

# **9** Variables and lifetime

- Variables and storage
- Simple vs composite variables
- Lifetime: global, local, heap variables
- Pointers
- Commands
- Expressions with side effects

9-1



- In functional PLs (and in mathematics), a variable stands for a fixed (but possibly unknown) value.
- In imperative and OO PLs, a variable contains a value. The variable may be inspected and updated as often as desired.
  - Such a variable can be used to model a real-world object whose state changes over time.



- To understand such variables, assume an abstract storage model:
  - A store is a collection of cells, each of which has a unique address.
  - Each cell is either *allocated* or *unallocated*.
  - Each allocated cell contains either a simple value or *undefined*.
  - An allocated cell can be inspected, unless it contains undefined.
  - An allocated cell can be updated at any time.





- A **simple variable** is one that contains a primitive value or pointer.
  - A simple variable occupies a single allocated cell.
- A composite variable is one that contains a composite value.
  - A composite variable occupies a group of adjacent allocated cells.



- A variable of a composite type has the same structure as a value of the same type. For instance:
  - A tuple variable is a tuple of component variables.
  - An array variable is a mapping from an index range to a group of component variables.
- Depending on the PL, a composite variable can be:
  - totally updated (all at once), and/or
  - **selectively updated** (one component at a time).



Declaration and updating of struct variables:

struct Date {int y, m, d;}
struct Date xmas, today;





- Every variable is created at some definite time, and destroyed at some later time when it is no longer needed.
- A variable's lifetime is the interval between its creation and destruction.
- A variable occupies cells only during its lifetime. When the variable is destroyed, these cells may be deallocated.
  - And these cells may subsequently be re-allocated to other variable(s).



- A global variable's lifetime is the program's entire run-time. It is created by a global declaration.
- A local variable's lifetime is an activation of a block. It is created by a declaration within that block, and destroyed on exit from that block.
- A heap variable's lifetime is arbitrary, but bounded by the program's run-time. It can be created at any time (by an allocator), and may be destroyed at any later time. It is accessed through a pointer.



Outline of C program:

}

. . .

```
void e () {
    int e1; -----local variable
```





 Global and local variables' lifetimes are nested. They can never be overlapped.





Outline of C program:

```
void main () {
    float m;
    ... r(3); ...
}
```

```
void r (int n) {
    int i;
    if (n > 1) {
        ... r(n-1); ...
    }
```



Lifetimes of global and local variables:



 Note: A local variable of a recursive procedure/function has several nested lifetimes.



## Example: C heap variables (1)

#### Outline of C program:

struct IntNode {int elem; IntList succ;}
typedef struct IntNode \* IntList;

```
IntList c (int h, IntList t) {
    // Return an IntList with head h and tail t.
    IntList ns =
        (IntList) malloc (sizeof IntNode);
    ns->elem = h; ns->succ = t;
    return ns;
```

```
void d (IntList ns) {
    ns->succ = ns->succ->succ;
```



## Example: C heap variables (2)

#### • Outline *(continued)*:

```
void main {
    IntList l1, l2;
    l1 = c(3, c(5, c(7, NULL)));
    l2 = c(2, l1);
    d(l1);
}
```



## Example: C heap variables (3)

• After initializing 11 and 12:



• After calling d (11):





# Example: C heap variables (4)



 Heap variables' lifetimes can overlap one another and local/global variables' lifetimes.



- An allocator is an operation that creates a heap variable, yielding a pointer to that heap variable.
   E.g.:
  - C's allocator is a library function, malloc().
  - Java's allocator is an expression of the form "new  $C(\ldots)$ ".
- A deallocator is an operation that explicitly destroys a designated heap variable. E.g.:
  - C's deallocator is a library function, free().
  - Java has no deallocator at all.



- A heap variable remains reachable as long as it can be accessed by following pointers from a global or local variable.
- A heap variable's lifetime extends from its creation until:
  - it is destroyed by a deallocator, or
  - it becomes unreachable, or
  - the program terminates.



- A pointer is a reference to a particular variable. (In fact, pointers are sometimes called references.)
- A pointer's referent is the variable to which it refers.
- A null pointer is a special pointer value that has no referent.
- A pointer is essentially the address of its referent in the store.
  - However, each pointer also has a *type*, and the type of a pointer allows us to infer the type of its referent.



- Pointers and heap variables can be used to represent recursive values such as lists and trees.
- But the pointer itself is a low-level concept. Manipulation of pointers is notoriously errorprone and hard to understand.
- For example, the C assignment "p->succ = q;" appears to manipulate a list, but which list? Also:
  - Does it delete nodes from the list?
  - Does it stitch together parts of two different lists?
  - Does it introduce a cycle?



# Dangling pointers (1)

- A dangling pointer is a pointer to a variable that has been destroyed.
- Dangling pointers arise when:
  - a pointer to a *heap variable* still exists after the heap variable is destroyed by a deallocator
  - a pointer to a *local variable* still exists at exit from the block in which the local variable was declared.
- A deallocator immediately destroys a heap variable; all existing pointers to that heap variable then become dangling pointers.
  - Thus deallocators are inherently unsafe.



# Dangling pointers (2)

- C is highly unsafe:
  - After a heap variable is destroyed, pointers to it might still exist.
  - At exit from a block, pointers to its local variables might still exist (e.g., if stored in global variables).
- Java is very safe:
  - It has no deallocator.
  - Pointers to local variables cannot be obtained.



## **Example: C dangling pointers**

#### Consider this C code:

date1->m = 1;

date1 -> d = 1;

free(date2); -

date2 - y = 2009;

DatePtr date2 = date1;

printf("%d4", date1->y);

```
struct Date {int y, m, d;}
typedef Date * DatePtr;
```

```
DatePtr date1 = (DatePtr) makes date1 point to
malloc(sizeof Date); a newly-allocated
date1->y = 2008;
```

```
makes date2 point to
that same heap
variable
```

```
deallocates that heap
variable – date1 and
date2 now contain
dangling pointers
```

```
-- behaves unpredictably
```



- A command (often called a statement) is a program construct that will be executed to update variables.
- Commands are characteristic of imperative and OO PLs (but not functional PLs).
- Simple commands:
  - A skip command is a command that does nothing.
  - An assignment command is a command that uses a value to update a variable.
  - A procedure call is a command that calls a proper procedure with argument(s). Its net effect is to update some variables.



- Compound commands:
  - A sequential command is a command that executes its sub-commands in sequence.
  - A conditional command is a command that chooses one of its sub-commands to execute.
  - An iterative command is a command that executes its sub-command repeatedly. This may be:
    - **definite iteration** (where the number of repetitions is known in advance)
    - **indefinite iteration** (where the number of repetitions is not known in advance).
  - A block command is a command that contains declarations of local variables, etc.



Java single assignment:

```
m = n + 1;
```

Java multiple assignment:

m = n = 0;

Java assignment combined with binary operator:

m += 7; ------ equivalent to "m = m+7;"

n /= b; ----- equivalent to "n = n/b;"



Java if-command:

if (x > y)
 out.print(x);
else
 out.print(y);

Java switch-command:

```
Date today = ...;
switch (today.m) {
    case 1: out.print("JAN"); break;
    case 2: out.print("FEB"); break;
    ...
    case 11: out.print("NOV"); break;
    case 12: out.print("DEC"); 
    Breaks are
    essential here!
```



Java while-command:

```
Date[] dates;
...
int i = 0;
while (i < dates.length) {
    out.print(dates[i]);
    i++;
}
```

Java for-commands (both forms):

```
for (int i = 0; i < dates.length; i++)
    out.print(dates[i]);</pre>
```

```
for (Date d : dates)
    out.print(d);
```



Java do-while-command:

```
static String render (int n) {
   String s = "";
   int m = n;
   do {
      char d = (char) (m % 10) + '0';
      s = d + s;
      m /= 10;
   } while (m > 0);
   return s;
}
```

 Here the loop condition is evaluated after each repetition of the loop body.



#### **Example: Java block command**

Java block command:

```
if (x > y) {
    int z = x;
    x = y;
    y = z;
}
```



- The primary purpose of evaluating an expression is to yield a value.
- In most imperative and OO PLs, evaluating an expression can also update variables – these are side effects.
- In C and Java, the body of a function is a command. If that command updates global or heap variables, calling the function has side effects.
- In C and Java, assignments are in fact expressions with side effects: "V = E" stores the value of E in V as well as yielding that value.



- The C function getchar(fp) reads a character and updates the file variable that fp points to.
- The following C code is correct:

```
char ch;
while ((ch = getchar(fp)) != NULL)
    putchar(ch);
```

The following C code is incorrect (why?):

```
enum Gender {FEMALE, MALE};
Gender g;
if (getchar(fp) == 'F') g = FEMALE;
else if (getchar(fp) == 'M') g = MALE;
else ...
```