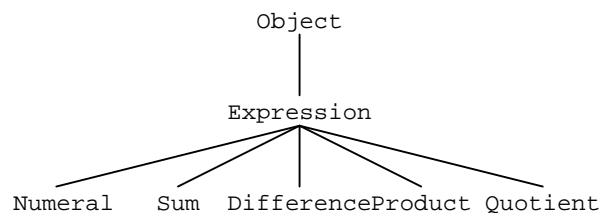


Solutions to Exercises in Chapter 14

14.4 The class hierarchy of a Java program, reflecting the subclass relationship between classes, can be represented by a tree.

- (a) The root node of the class hierarchy tree corresponds to the `Object` class.
- (b) The class hierarchy a tree because Java enforces single inheritance, i.e., each class (except `Object`) has exactly one superclass.
- (c) The class hierarchy for Program 14.12 is:



- (d) If Java interfaces are included, the 'hierarchy' no longer a tree because a class may implement any number of interfaces.

14.5 An implementation of ordered trees is outlined in Program S14.1.

```
public class LinkedOrderedTree implements Tree {  
    // Each LinkedOrderedTree object is an ordered tree whose  
    // elements are arbitrary objects.  
  
    // This tree is represented by a reference to its root node (root), which is  
    // null if the tree is empty. Each tree node contains links to its first and last  
    // children, to its parent, and to its next sibling.  
    private LinkedOrderedTree.Node root;  
  
    //////////// Constructor ////////////  
  
    public LinkedOrderedTree () {  
        // Construct a tree, initially empty.  
        root = null;  
    }  
  
    //////////// Accessors ////////////  
  
    ...  
  
    //////////// Transformers ////////////  
  
    public void makeRoot (Object elem) {  
        // Make this tree consist of just a root node containing element elem.  
        root = new LinkedOrderedTree.Node(elem);  
    }  
}
```

Program S14.1 Outline implementation of ordered trees using linked data structures
(continued on next page).

```

public Tree.Node addChild (Tree.Node node,
                          Object elem) {
// Add a new node containing element elem as the last child of node in
// this tree, and return the new node. The new node has no children of its
// own.
    LinkedOrderedTree.Node parent =
        (LinkedOrderedTree.Node)node;
    LinkedOrderedTree.Node newChild =
        new LinkedOrderedTree.Node(elem);
    newChild.parent = parent;
    if (parent.firstChild == null)
        parent.firstChild = newChild;
    else
        parent.lastChild.nextSib = newChild;
    parent.lastChild = newChild;
    return newChild;
}

public void remove (Tree.Node node) {
// Remove node from this tree, together with all its descendants.
    if (node == root) {
        root = null;
        return;
    }
    LinkedOrderedTree.Node parent = node.parent;
    if (node == parent.firstChild) {
        parent.firstChild = node.nextSib;
        if (parent.firstChild == null)
            parent.lastChild = null;
    } else {
        LinkedOrderedTree.Node prevSib =
            parent.firstChild;
        while (prevSib.nextSib != node)
            prevSib = prevSib.nextSib;
        prevSib.nextSib = node.nextSib;
        if (prevSib.nextSib == null)
            parent.lastChild = prevSib;
    }
}

////////// Iterator //////////

...

////////// Inner class definition for tree nodes //////////

private static class Node implements Tree.Node {
    // Each LinkedOrderedTree.Node object is a node of an
    // ordered tree, and contains a single element.

    // This tree node consists of an element (element), a link to its first
    // and last children (firstChild, lastChild) a link to its parent
    // (parent), and a link to its next sibling (nextSib).
    private Object element;
    private LinkedOrderedTree.Node firstChild,
        lastChild, parent, nextSib;
    ...
}
}

```

Program S14.1 Outline implementation of ordered trees using linked data structures
(continued).

14.6 The following methods visit, in pre-order, all of the nodes in a given tree:

```
static void preOrderTraverse (Tree tree) {
    if (tree.root() != null)
        preOrderTraverseSubtree(tree, tree.root());
}

static void preOrderTraverseSubtree (Tree tree,
    Tree.Node parent) {
    ... // Visit parent.
    Iterator children = tree.children(parent);
    while (children.hasNext()) {
        Tree.Node child =
            (Tree.Node)children.next();
        preOrderTraverseSubtree(tree, child);
    }
}
```

14.8 Methods to visit, in *post-order*, all of the nodes in a given tree would be similar to the methods of Exercise 14.6, except that the code to visit parent must *follow* the while-loop that traverses the children.

14.10 To visit the nodes of *tree* in depth order:

1. Make *node-queue* contain only the root node of *tree*.
2. While *node-queue* is nonempty, repeat:
 - 2.1. Remove the front element of *node-queue* into *node*.
 - 2.2. Visit *node*.
 - 2.3. Add all the children of *node* to the rear of *node-queue*.
3. Terminate.

Implementation (using the `java.util.LinkedList` representation of the node queue):

```
static void depthOrderTraverse (Tree tree) {
    LinkedList nodeQueue = new LinkedList();
    nodeQueue.addLast(tree.root());
    while (! nodeQueue.isEmpty()) {
        Tree.Node node =
            (Tree.Node)nodeQueue.removeFirst();
        ... // Visit node.
        Iterator children = tree.children(node);
        while (children.hasNext()) {
            Tree.Node child =
                (Tree.Node)children.next();
            nodeQueue.addLast(child);
        }
    }
}
```

14.11 An implementation of unordered trees using arrays is outlined in Program S14.2.

The `addChild` operations has time complexity $O(1)$. If c is the maximum number of children per node, the `remove` operation has time complexity $O(c)$.

```

public class ArrayUnorderedTree implements Tree {
    // Each ArrayUnorderedTree object is an unordered tree whose
    // elements are arbitrary objects.

    // This tree is represented by a reference to its root node (root), which is
    // null if the tree is empty. Each tree node contains an array of children.
    private ArrayUnorderedTree.Node root;

    //////////// Constructor ////////////

    public ArrayUnorderedTree () {
        // Construct a tree, initially empty.
        root = null;
    }

    //////////// Accessors ////////////

    public Tree.Node root () {
        // Return the root node of this tree, or null if this tree is empty.
        return root;
    }

    public Tree.Node parent (Tree.Node node) {
        // Return the parent of node in this tree, or null if node is the root node.
        return node.parent;
    }

    public int childCount (Tree.Node node) {
        // Return the number of children of node in this tree.
        ArrayUnorderedTree.Node parent =
            (ArrayUnorderedTree.Node)node;
        return parent.childCount;
    }

    //////////// Transformers ////////////

    public void makeRoot (Object elem) {
        // Make this tree consist of just a root node containing element elem.
        root = new ArrayUnorderedTree.Node(elem);
    }

    public Tree.Node addChild (Tree.Node node,
                               Object elem) {
        // Add a new node containing element elem as a child of node in this
        // tree, and return the new node. The new node has no children of its own.
        ArrayUnorderedTree.Node parent =
            (ArrayUnorderedTree.Node)node;
        ArrayUnorderedTree.Node newChild =
            new ArrayUnorderedTree.Node(elem);
        newChild.parent = parent;
        if (parent.childCount == parent.children.length)
            parent.expand();
        parent.children[parent.childCount++] = newChild;
        return newChild;
    }
}

```

Program S14.2 Outline implementation of unordered trees using arrays
(continued on next page).

```

public void remove (Tree.Node node) {
// Remove node from this tree, together with all its descendants.
    if (node == root) {
        root = null;
        return;
    }
    ArrayUnorderedTree.Node parent = node.parent;
    parent.childCount--;
    int i = 0;
    while (parent.children[i] != node) i++;
    while (i < parent.childCount) {
        parent.children[i] = parent.children[i+1];
        i++;
    }
}

////////// Iterator //////////

...

////////// Inner class definition for tree nodes //////////

private static class Node implements Tree.Node {
    // Each ArrayUnorderedTree.Node object is a node of an
    // unordered tree, and contains a single element.

    // This tree node consists of an element (element), a link to its parent
    // (parent), an array of links to its children (children), and the
    // number of children (childCount).
    private Object element;
    private ArrayUnorderedTree.Node parent;
    private ArrayUnorderedTree.Node[] children;
    private int childCount;

    private Node (Object elem) {
        // Construct a tree node, containing element elem, that has no parent
        // and no children.
        this.element = elem;
        this.parent = null;
        this.children = new ArrayUnorderedTree.Node[4];
        this.childCount = 0;
    }

    ...

    public void expand () {
        // Increase the length of this node's array of links to children.
        ...
    }
}
}

```

Program S14.2 Outline implementation of unordered trees using arrays (*continued*).

- 14.14** In the linked (or array) implementation of an unordered tree, the explicit reference to a node's parent could be removed, but the `parent` operation must then search the tree to find the node's parent. This search can be done by a pre-order traversal, terminating when the parent is found:

```

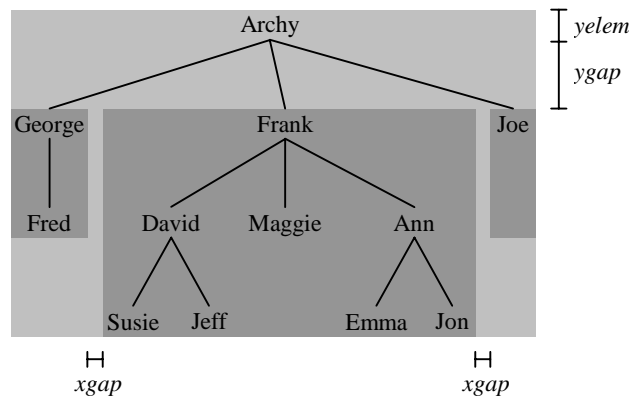
public Tree.Node parent (Tree.Node node) {
    // Return the parent of node in this tree, or null if node is the root
    // node.
    if (root == node)
        return null;
    else
        return findParent(node, root);
}

private Tree.Node findParent (
    Tree.Node node,
    Tree.Node ancestor) {
    // Return the parent of node in this tree, assuming that ancestor
    // is a parent or grandparent or ... of node.
    Iterator children = tree.children(ancestor);
    while (children.hasNext()) {
        Tree.Node child =
            (Tree.Node)children.next();
        if (child == node) return ancestor;
        Tree.Node parent =
            findParent(node, child);
        if (parent != null) return parent;
    }
    return null;
}

```

The parent operation now has time complexity $O(n)$, as does any other operation that must call the parent operation.

14.20 In the drawing of a tree, let each subtree's *bounding rectangle* be the smallest rectangle that encloses all that subtree's nodes. The following example shows a family tree and some of the bounding rectangles:



Here is one simple idea for drawing a tree. Consider a node N and its subtrees. Place the subtrees' bounding rectangles side by side, with their tops aligned, leaving a small gap ($xgap$ above) between neighboring rectangles. Draw node N 's element centered above these rectangles, leaving a small gap ($ygap$ above). Then the bounding rectangle for the tree whose top node is N is the smallest rectangle that encloses node N 's element and all the subtrees' rectangles.

To draw *tree*:

1. Draw the subtree whose topmost node is *tree*'s root, with the top left of its bounding rectangle at $(0, 0)$.
2. Terminate.

To draw the subtree whose topmost node is N , with the top left of its bounding rectangle at (x, y) :

1. Let c be the number of children of node N .
2. If $c = 0$:
 - 2.1. Set $width$ to the width of node N 's element when drawn.
 - 3.4. Set x_{top} to $x + width/2$.
 - 2.2. Draw node N 's element centered at (x_{top}, y) .
3. If $c > 0$:
 - 3.1. Set x_{left} to x , and set y_{child} to $y + y_{elem} + y_{gap}$.
 - 3.2. For $i = 1, \dots, c$, repeat:
 - 3.2.1. Draw the subtree whose topmost node is the i th child of N , with the top left of its bounding rectangle at (x_{left}, y_{child}) , and let its width be w .
 - 3.2.2. Set $x_{child}[i]$ to $x_{left} + w/2$.
 - 3.2.3. Increment x_{left} by $w + x_{gap}$.
 - 3.3. Set $width$ to $x_{left} - x_{gap} - x$.
 - 3.4. Set x_{top} to $x + width/2$.
 - 3.5. Draw node N 's element centered at (x_{top}, y) .
 - 3.6. For $i = 1, \dots, c$, repeat:
 - 3.6.1. Draw a straight line from $(x_{top}, y + y_{elem})$ to $(x_{child}[i], y_{child})$.
4. Terminate with answer $width$.