



Scalable Persistent Storage for Erlang: Theory and Practice

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Outline

- What does this research aim to do as a part of the RELEASE Project?
- General principles of scalable DBMSs
- NoSQL DBMSs for Erlang
- Measuring the Reliability of Riak
- Scalability of Riak in Practice
- Investigating the scalability of distributed Erlang
- Conclusion & Future work



RELEASE project

• RELEASE is an European project aiming to scale Erlang onto commodity architectures with 100000 cores.





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Scaling Erlang

Scaling Erlang in three levels:

- Distributed Erlang (SD Erlang)
- In-memory Data Structure (ETS table)
- Scalable Persistent Storage for Erlang

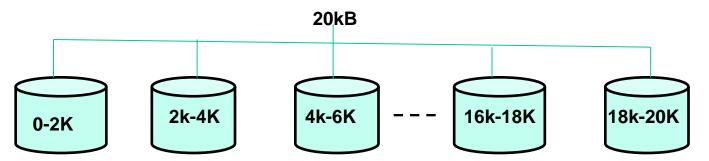


General principles of scalable DBMSs

Data Fragmentation

- 1. Decentralized model (e.g. P2P model)
- 2. Systematic load balancing (make life easier for developer)
- 3. Location transparency

e.g. 20k data is fragmented among 10 nodes

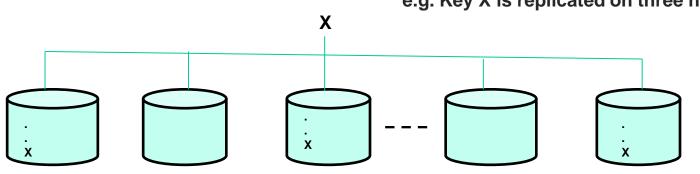




General principles of scalable DBMSs

Replication

- 1. Decentralized model (e.g. P2P model)
- 2. Location transparency
- 3. Asynchronous replication (write is considered complete as soon as on node acknowledges it)

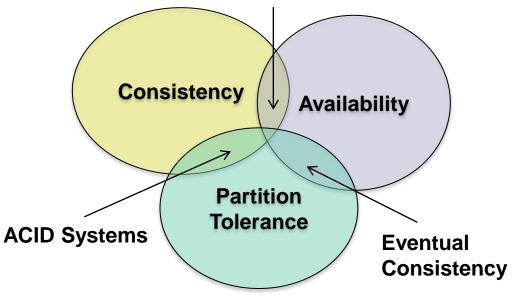






General principles of scalable DBMSs

Not achievable because network failures are inevitable



CAP theorem: cannot simultaneously guarantee:

•Partition tolerance: system continues to operate despite nodes can't talk to each other

•Availability: guarantee that every request receives a response

•Consistency: all nodes see the same data at the same time

Solution: Eventual consistency and reconciling conflicts via data versioning

ACID=Atomicity, Consistency, Isolation, Durability



NoSQL DBMSs for Erlang

	Mnesia	CouchDB	Riak	Cassandra
Fragmentation	•Explicit placement •Client-server •Automatic by using a hash function	•Explicit placement •Multi-server •Lounge is not part of each CouchDB node	 Implicit placement Peer to peer Automatic by using consistent hash technique 	 Implicit placement Peer to peer Automatic by using consistent hash technique
Replication	•Explicit placement •Client-server ? •Asynchronous (Dirty operation)	 Explicit placement Multi-server Asynchronous 	 Implicit placement Peer to peer Asynchronous 	 Implicit placement Peer to peer Asynchronous
Partition tolerance	•Strong consistency	•Eventual consistency •Multi-Version Concurrency Control for reconciliation	•Eventual consistency •Vector clocks for reconciliation	•Eventual consistency •Use timestamp to reconcile
Storage limitation	•The largest possible Mnesia table is 4Gb	•No limitation	 Bitcask has memory limitation LevelDB has no limitation 	•No limitation



Initial Evaluation Results



Initial Evaluation

- Mnesia
- CouchDB
- Riak
- Cassendra

Scalable persistent storage for SD Erlang can be provided by Riak (or Cassandra)

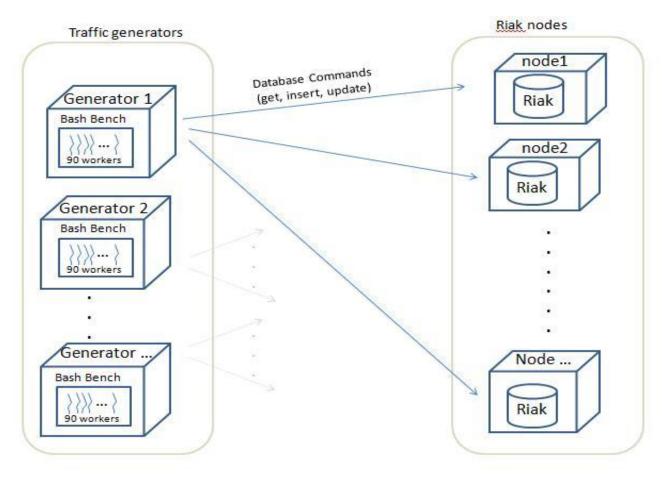


Availability and Scalability of Riak in Practice

- Basho Bench, a benchmarking tool for Riak
- We use Basho Bench on 348-node Kalkyl cluster
- How does Riak cope with node failure? (Availability)
- How adding more Riak nodes affect the throughput? (Scalability)
- There are two kinds of nodes in a cluster:
 - Traffic generators
 - Riak nodes



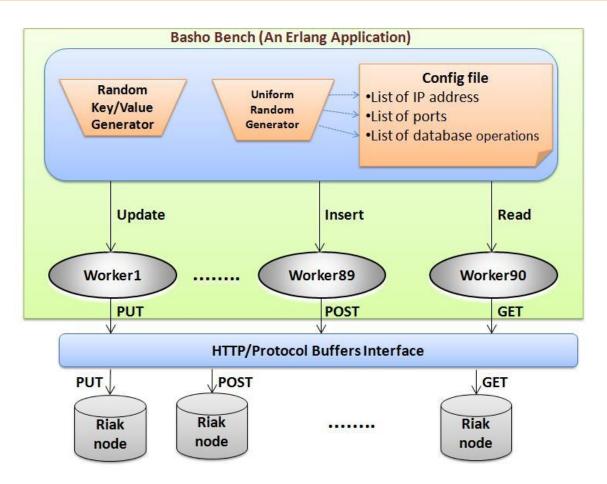
Node Organisation



We use one traffic generator per 3 Riak nodes

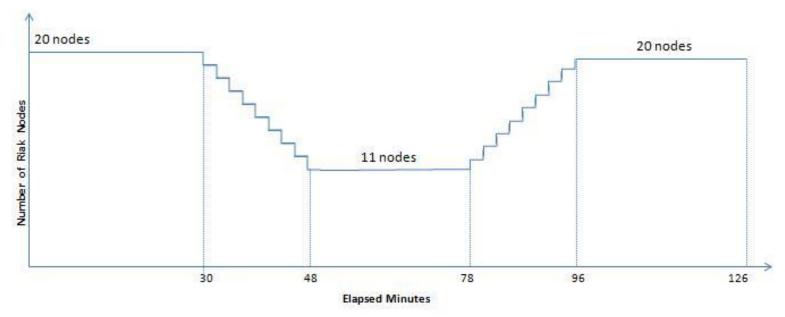


Traffic Generator





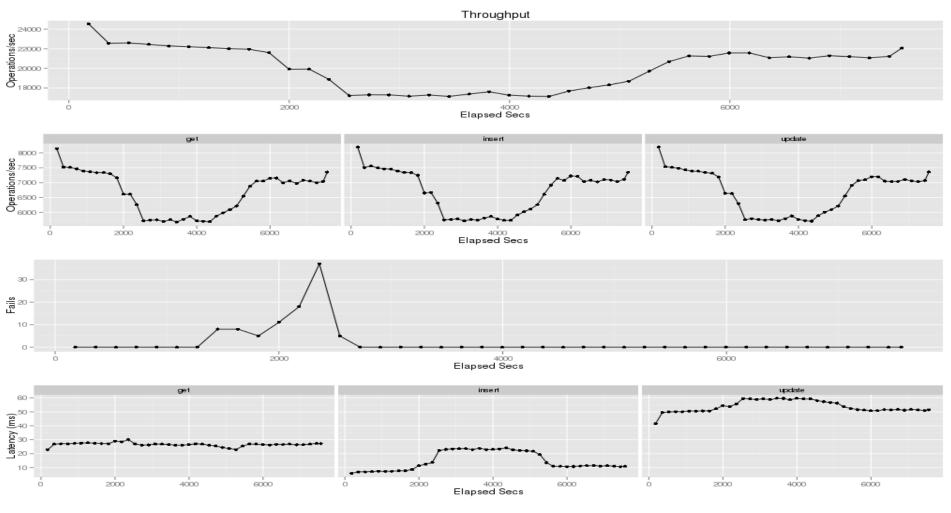
Riak Availability



Time-line shows Riak cluster losing nodes



Riak Availability



How Riak deals with failures



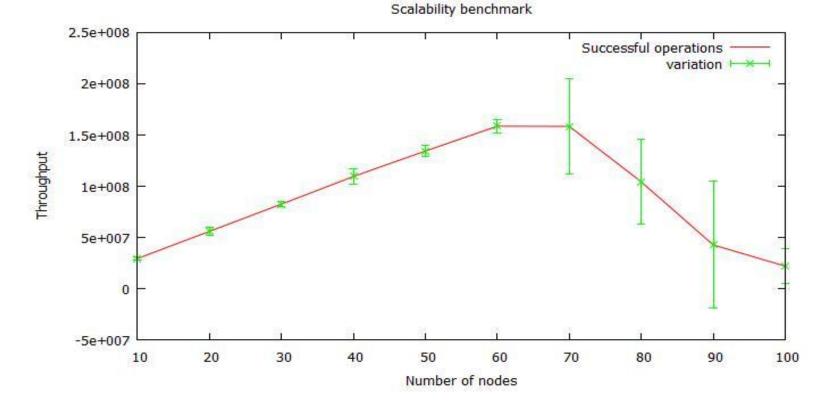
Observation

- Number of failures (37)
- Number of successful operations (approximately 3.41 million)

•When failed nodes come back up, the throughput has grown which means Riak has a good elasticity.



Riak Scalability

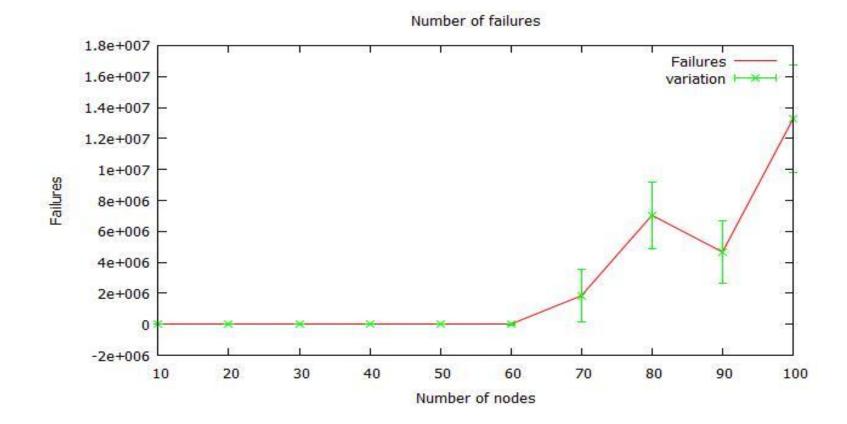


Benchmark on 348-node Kalkyl cluster at Uppsala University

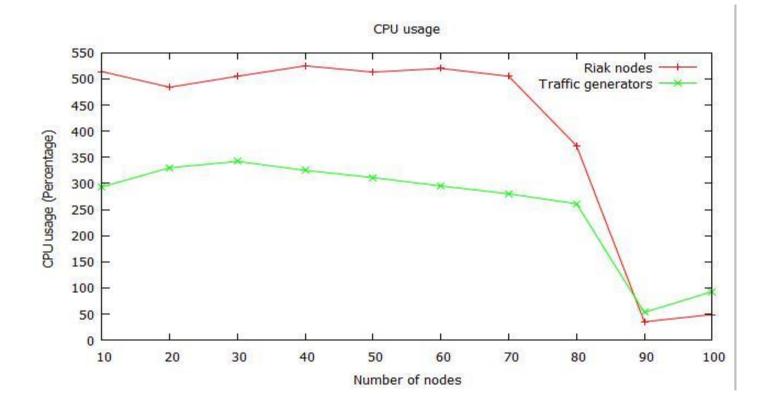


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Failure



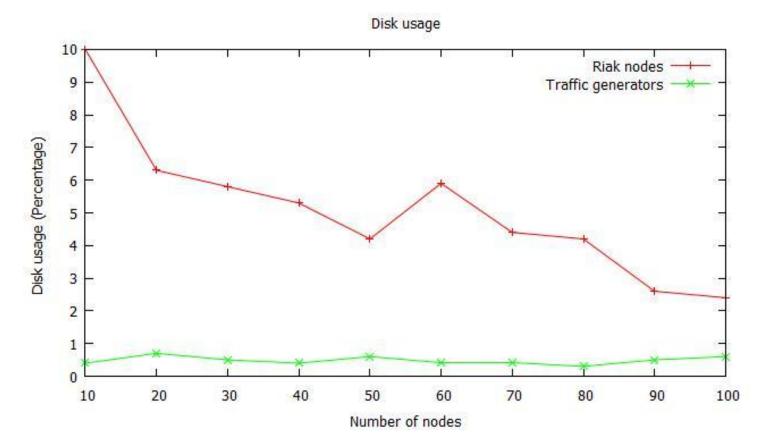




CPU Usage



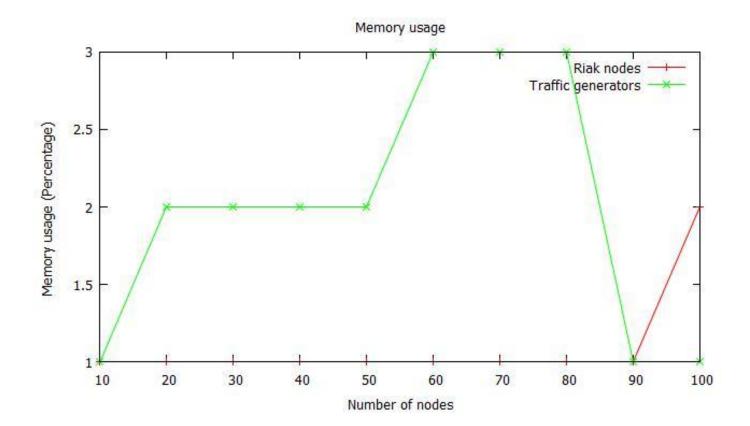
Profiling DISK



DISK Usage



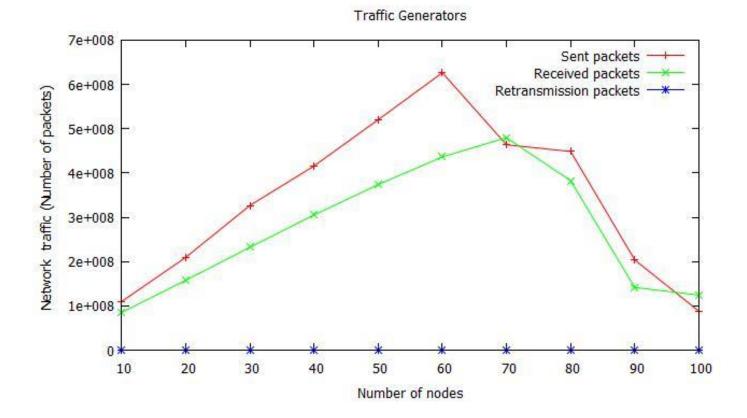
Profiling RAM



Memory Usage



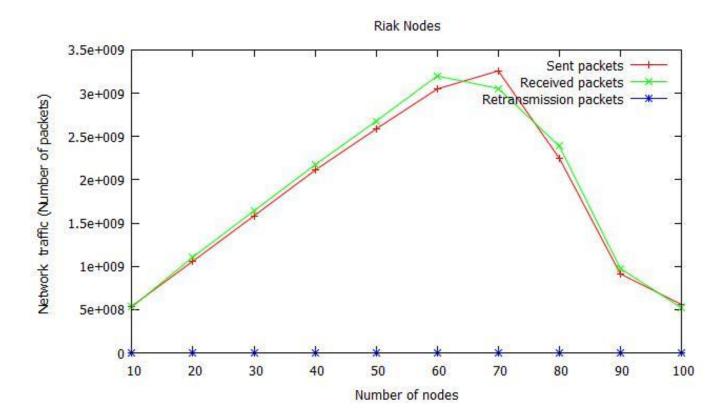
Profiling-Network (Generator)



Network Traffic of Generator Nodes



Profiling-Network (Riak)



Network Traffic of Riak Nodes



Bottleneck for Riak Scalability

The results of profiling *CPU*, *RAM*, *Disk*, and *Network* reveal that they can't be bottleneck for Riak scalability.

Is Riak scalability limits due to limits in distributed Erlang? To find it, we need to measure the scalability of distributed Erlang.



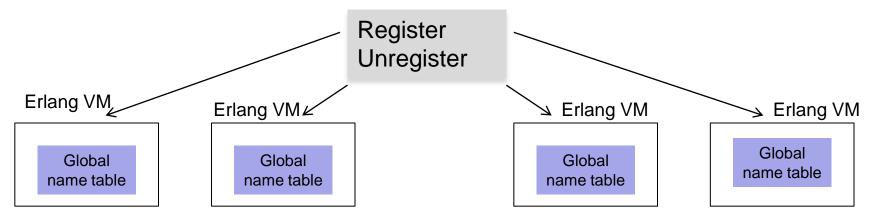
DEbench

- We design Debench for measuring the scalability of distributed Erlang
- Based on Basho Bench
- Measures the Throughput and Latency of Distributed Erlang commands



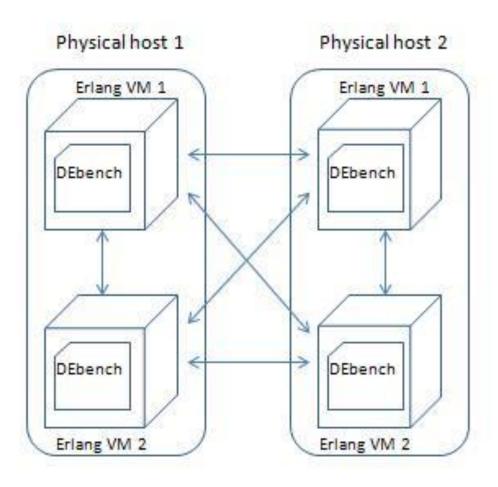
Distributed Erlang Commands

- Spawn: a peer to peer command
- •*register_name* : global name tables located on every node
- unregister_name : global name tables located on every node
- where is_name : a lookup in the local table



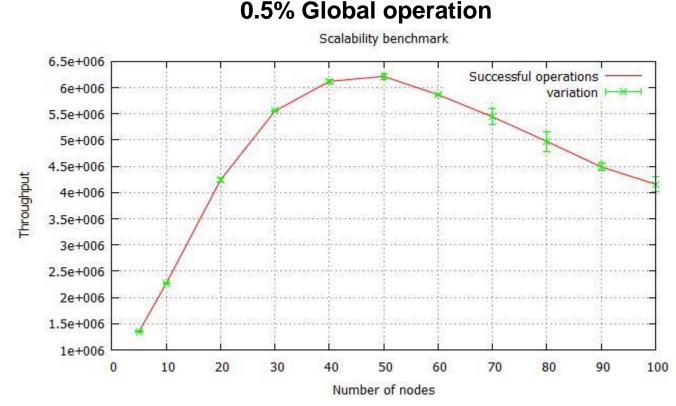


DEbench P2P Nodes





Scalability of Distributed Erlang

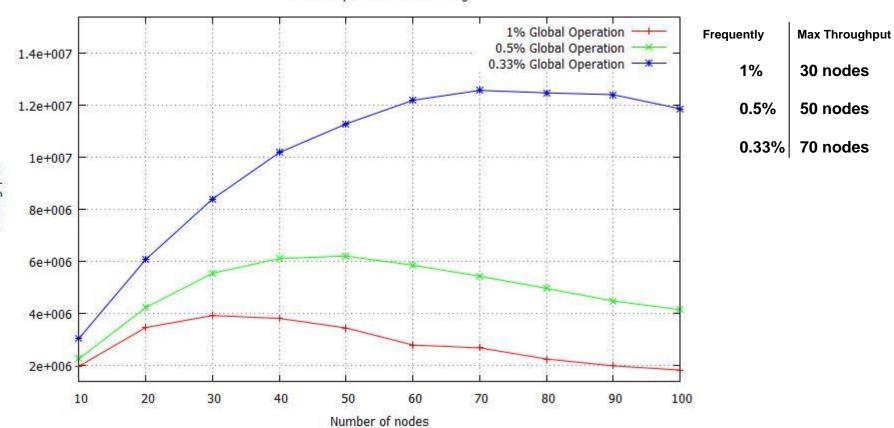


Throughput peaks at 50 nodesLittle improvement beyond 40 nodes



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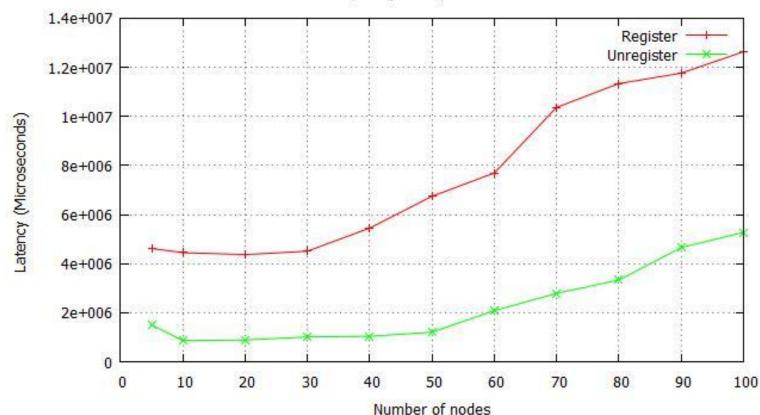
Frequency of Global Operation



Scalability of Distributed Erlang



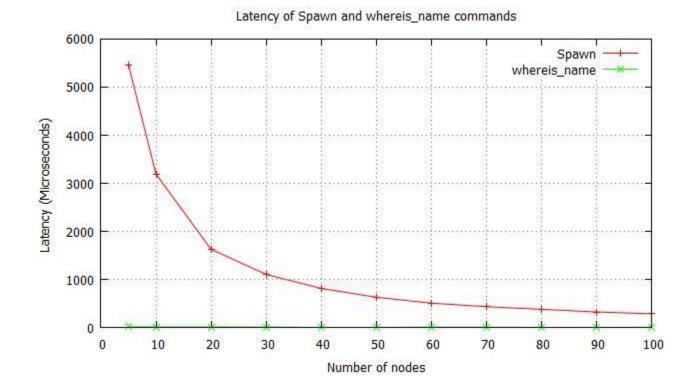




Latency for global update commands

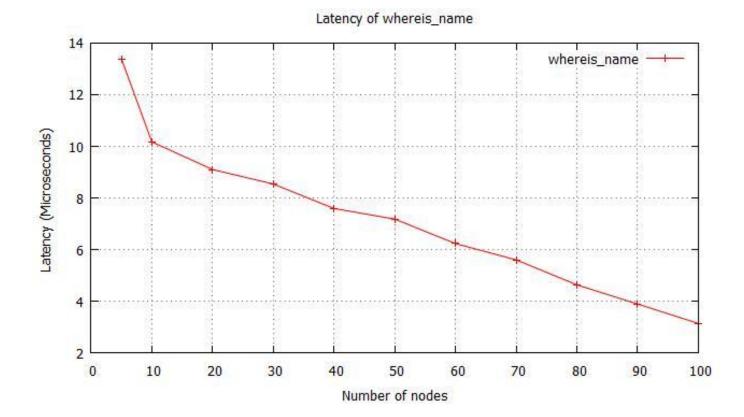
Latency for register and unregister for 2% global update





Latency of spawn





Latency of whereis_name



Conclusion and Future works

•Our benchmark confirms that Riak is highly available and fault-tolerant.

- •We have discovered the scalability limits of Riak is ~60 nodes
- •Global operation limits the scalability of distributed Erlang.

 In the RELEASE, we are working to scale up Distributed Erlang by designing and implementing Scalable Distributed Erlang (SD Erlang)



Thank you!

