

# Deleting Edges to Save Cows: Using Graph Theory to Control the Spread of Disease in Livestock

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Joint work with Jessica Enright (University of Stirling)

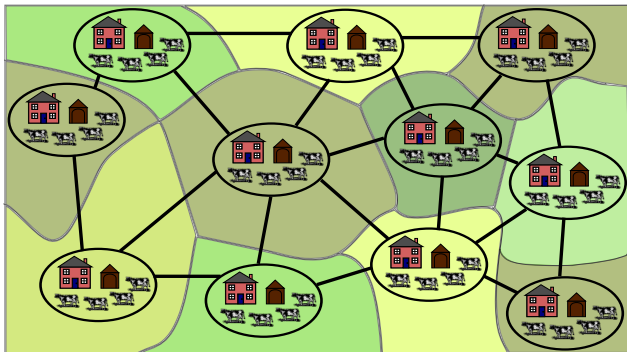


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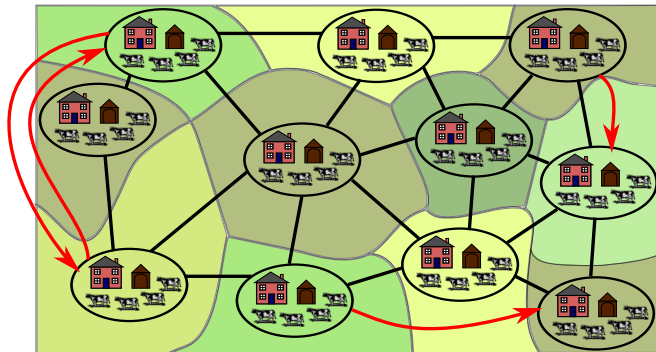
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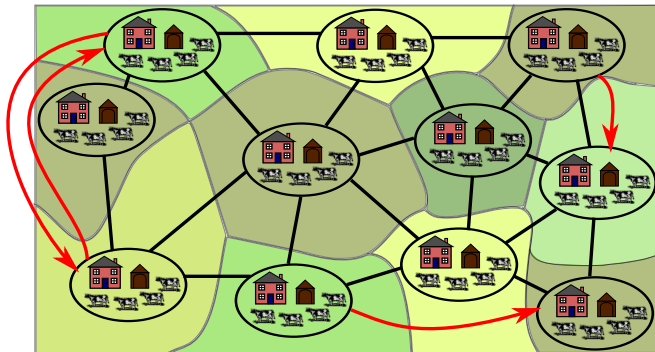


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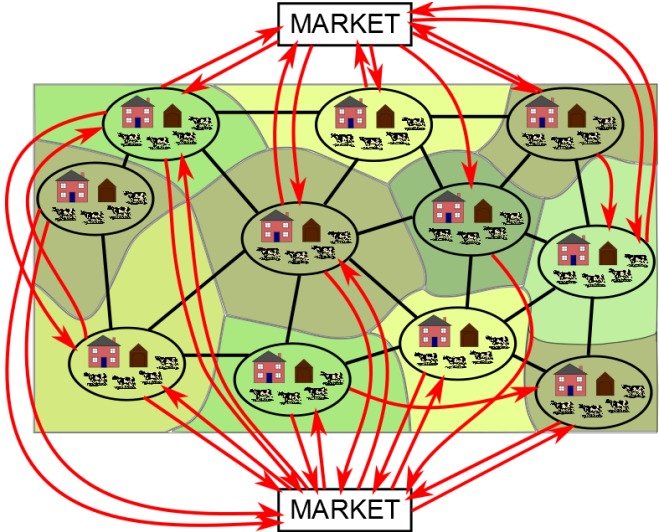
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**Cost of modifications** The cost of deleting individual vertices/edges may vary; this can be captured with a weight function on vertices and/or edges.

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- ▶ consider the total number of animals in e.g. a connected component, rather than just the number of animal holdings
- ▶ place more or less strict restrictions on individual animal holdings

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**PROBLEM:** There is no polynomial-time algorithm to solve this problem in general unless  $P=NP$  (even if  $h = 3$ ).



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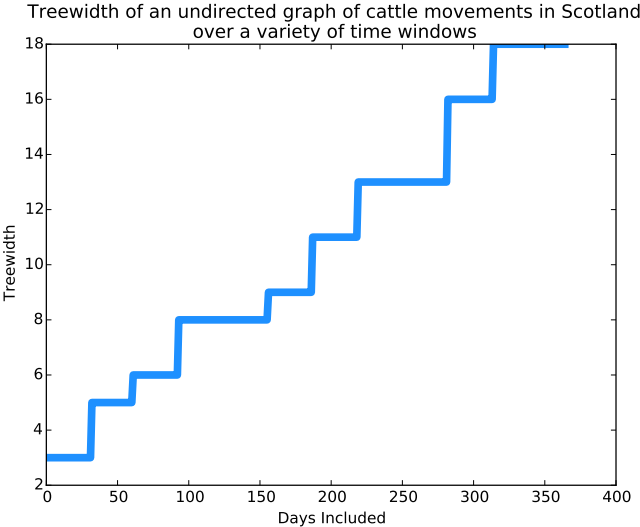
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- ▶ Animal trade networks are very unlikely to form trees, but they might have some similarities to trees.
- ▶ The **treewidth** of a graph is a measure of how “tree-like” a graph is, in a specific sense. Trees have treewidth equal to 1, and cycles have treewidth 2.
- ▶ Often algorithmic problems can be solved more efficiently on graphs with small treewidth.

# (Some) cattle trade networks have small treewidth

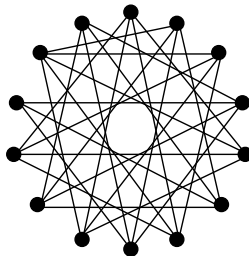
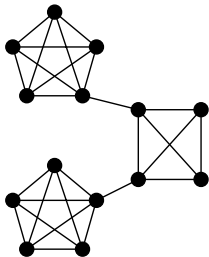


# New results

## Theorem (Enright and M., 2015)

*Suppose we are given a (weighted) graph  $G$  on  $n$  vertices which has treewidth  $w$ . We can determine the least costly set of edges to delete, so that the remaining graph has no connected component with more than  $h$  vertices, in time  $O((wh)^{2w} n)$ .*

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*Suppose we are given a (weighted) graph  $G$  on  $n$  vertices which has treewidth  $w$ . We can determine the least costly set of edges to delete, so that the remaining graph contains no graph from the set  $\mathcal{F}$  as a subgraph, in time  $2^{O(|\mathcal{F}|w^r)}(n + 2^r)$ , if no element of  $\mathcal{F}$  has more than  $r$  vertices.*





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**Thank you**