



University
of Glasgow

The Judgment of Forseti

Economic Utility for Dynamic Heap Sizing of Multiple Runtimes

Jeremy.Singer@glasgow.ac.uk

digested talk



Jeremy Singer @jsinger_compsci · 6s

#forseti optimizes whole-system thruput by tweaking heap sizes of all executing JVMs simultaneously dx.doi.org/10.1145/275416... #isimm2015



David.Vengerov @ oracle.com



Callum.Cameron @ glasgow.ac.uk

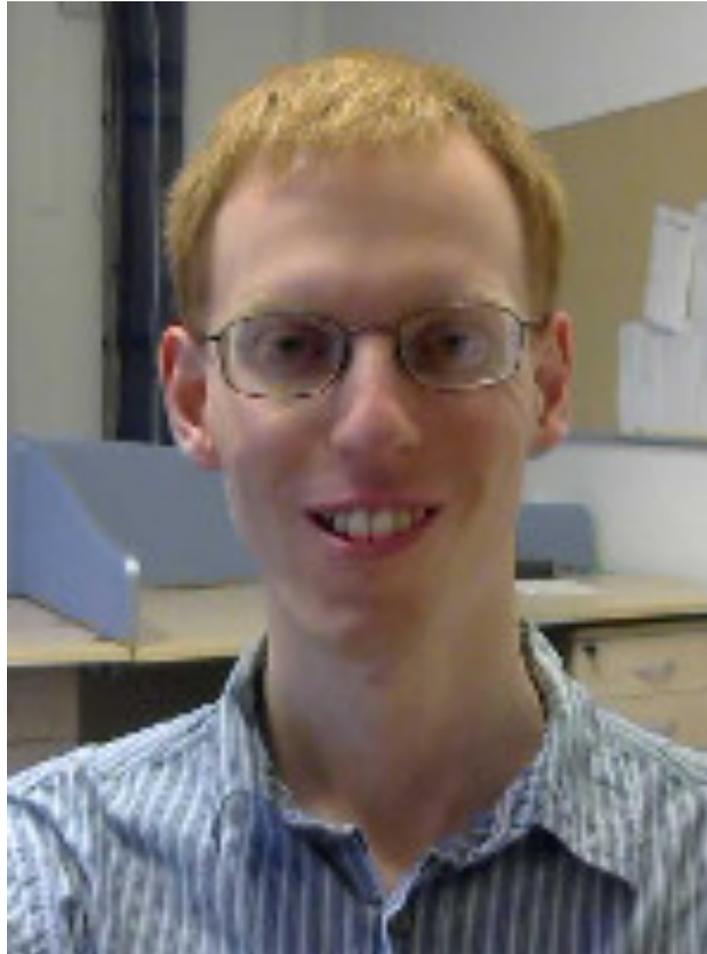






image: wikipedia.org

Motivation

dynamic memory
resource
allocation in
datacenters



image: wikipedia.org

dynamic memory
resource
allocation in
smartphones



image: wikipedia.org

Requirements

- Satisfy users
- Be economical

Characteristics of VM tasks

- elastic memory usage
- phased behavior

Automatic Mem Mgt

- a.k.a. *Garbage Collection (GC)*
- Automatically deallocate a block of memory when it is no longer reachable
- *Runtime Heap* grows/shrinks on a demand basis

Key Heap Metrics

- *Live size* – current amount of *live* data
- *Current heap occupancy* – current amount of *allocated* data (live and dead)
- *Heap limit* – max permitted value for current heap occupancy

What is the
optimal
heap limit?

Lots of possibilities

How do you find the best settings for your system? ... for your application?

1. domain expertise
2. exhaustive search
3. mathematical model

State-of-the-art: Domain Expertise

```
Java -Xmx12g -XX:MaxPermSize=64M -XX:PermSize=32M-XX:MaxNewSize=2g  
-XX:NewSize=1g -XX:SurvivorRatio=128 -XX:+UseParNewGC  
-XX:+UseConcMarkSweepGC -XX:MaxTenuringThreshold=0  
-XX:CMSInitiatingOccupancyFraction=60 -XX:+CMSParallelRemarkEnabled  
-XX:+UseCMSInitatingOccupancyOnly -XX:ParallelGCThreads=12  
-XX:LargePageSizeInBytes=256m ...
```



State-of-the-art: Exhaustive Search

The Taming of the Shrew: Increasing Performance by Automatic Parameter Tuning for Java Garbage Collectors

Philipp Lengauer
Christian Doppler Laboratory MEVSS
Johannes Kepler University Linz, Austria
philipp.lengauer@jku.at

Hanspeter Mössenböck
Institute for System Software
Johannes Kepler University Linz, Austria
hanspeter.moessenboeck@jku.at

ABSTRACT

Garbage collection, if not tuned properly, can considerably impact application performance. Unfortunately, configur-

However, while object allocations produce a direct and easy to understand performance impact, the costs of garbage collections are easily overlooked. Programmers are often unaware of the proportion their application spends on collect-

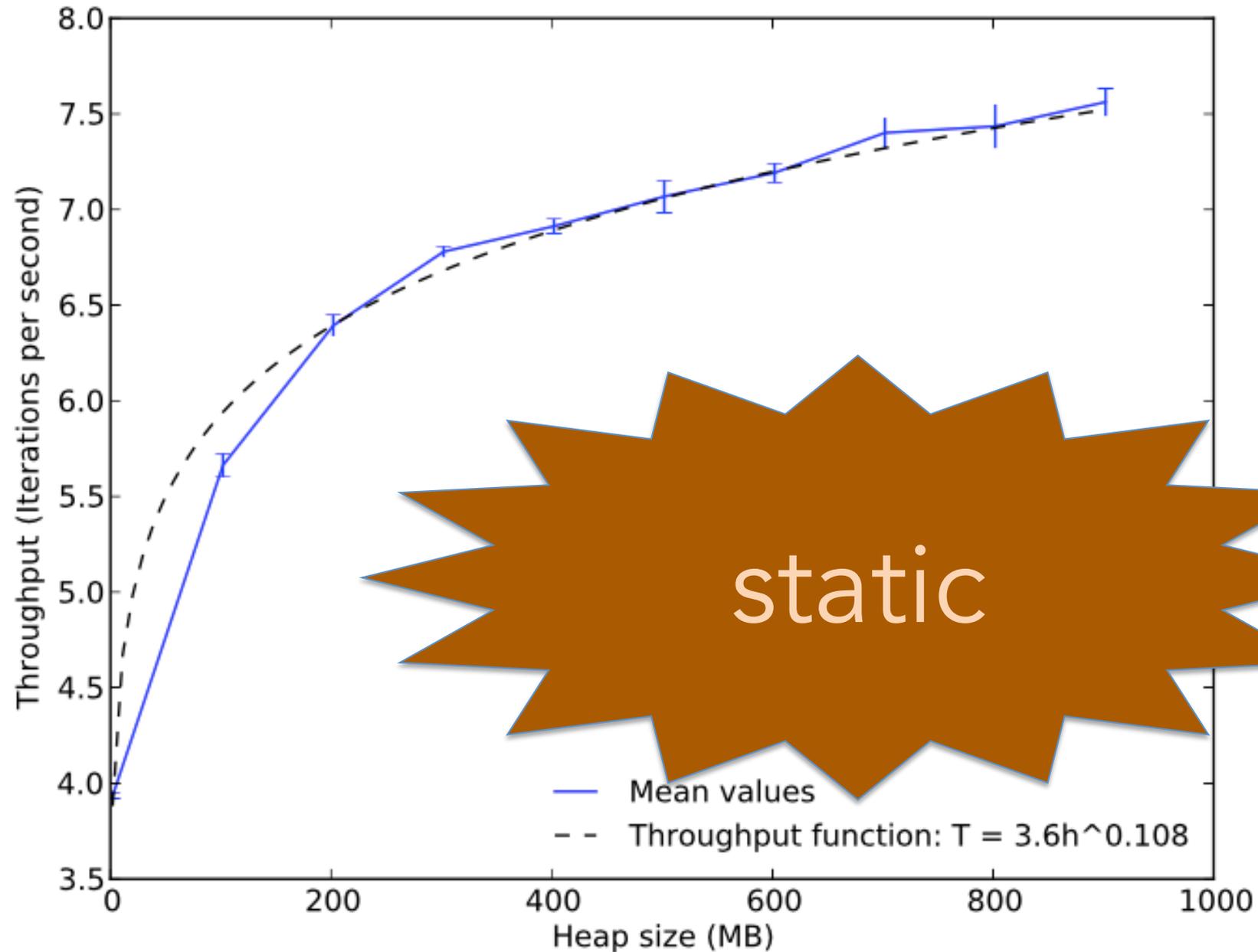
- around **300** GC parameters
- search parameter space for **4 hours**
- select **best** configuration

[ICPE 2014]

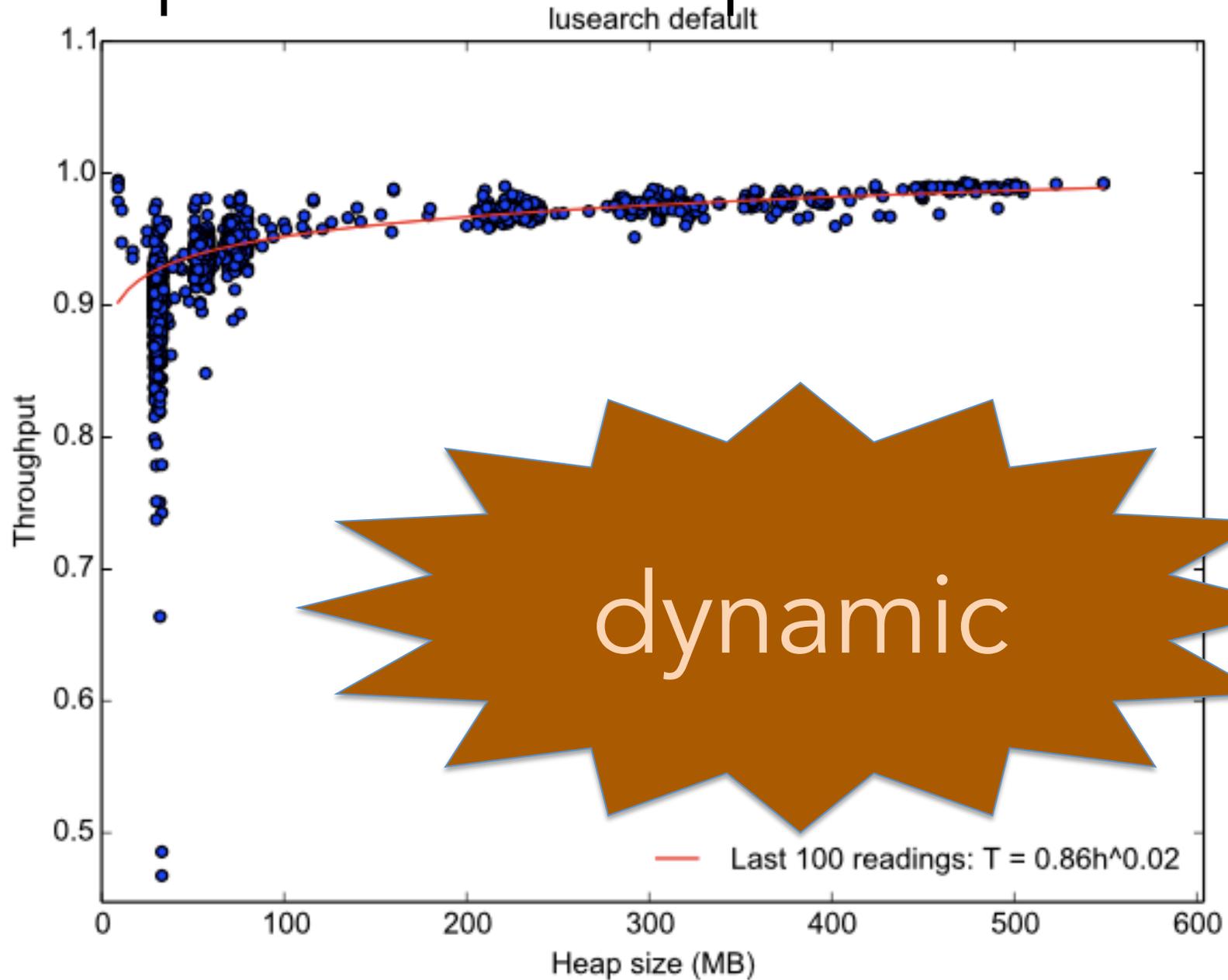
State-of-the-art: Mathematical Model

- decision tree
 - machine learning [ISMM 2007]
- supply/demand curve
 - economics [ISMM 2010]
- differential equations
 - control theory [ISMM 2013]

heap limit affects performance



heap limit affects performance



Math Model based on *economic utility*

For a single VM:
the utility function has form

$$U(h) = ah^b$$

with $a > 0$, $0 < b < 1$

Math Model based on *economic utility*

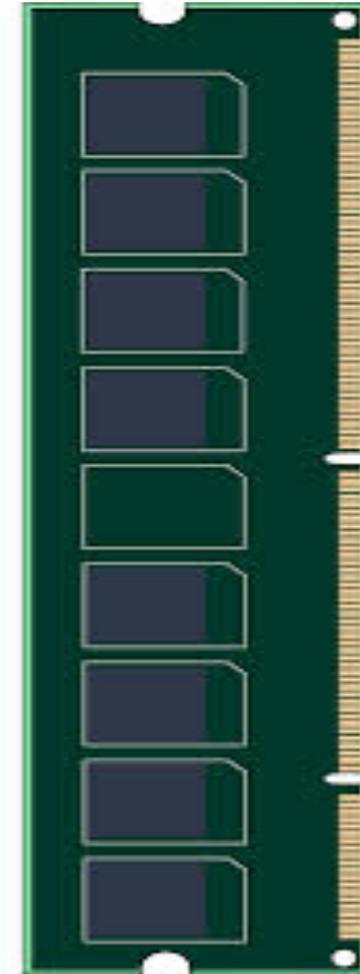
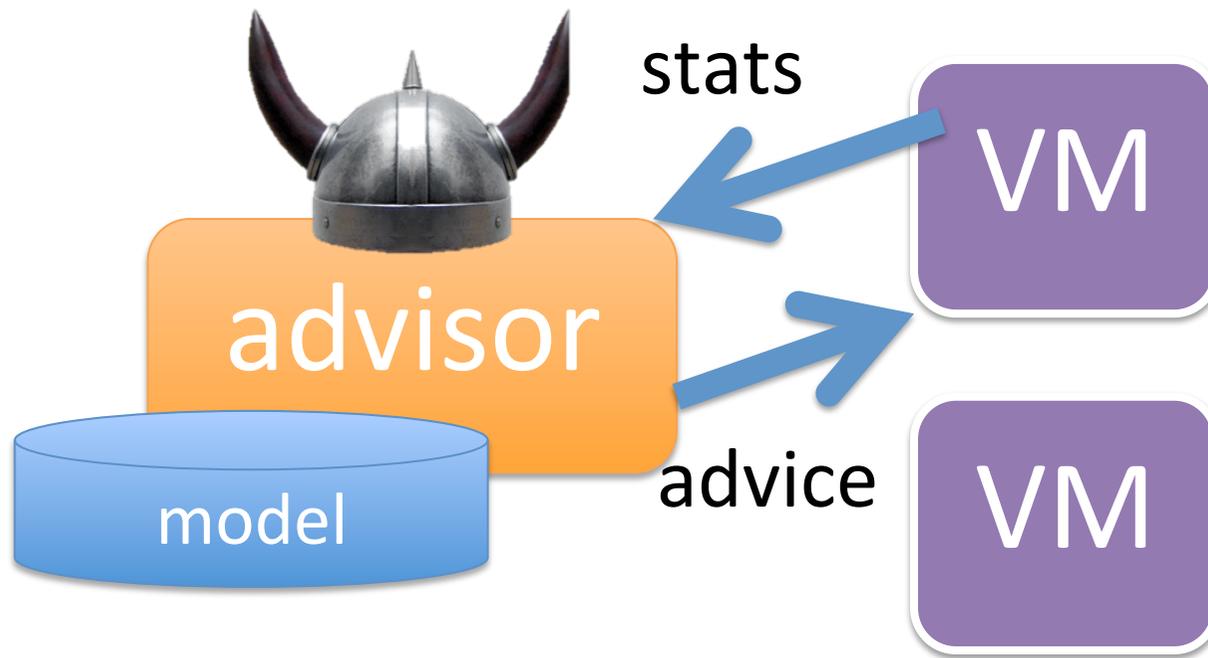
overall utility function for whole-system:
product of individual utilities

$$U(h_1, \dots, h_n) = \prod_i a_i h_i^{b_i}$$

Math Model based on *economic utility*

- *maximise* overall utility function
 - possibly via analytic solution [ICOOOLPS'14]
 - here we use *numeric* optimization

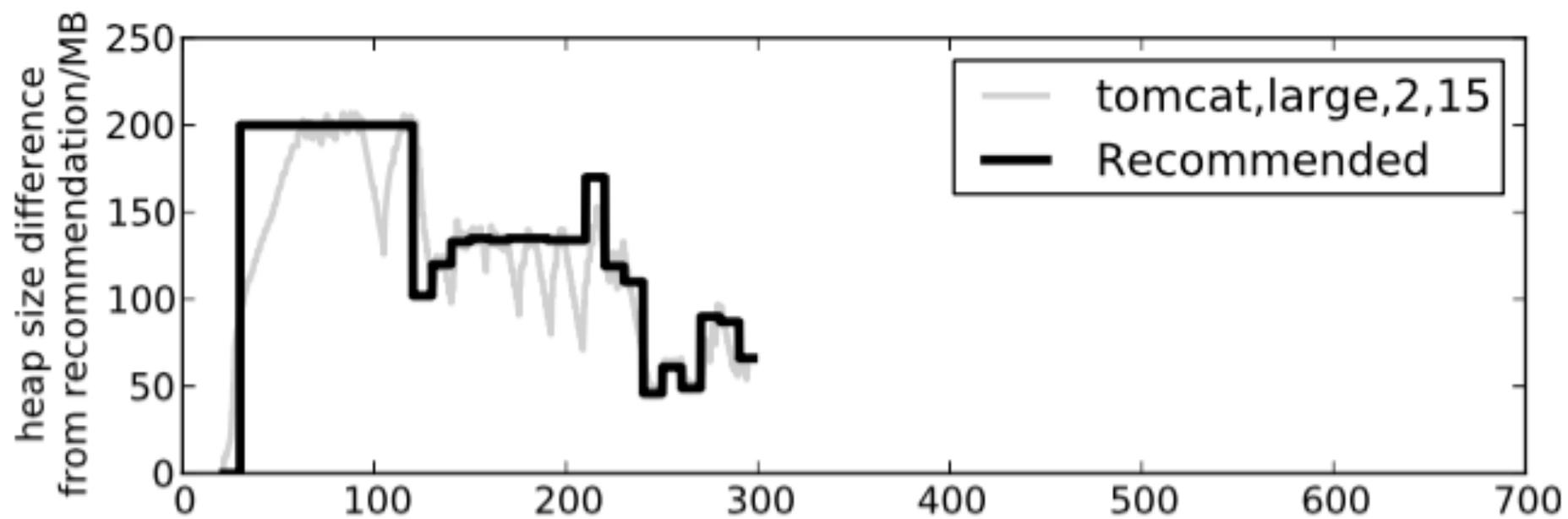
Forseti concept

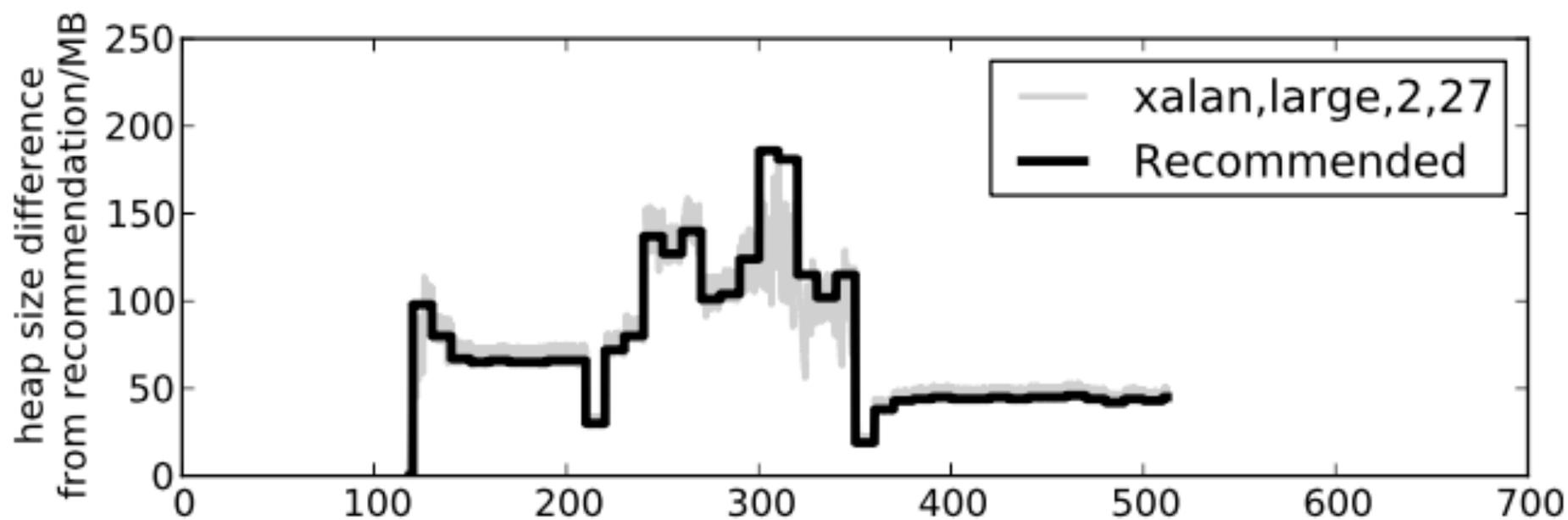


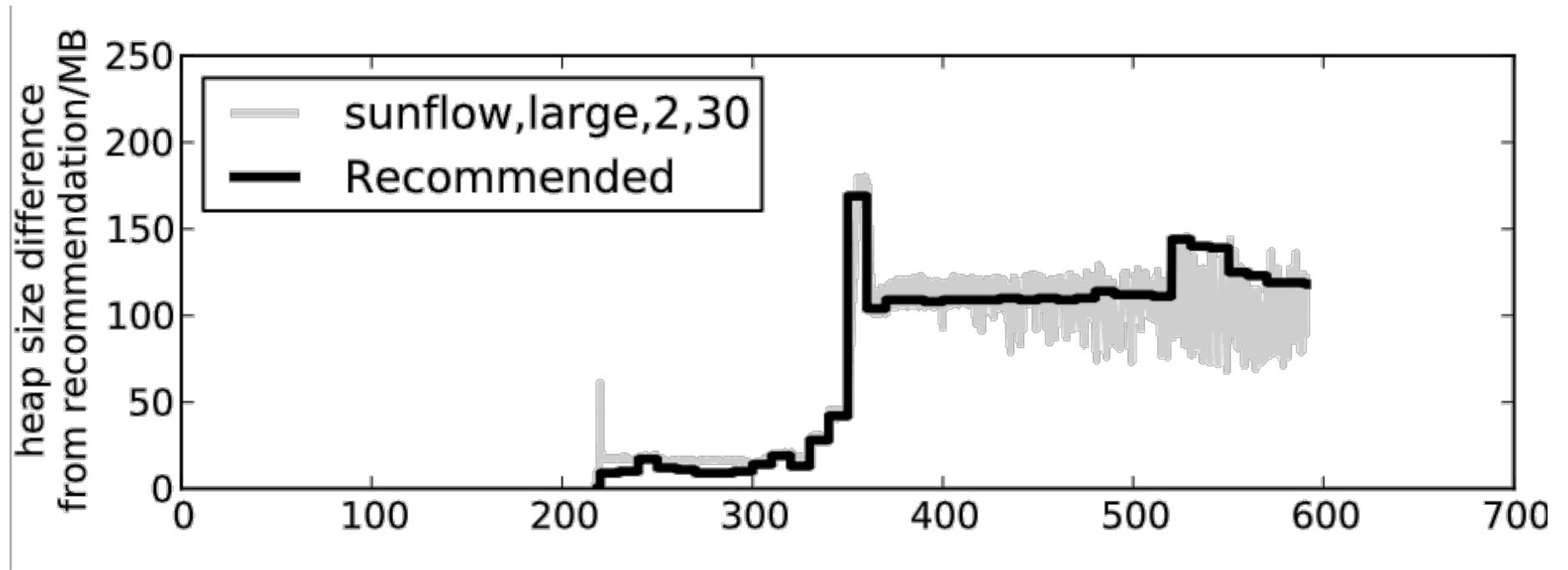
Evaluation

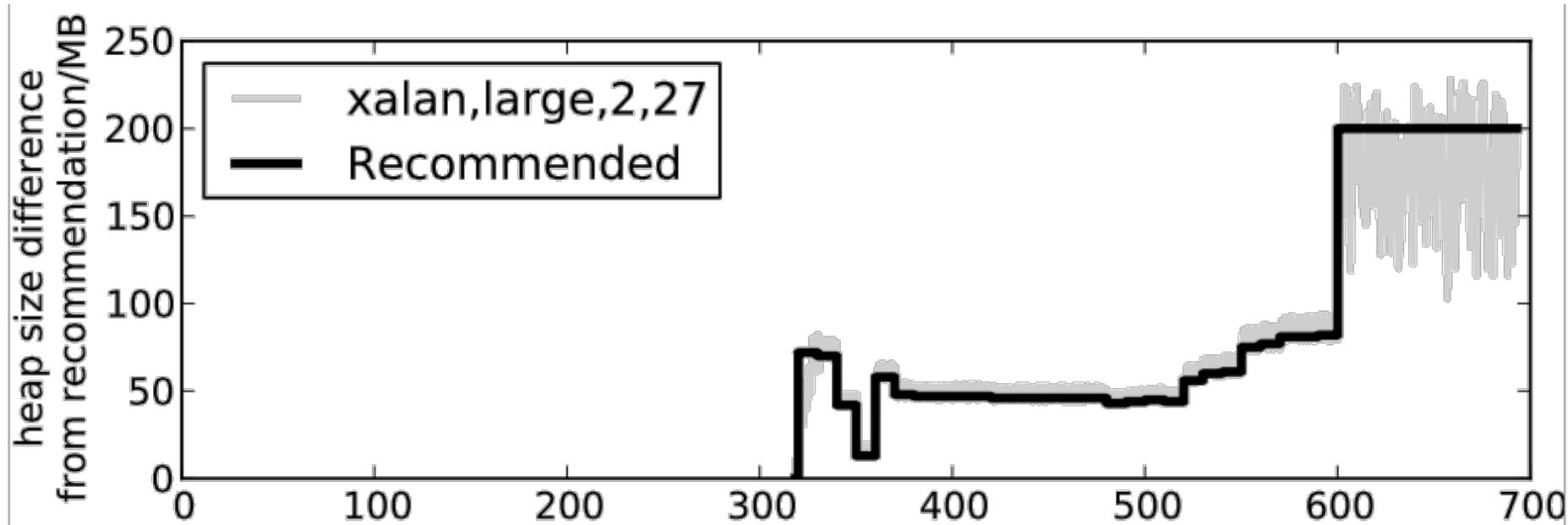
Experiment 1:

- run 4 DaCapo benchmarks
- staggered start times
- set target total mem usage to *200MB*

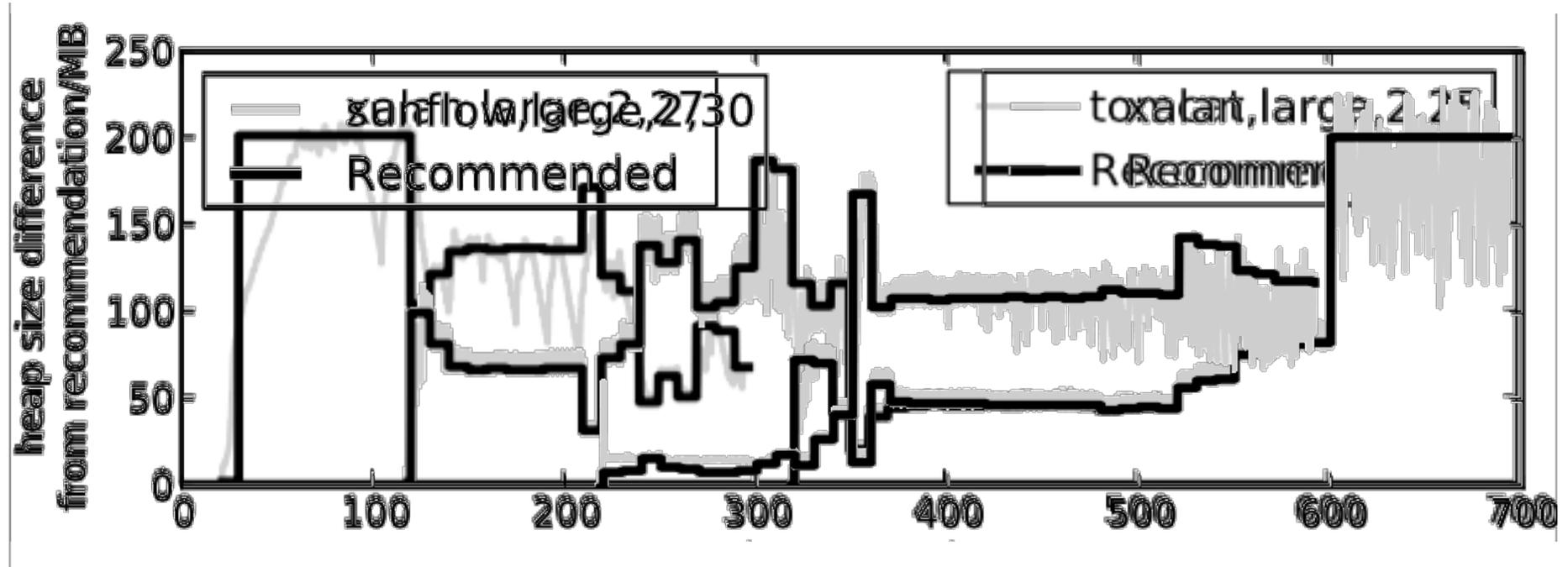




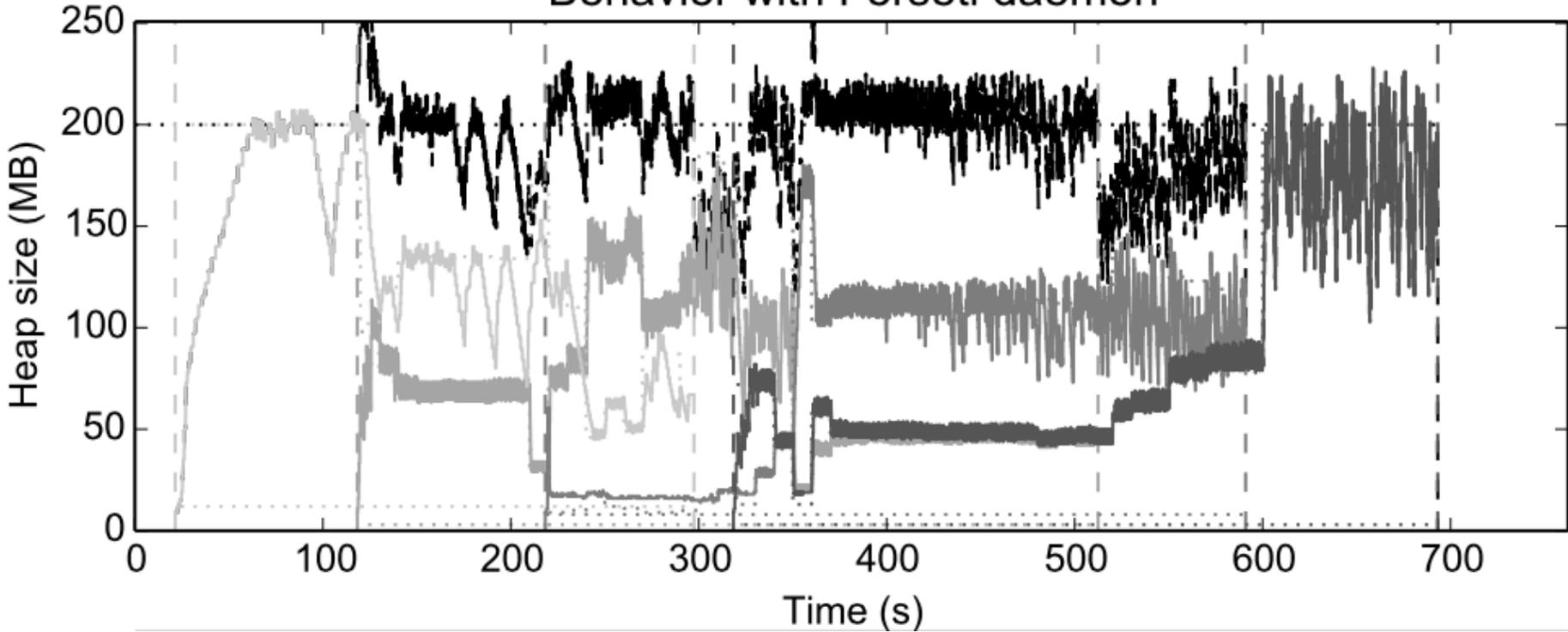




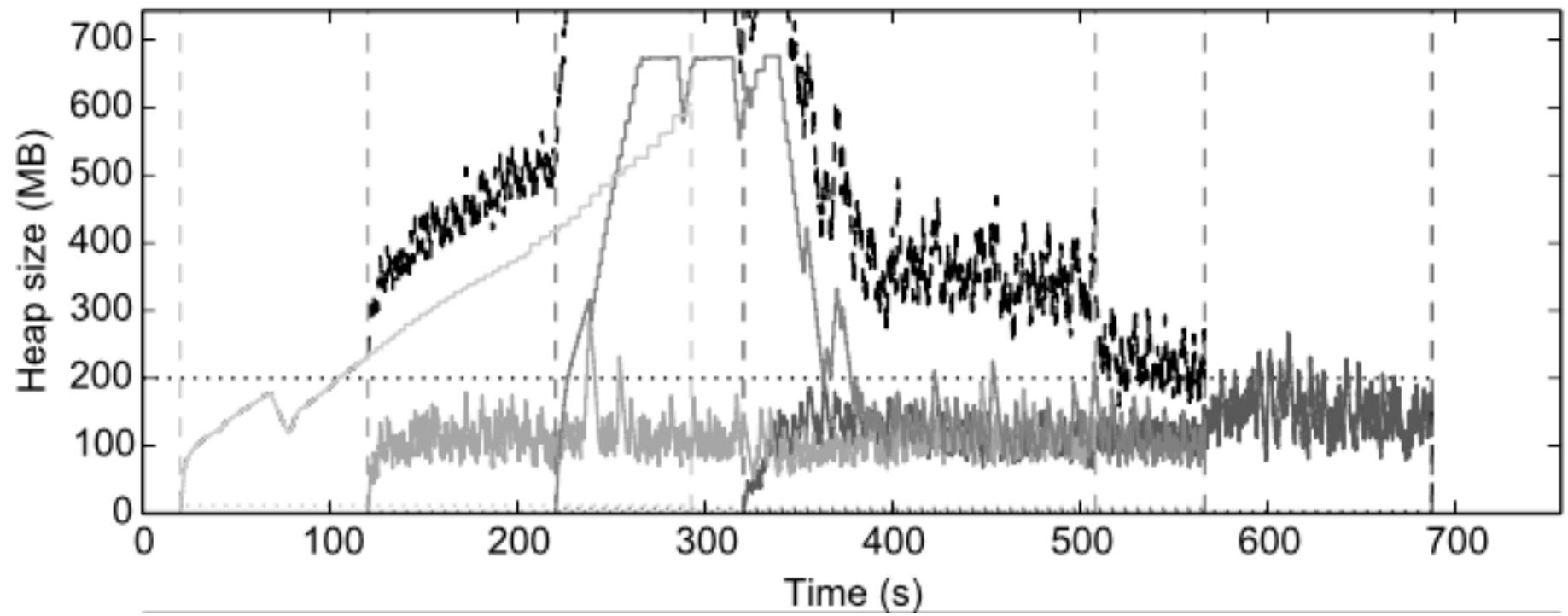
OHP slide reminiscence...



Behavior with Forseti daemon



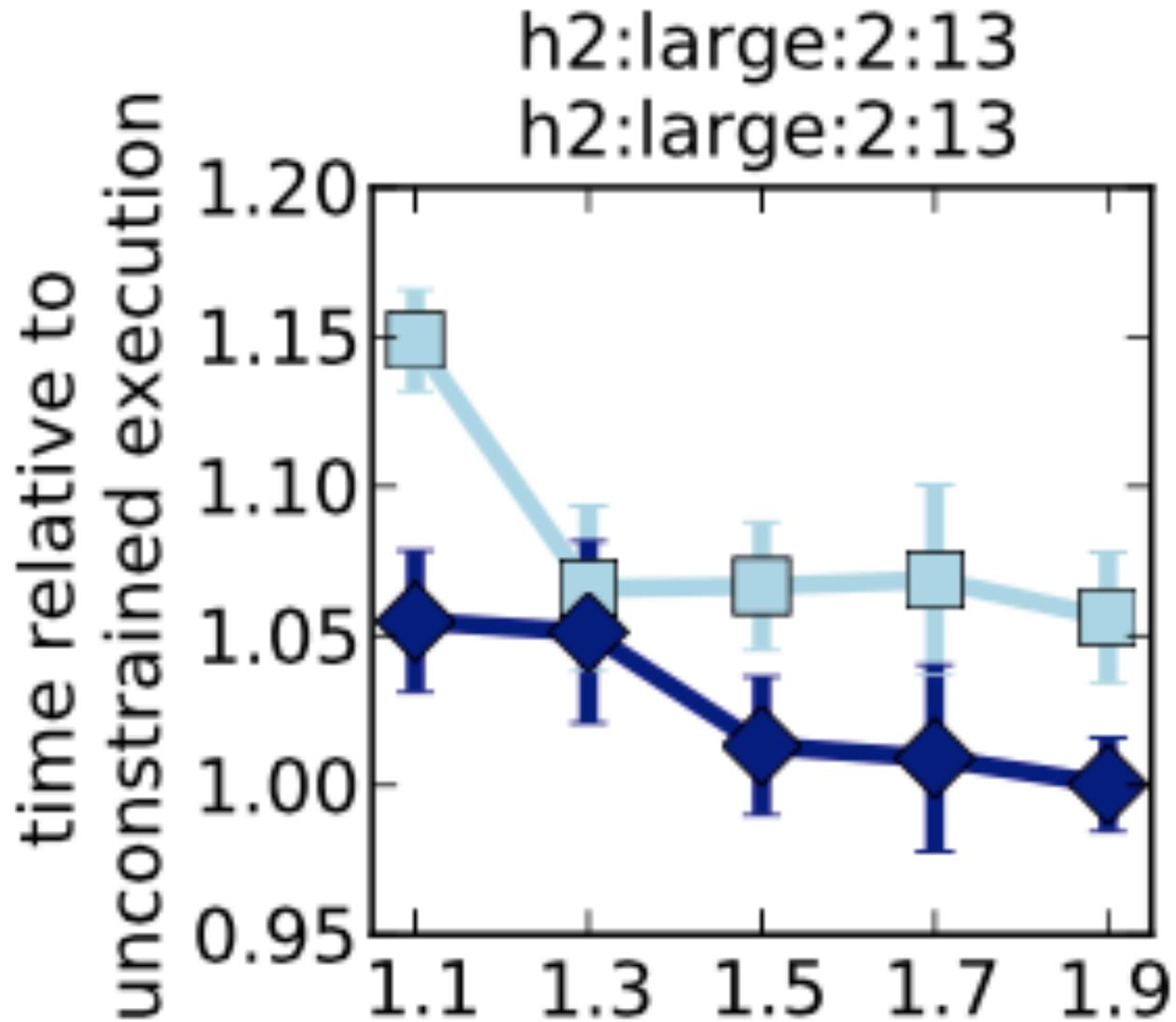
Behavior without Forseti daemon



Experiment 2:

- run pairs of Java benchmarks concurrently
- Set target total mem usage to $1.1..1.9 \times \Sigma \text{ minheaps}$
- compare execution time with
 - Forseti
 - Static fixed heaps
 - Less constrained sizing

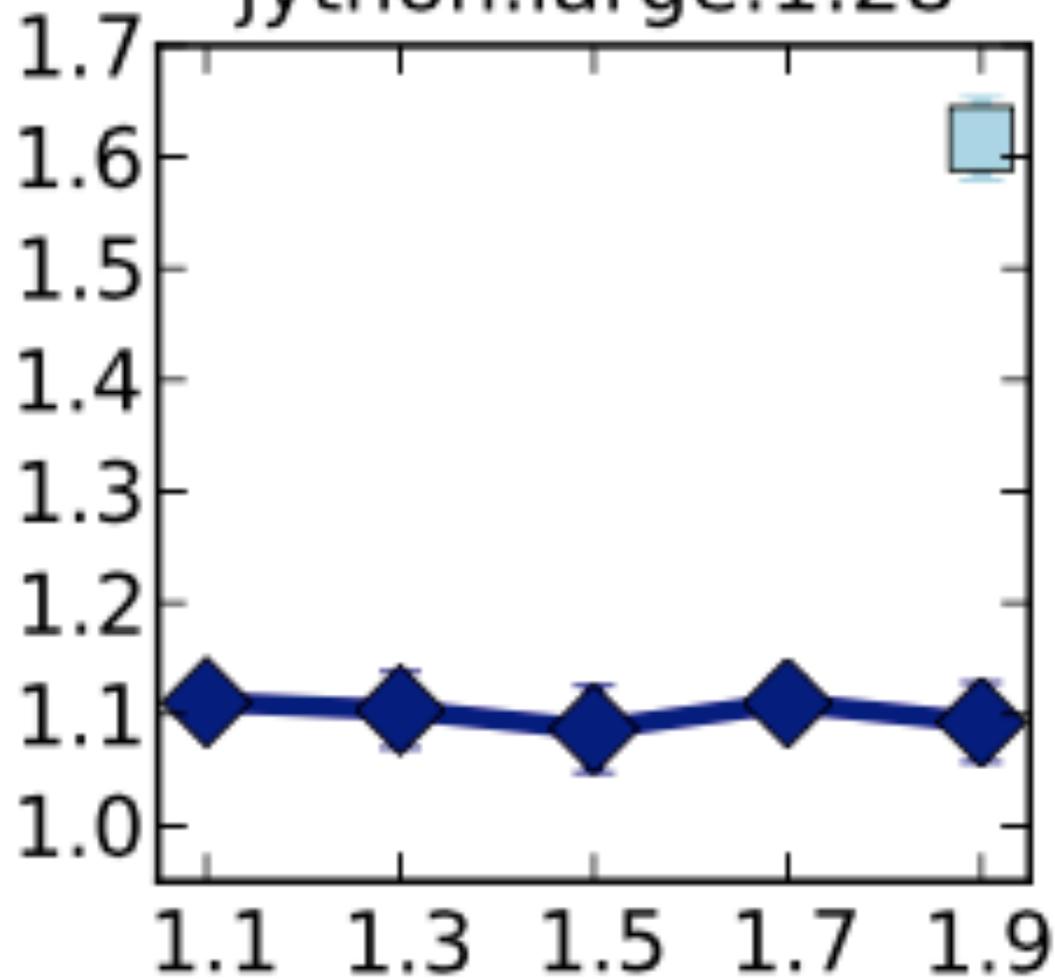
Relative Performance



Relative Performance

h2:large:2:13

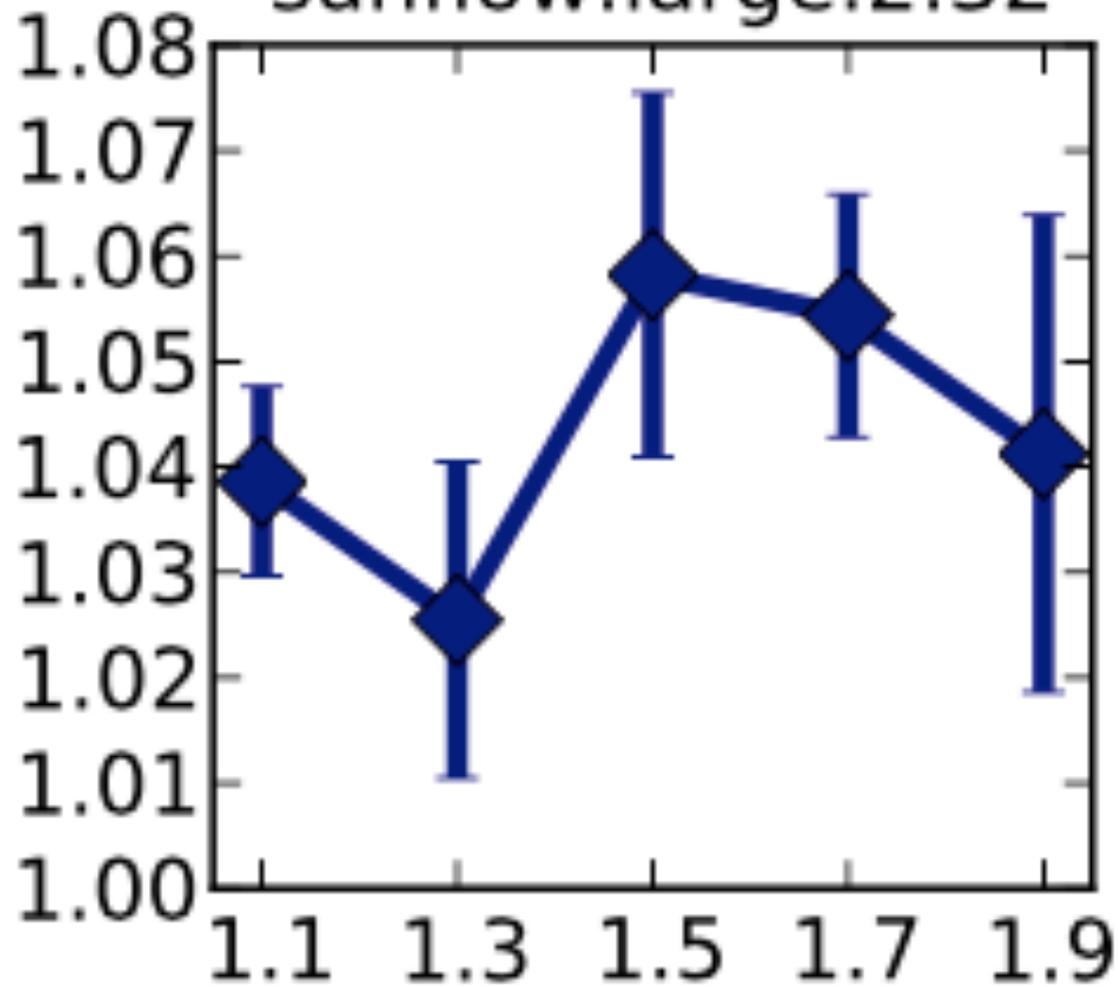
jython:large:1:28



Relative Performance

h2:large:2:13

sunflow:large:2:32



Overheads

In all reported experiments, the time overhead for running the Forseti daemon is small. We analyzed the 6104 experimental runs completed for this paper:

- mean experiment wall clock time is 412 seconds (max is 2300 seconds).
- mean daemon CPU time is 1.00 seconds (max is 5.94 seconds).
- mean daemon memory footprint is 23MB (max is 29MB).

Conclusions

Garbage Collectors require *Holistic Systems Optimization*

- Model must consider all VMs in system
- Optimize holistically, not in isolation
- Is this a new OS service?
- Generalizability?

end of presentation

Trust Issues

- Can Forseti trust VM readings?
 - spoofing or denial-of-service attacks
 - use other metrics (performance counters)
- Can VMs ignore Forseti advice?
 - yes, at present
 - but Forseti could interact with the OS mem mgr to be more aggressive

Alternative Utility Functions

Benthamite (utilitarian – maximize total utility)

$$W = \sum_{i=1}^n Y_i$$

Rawlsian (consider least well-off individual)

$$W = \min(Y_1, Y_2, \dots, Y_n)$$

Related Work

- Alonso and Appel [SIGMETRICS 1990]
 - their advisor daemon works to prevent paging in SML/NJ runtimes due to excessive heap growth
 - advice not based on economic utility model explicitly

get our code

[https://bitbucket.org/jsinger/
economics_memory_code](https://bitbucket.org/jsinger/economics_memory_code)

Throughput for staggered multi-VM experiment

