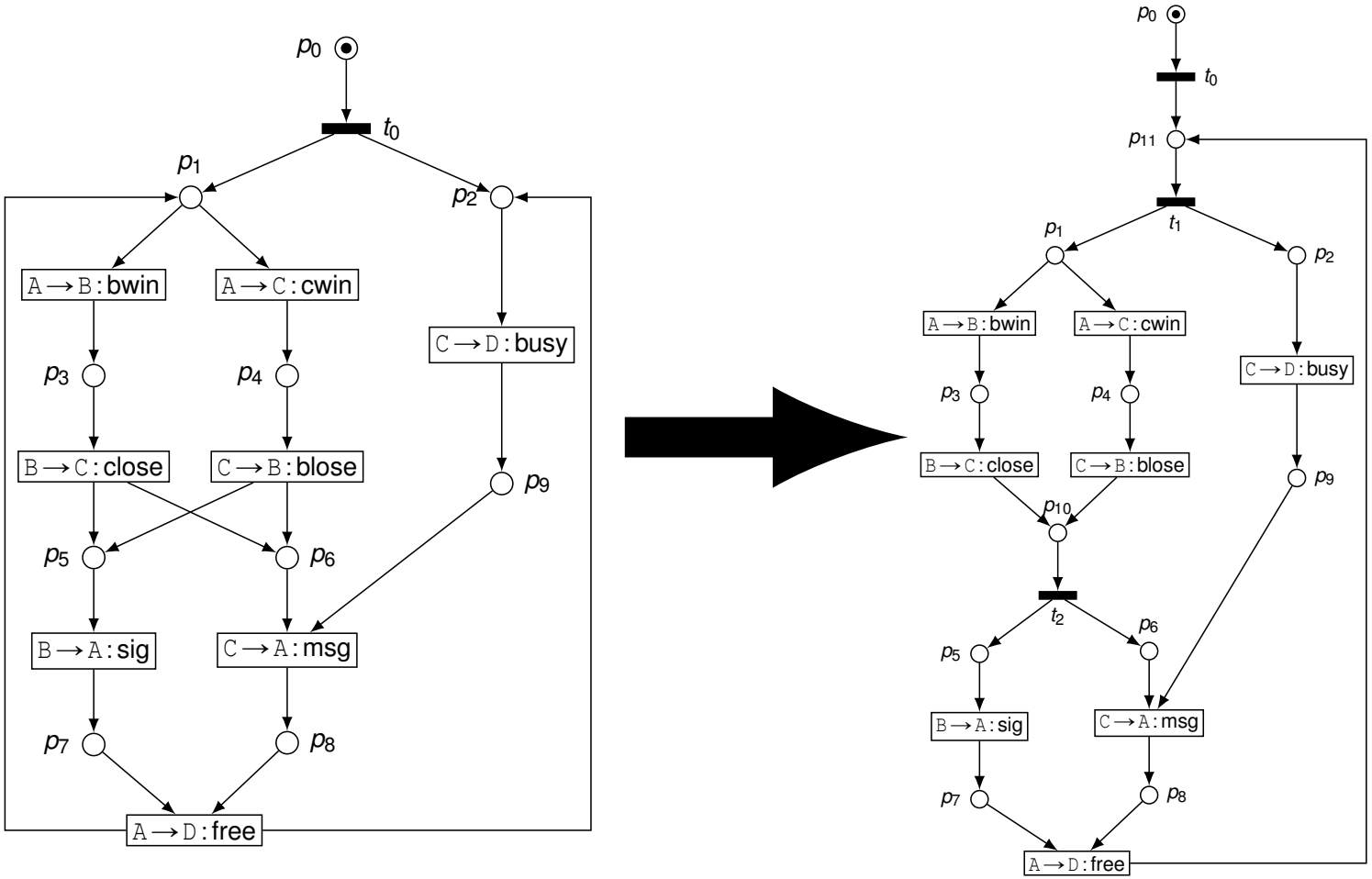


We use the work of Cortadella et al. (1998), based on the **theory of regions**, to derive a *safe* and *extended free-choice* Petri net from the Synchronous Transition System.

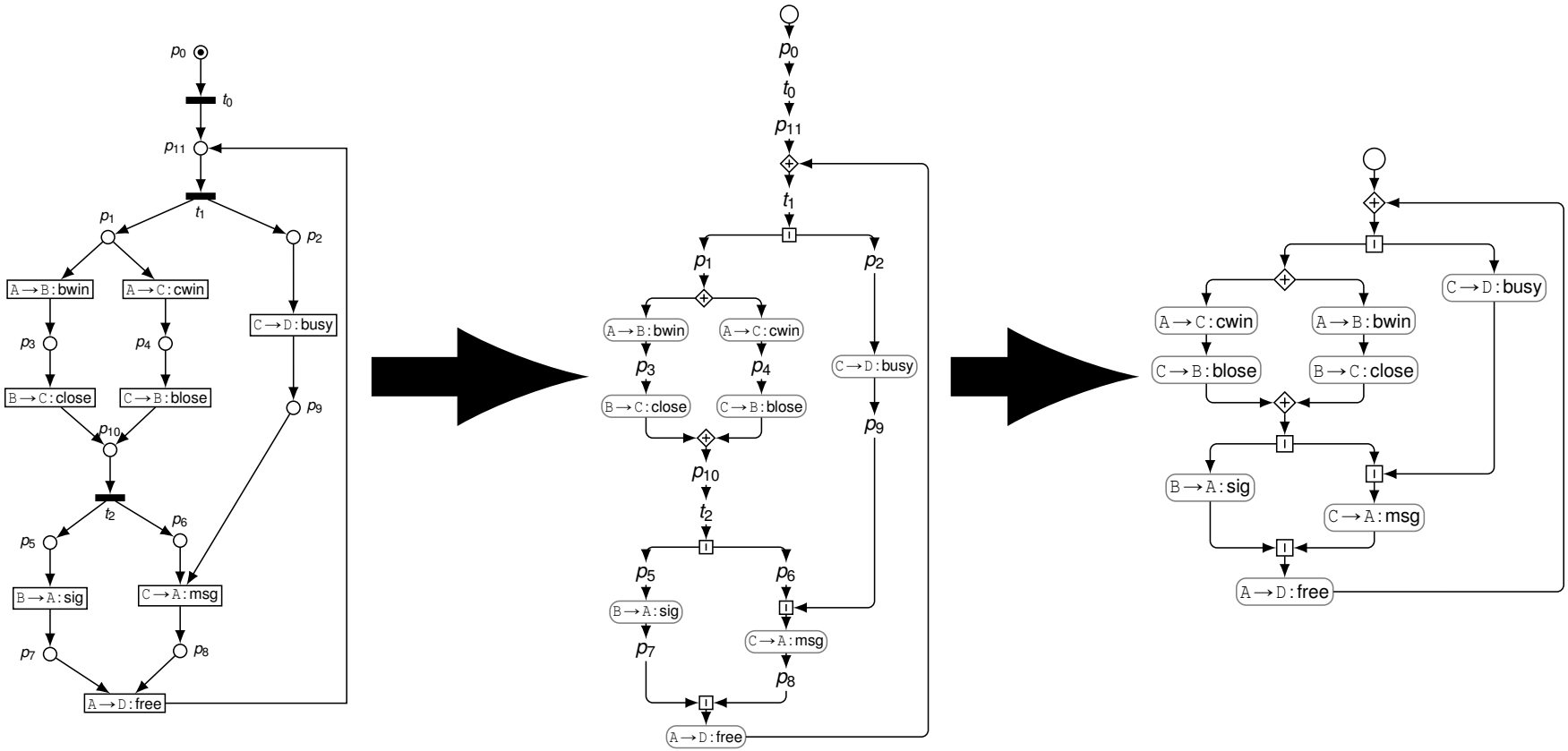
2: From Petri Net to One-source Petri Net



3: From One-source Petri Net to Joined Petri Net



4 & 5 : From Joined Petri Net to Global Graph



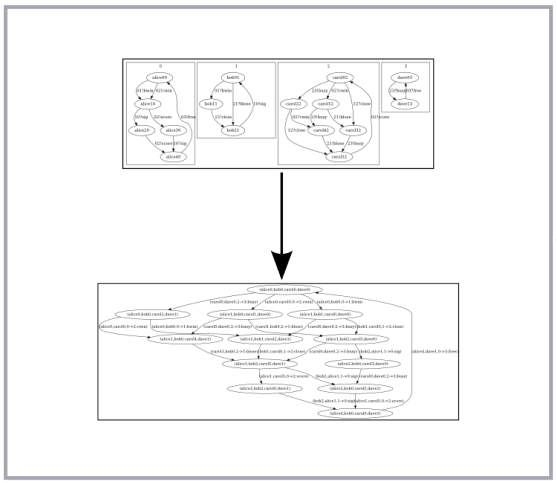
Prototype Implementation

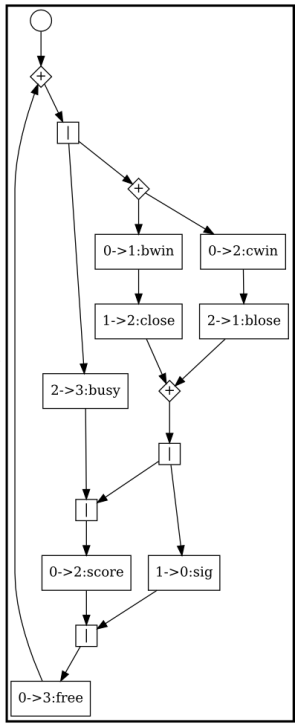
```
.outputs
.state graph
alice0 1 ! bwin alice1
alice0 2 ! cwin alice1
alice1 1 ? sig alice2
alice1 2 ! score alice3
alice2 2 ! score alice4
alice3 1 ? sig alice4
alice4 3 ! free alice0
.marking alice0
.end

.outputs
.state graph
bob0 0 ? bwin bob1
bob1 2 ! close bob2
bob0 2 ? blose bob2
bob2 0 ! sig bob0
.marking bob0
.end

.outputs
.state graph
carol0 1 ? close carol3
carol0 3 ! busy carol2
carol0 0 ? cwin carol1
carol1 3 ! busy carol4
carol1 1 ! blose carol3
carol2 1 ? close carol5
carol2 0 ? cwin carol4
carol3 3 ! busy carol5
carol4 1 ! blose carol5
carol5 0 ? score carol0
.marking carol0
.end

.outputs
.state graph
dave0 2 ? busy dave1
dave1 0 ? free dave0
.marking dave0
.end
```





https://bitbucket.org/julien-lange/gmc-synthesis

Conclusions and Future Work

- Summing up:
  - An effective way of checking properties of CFSMs, and whether one can construct a global graph from them.
  - An algorithm based on the theory of regions.
  - A CFSMs characterisation of well-formed *graphical global types*.
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- Recent extension to support Communicating Timed Automata.

Thanks!

Any questions?

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