

Priority Queue

What is that?

Implementation with linked list with $O(n)$ behaviour

The Heap ($O(\log(n))$)

An implementation using an array

Your mission ...

- Store a collection of prioritized elements
 - Elements are comparable
- Allow insertion of an element
- Can only remove the element with highest priority
 - Comes first in order

```
Q.insert(e)
e = Q.removeMin()
Q.size()
Q.isEmpty()
Q.min()
```

We present 3 implementations

Example applications of a priority queue

- Dispatching processes in a computer
- Hospital waiting lists
- Standby passengers for a flight
- Queuing at call centres
- Internal data structure for another algorithm (graph algorithm)
- ...

- Store a collection of prioritized elements
 - ***Elements are comparable***
- Allow insertion of an element
- Can only remove the element with highest priority
 - Comes first in order

Two examples on ***comparing things***

This is a Vertex

```
import java.util.*;

public class Vertex implements Comparable<Vertex> {

    int index, degree, colour, saturation, nebDeg;
    boolean[] domain;

    public Vertex (int index,int degree) {
        this.index = index;
        this.degree = degree;
        nebDeg = 0;
    }
}
```

This is a Comparator for vertices

```
File Edit Format View Help
import java.util.*;

public class MCRComparator implements Comparator {

    public int compare(Object o1, Object o2){
        Vertex u = (Vertex) o1;
        Vertex v = (Vertex) o2;
        if (u.degree < v.degree ||
            u.degree == v.degree && u.nebDeg < v.nebDeg ||
            u.degree == v.degree && u.nebDeg == v.nebDeg && u.index > v.index) return 1;
        return -1;
    }
    //
    // to sort vertices by decreasing degree, tie breaking on neighbourhood degree (nebDeg)
    //
}
```

Using the comparator to sort an array of Vertex

```
boolean conflicts(int v,ArrayList<Integer> colourClass){
    for (int i=0;i<colourClass.size();i++){
        int w = colourClass.get(i);
        if (A[v][w] == 1) return true;
    }
    return false;
}

void orderVertices(ArrayList<Integer> Colord){
    Vertex[] V = new Vertex[n];
    for (int i=0;i<n;i++) V[i] = new Vertex(i,degree[i]);
    for (int i=0;i<n;i++)
        for (int j=0;j<n;j++)
            if (A[i][j] == 1) V[i].nebDeg = V[i].nebDeg + degree[j];
    if (style == 1) Arrays.sort(V);
    if (style == 2) minwidthorder(V);
    if (style == 3) Arrays.sort(V,new MCRComparator());
    for (Vertex v : V) Colord.add(v.index);
}

void minwidthorder(Vertex[] V){
    ArrayList<Vertex> L = new ArrayList<Vertex>(n);
    Stack<Vertex> S = new Stack<Vertex>();
    for (Vertex v : V) L.add(v);
    while (!L.isEmpty()){
        Vertex v = L.get(0);
        for (Vertex u : L) if (u.degree < v.degree) v = u;
        S.push(v); L.remove(v);
        for (Vertex u : L) if (A[u.index][v.index] == 1) u.degree--;
    }
    int k = 0;
    while (!S.isEmpty()) V[k++] = S.pop();
}
```

Another example: a Car

```
public class Car {  
    String make, model;  
    public Car(String s1,String s2){  
        make = s1; model = s2;  
    }  
    public String make(){return make;}  
    public String model(){return model;}  
    public String toString(){return make + " " + model;}  
}
```

This is a CarComparator

```
import java.util.*;

public class CarComparator implements Comparator<Car> {

    public int compare(Car a, Car b){
        int c1 = a.make().compareTo(b.make());
        int c2 = a.model().compareTo(b.model());
        if (c1 == 0) return c2;
        return c1;
    }
    ///
    /// make is most significant
    ///
}
```

Using the CarComparator

```
import java.util.*;
public class Test2 {
    public static void main(String args[]){
        TreeSet<Car> S = new TreeSet<Car>(new CarComparator());

        Car c1 = new Car("Citroen", "C1");
        Car c2 = new Car("Ford", "Mustang");
        Car c3 = new Car("Ferarri", "GTO");
        Car c4 = new Car("Cadillac", "Elderado");
        Car c5 = new Car("Ford", "Mustang");

        S.add(c1);
        S.add(c2);
        S.add(c3);
        S.add(c4);
        S.add(c5);

        System.out.println(S);
    }
}
```

```
Z:\public_html\ads2\java\compare>javac Test2.java
```

```
Z:\public_html\ads2\java\compare>java Test2
```

```
[Cadillac Elderado, Citroen C1, Ferarri GTO, Ford Mustang]
```

```
Z:\public_html\ads2\java\compare>
```

We might use a linked list

- To insert we add to the front of the list
- To find the minimum we must iterate over entire the list
- To remove the minimum we must find the minimum and remove it
- Maintain a counter of number of elements in the list

Method	Time
size	$O(1)$
isEmpty	$O(1)$
insert	$O(1)$
removeMin	$O(n)$
min	$O(n)$

We might use a linked list

- The list is **maintained in non-decreasing order**
- To insert we scan to find position and splice in (see below)
- To find the minimum we deliver the first element in the list
- To remove the minimum we return and remove the first element

```
public void insert(E s){
    if (head == null || head.getElement().compareTo(s) > 0)
        head = new Node<E>(s,head);
    else {
        Node<E> cursor = head;
        Node<E> next = cursor.getNext();
        while (next != null && next.getElement().compareTo(s) <= 0 ){
            cursor = next;
            next = next.getNext();
        }
        cursor.setNext(new Node<E>(s,next));
    }
    size++;
}
```

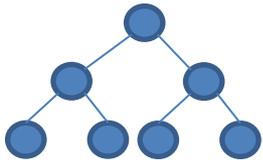
Method	Time
size	O(1)
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An alternative
THE HEAP

- a heap H is a binary tree

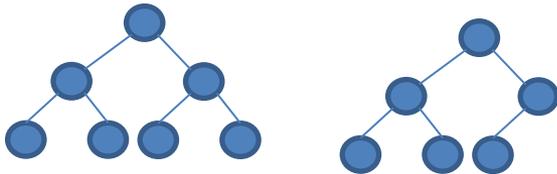
- a heap H is a binary tree
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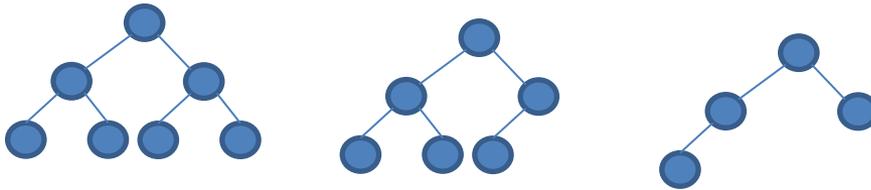
- Fill up level d before moving to level $d+1$
- each level but the last must be full
 - in last level fill from left to right

- a heap H is a binary tree
- H is a **complete** binary tree



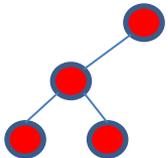
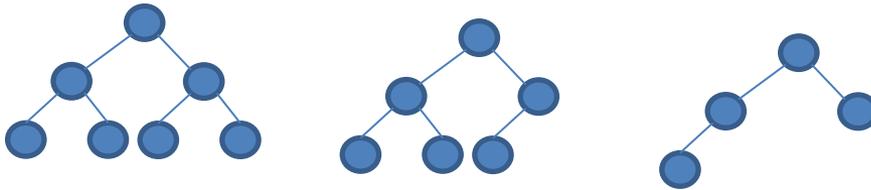
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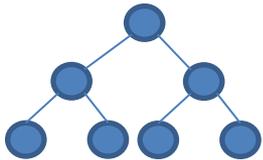
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Not a heap!

- Fill up level d before moving to level $d+1$
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- a heap H is a binary tree
- H is a ***complete*** binary tree
- ***heap order property*** is maintained

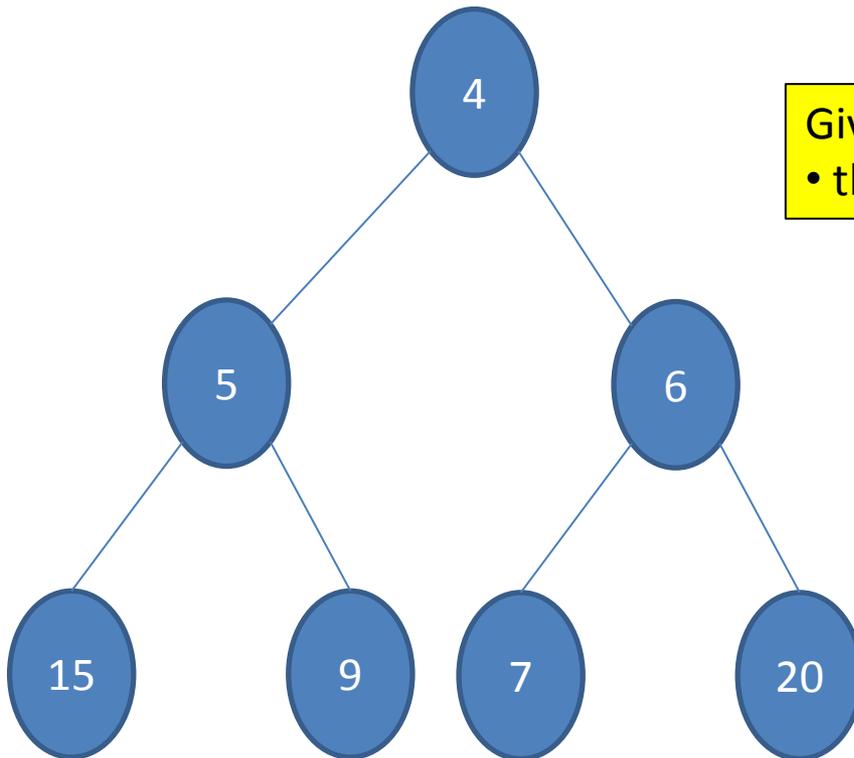


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Given a node v (not the root)

- the parent of v is less than or equal to v

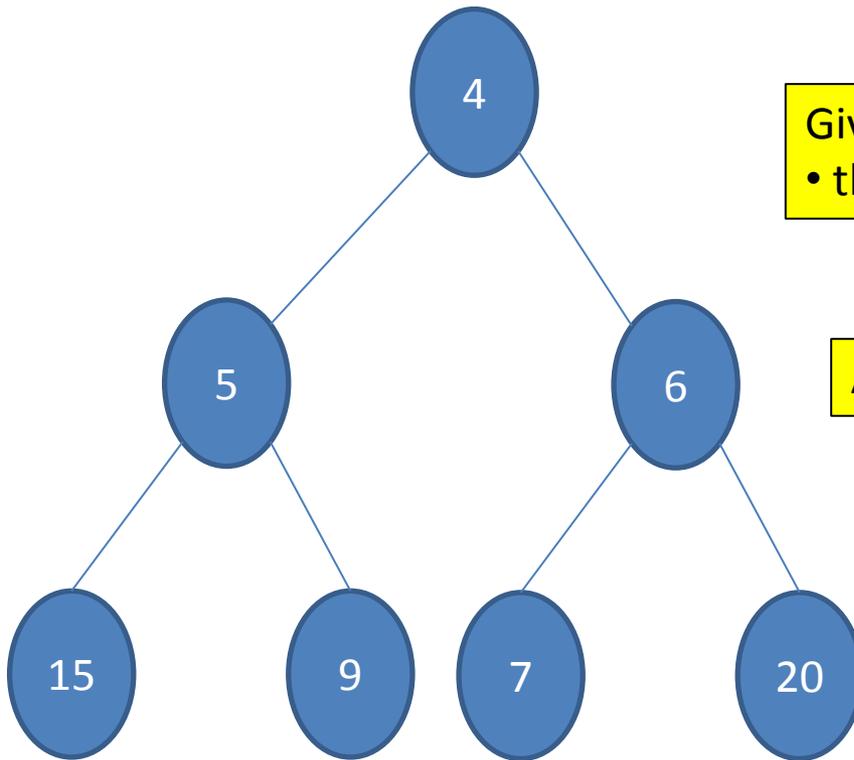
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Given a node v (not the root)

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Given a node v (not the root)

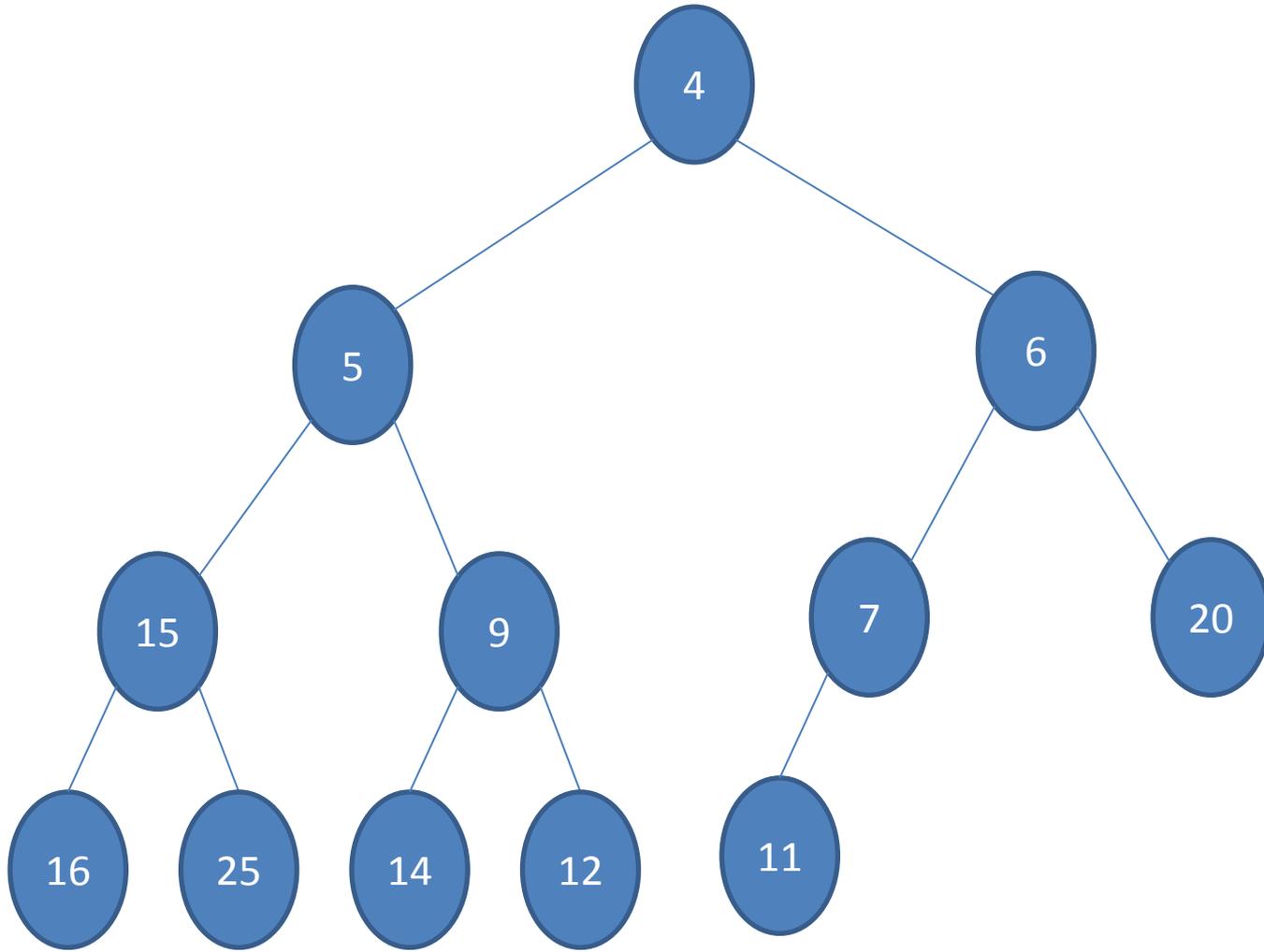
- the parent of v is less than or equal to v

A heap H with n nodes has height $O(\log(n))$

Example: adding to a heap

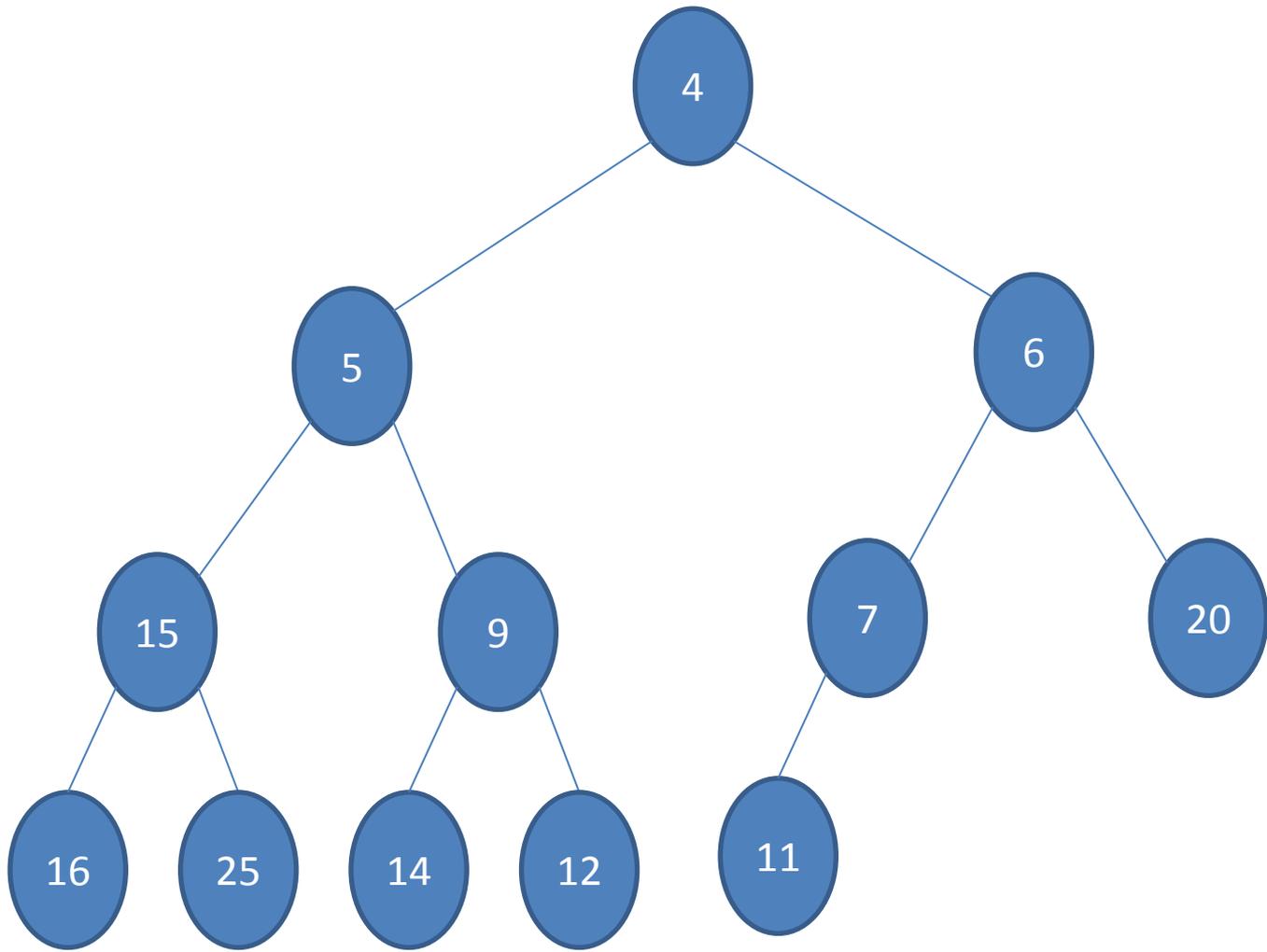
Example: adding to a heap

heap



Example: adding to a heap

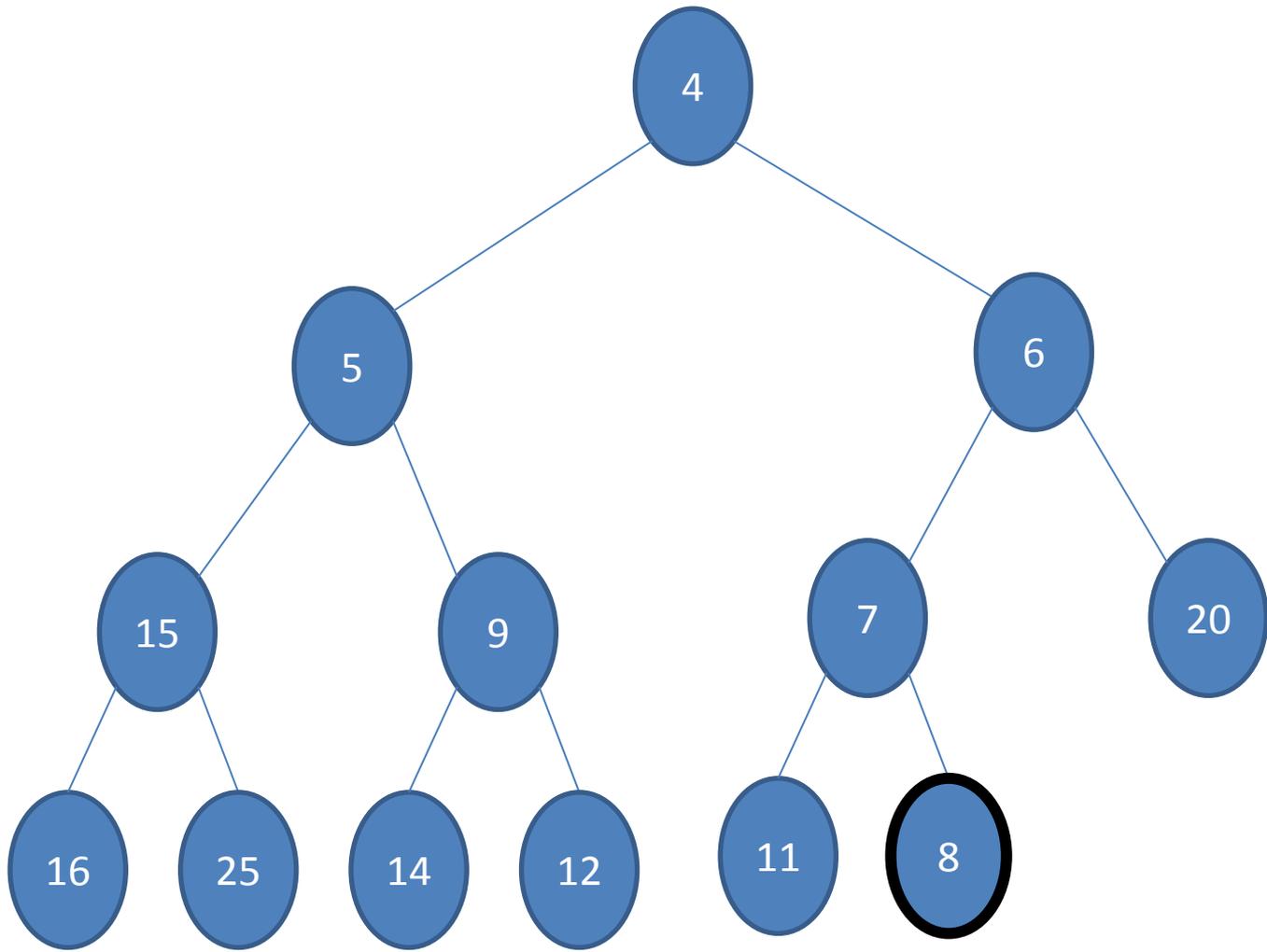
heap



Insert 8

Example: adding to a heap

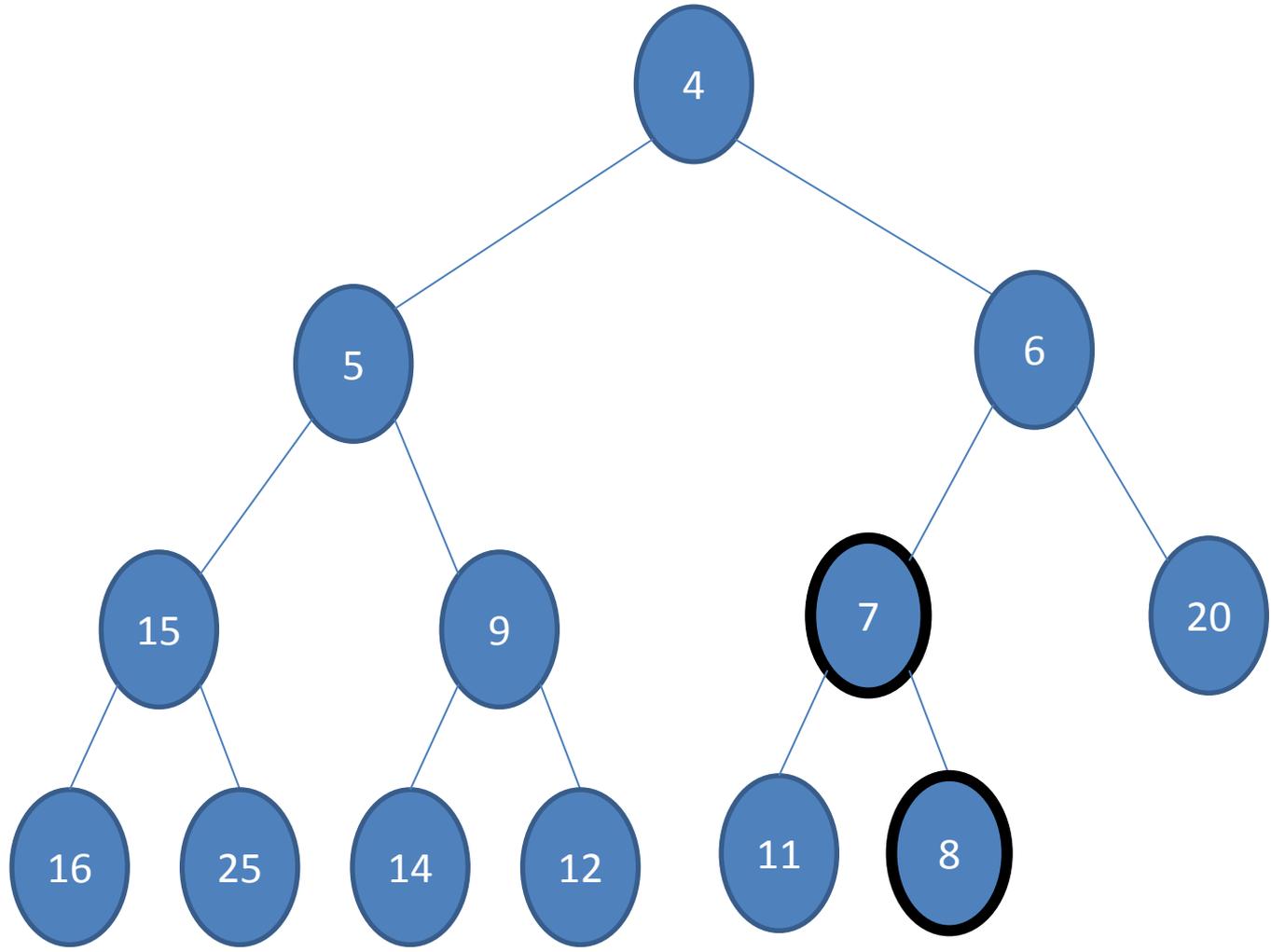
heap



Insert 8

Example: adding to a heap

heap

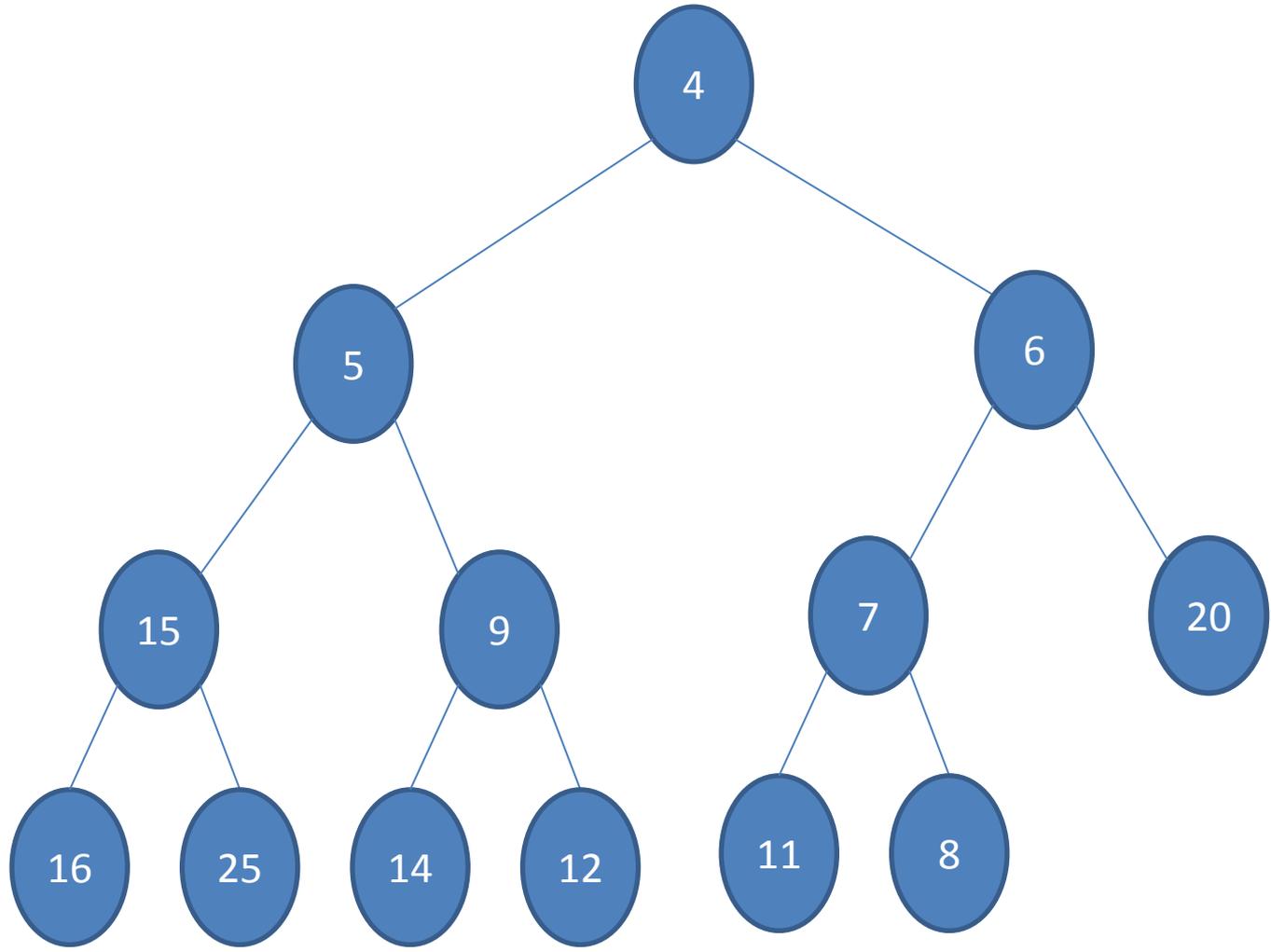


Insert 8

8 is greater than parent (7) ... done

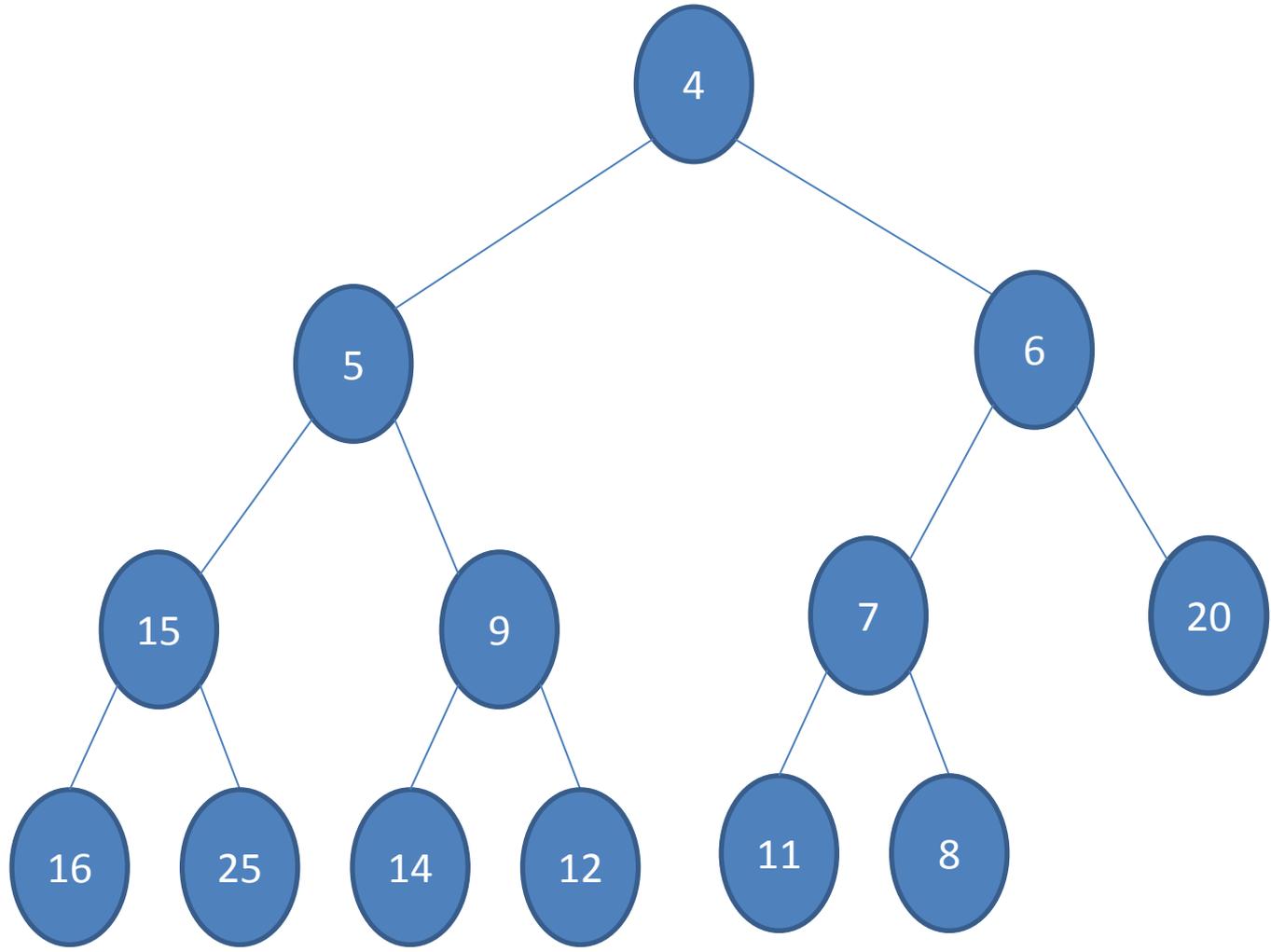
Example: adding to a heap

heap



Example: adding to a heap

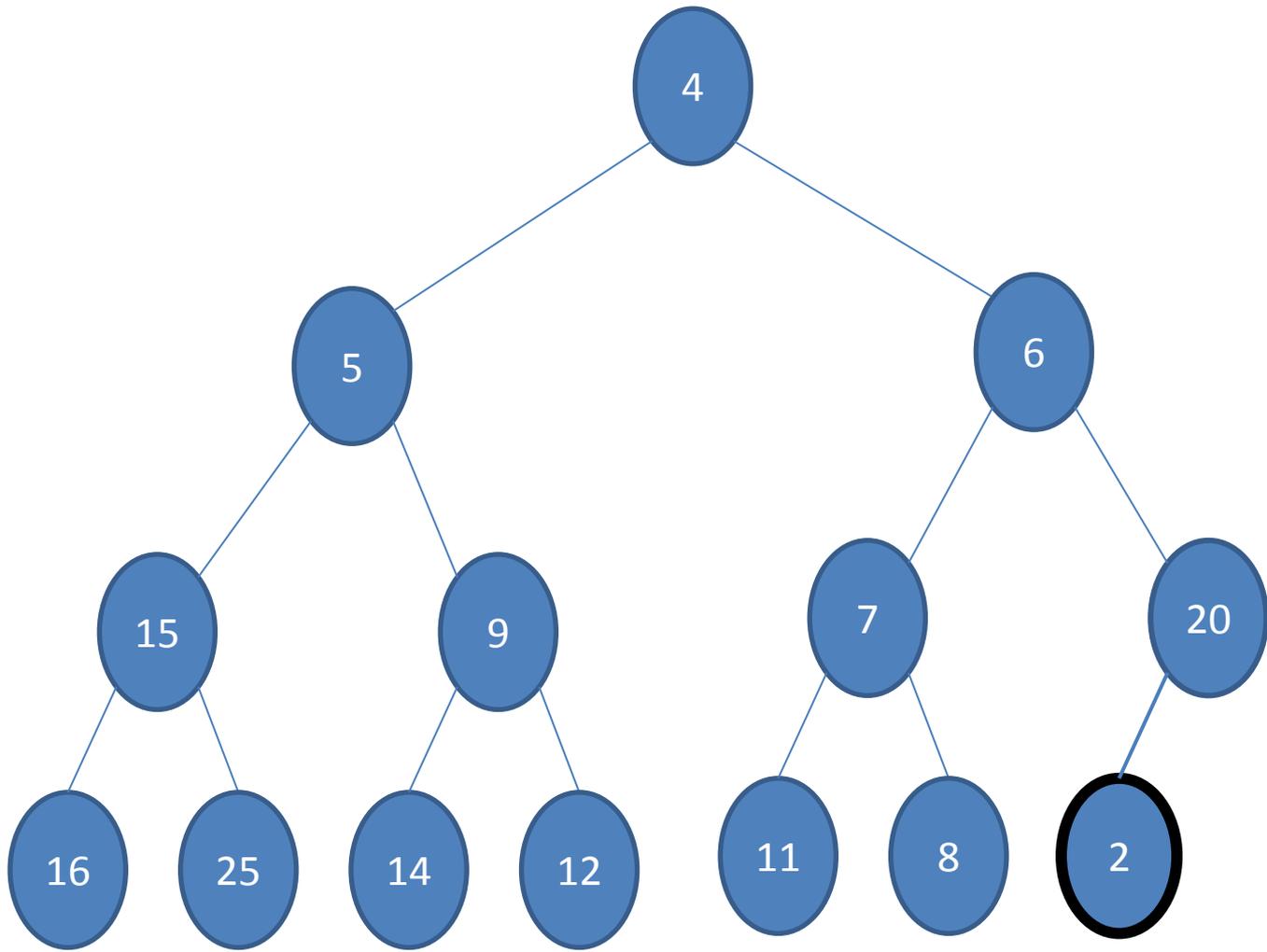
heap



Insert 2

Example: adding to a heap

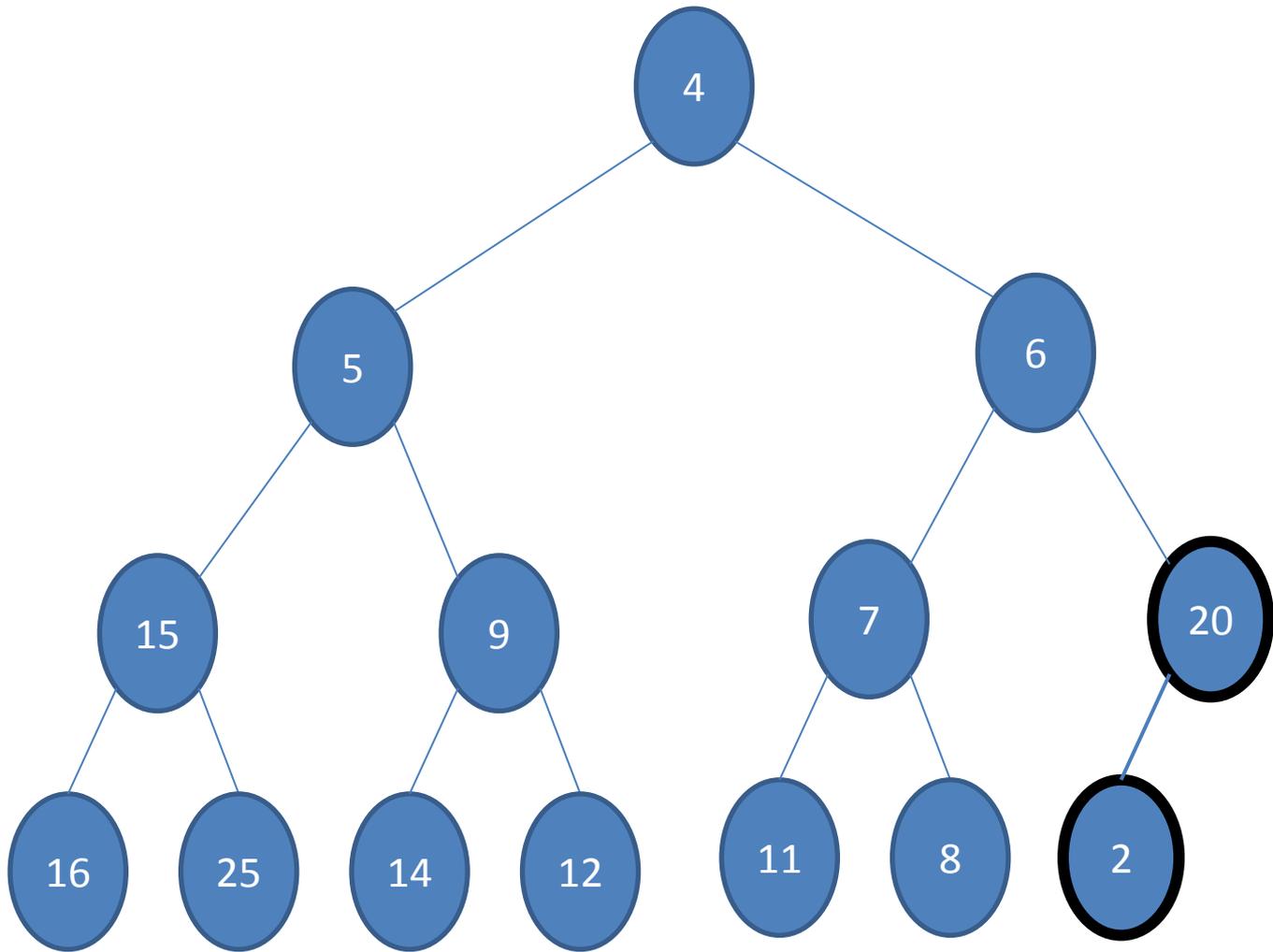
heap



Insert 2

Example: adding to a heap

heap

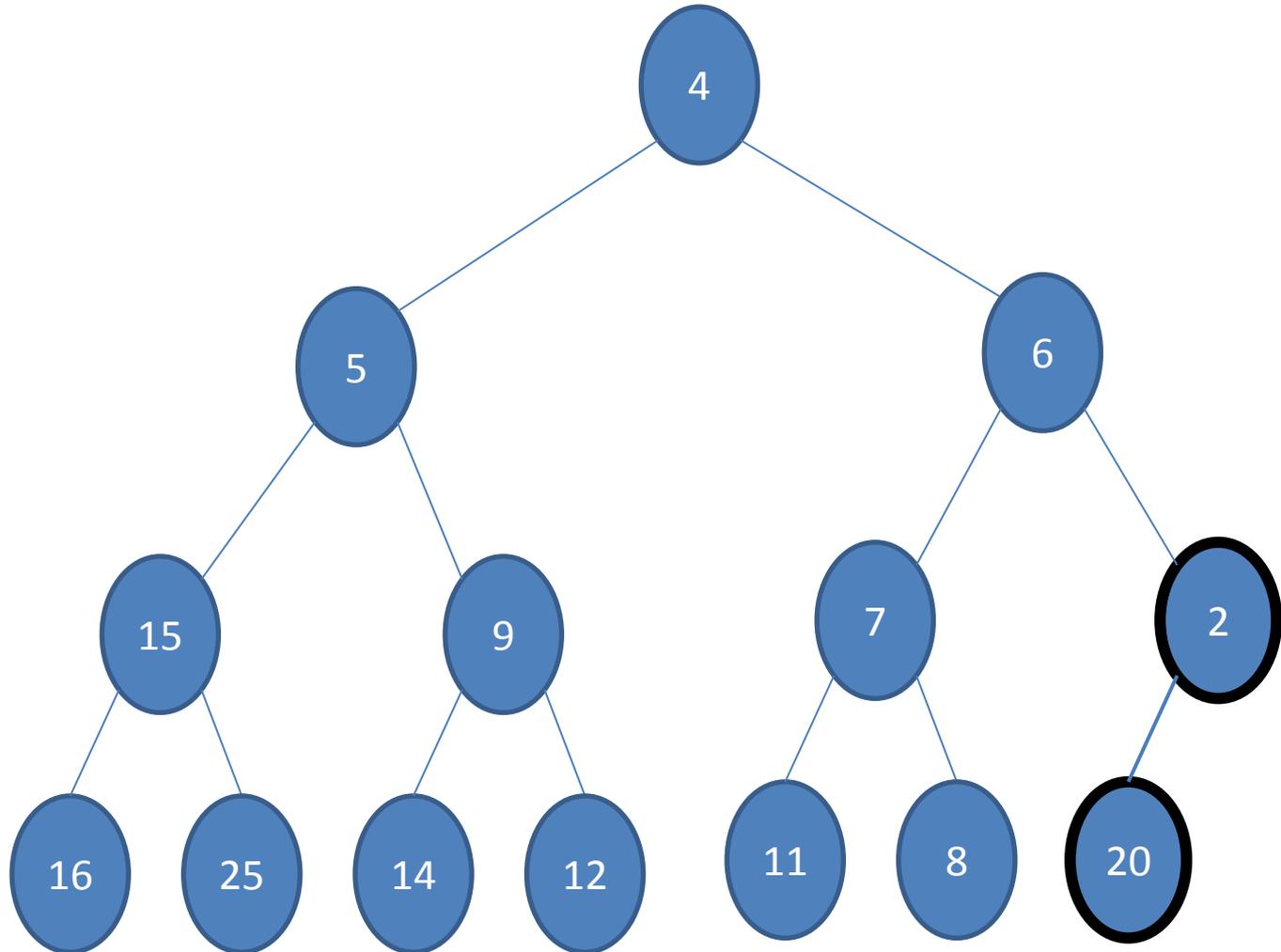


Insert 2

2 is less than parent 20

Example: adding to a heap

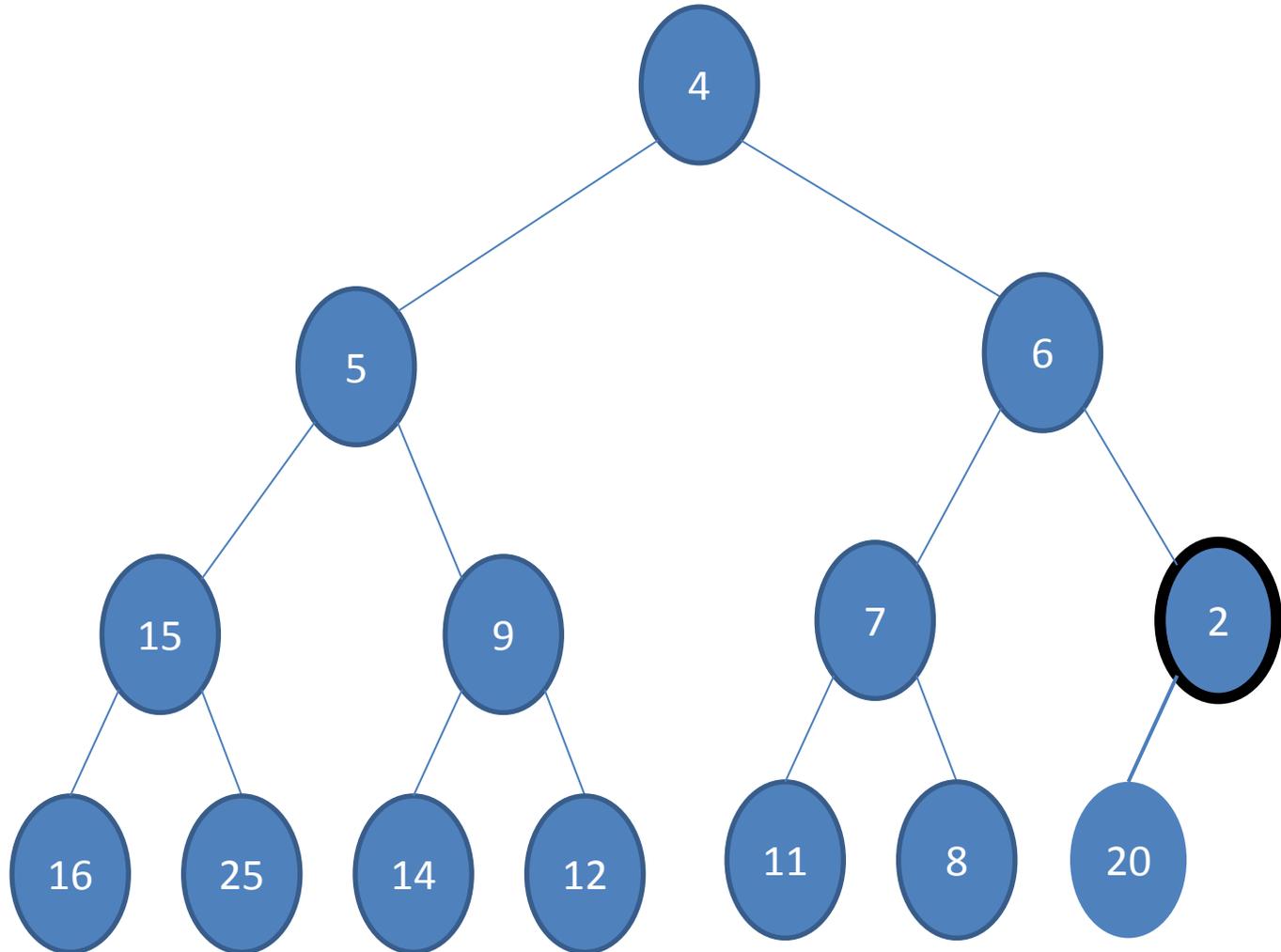
heap



Insert 2

Example: adding to a heap

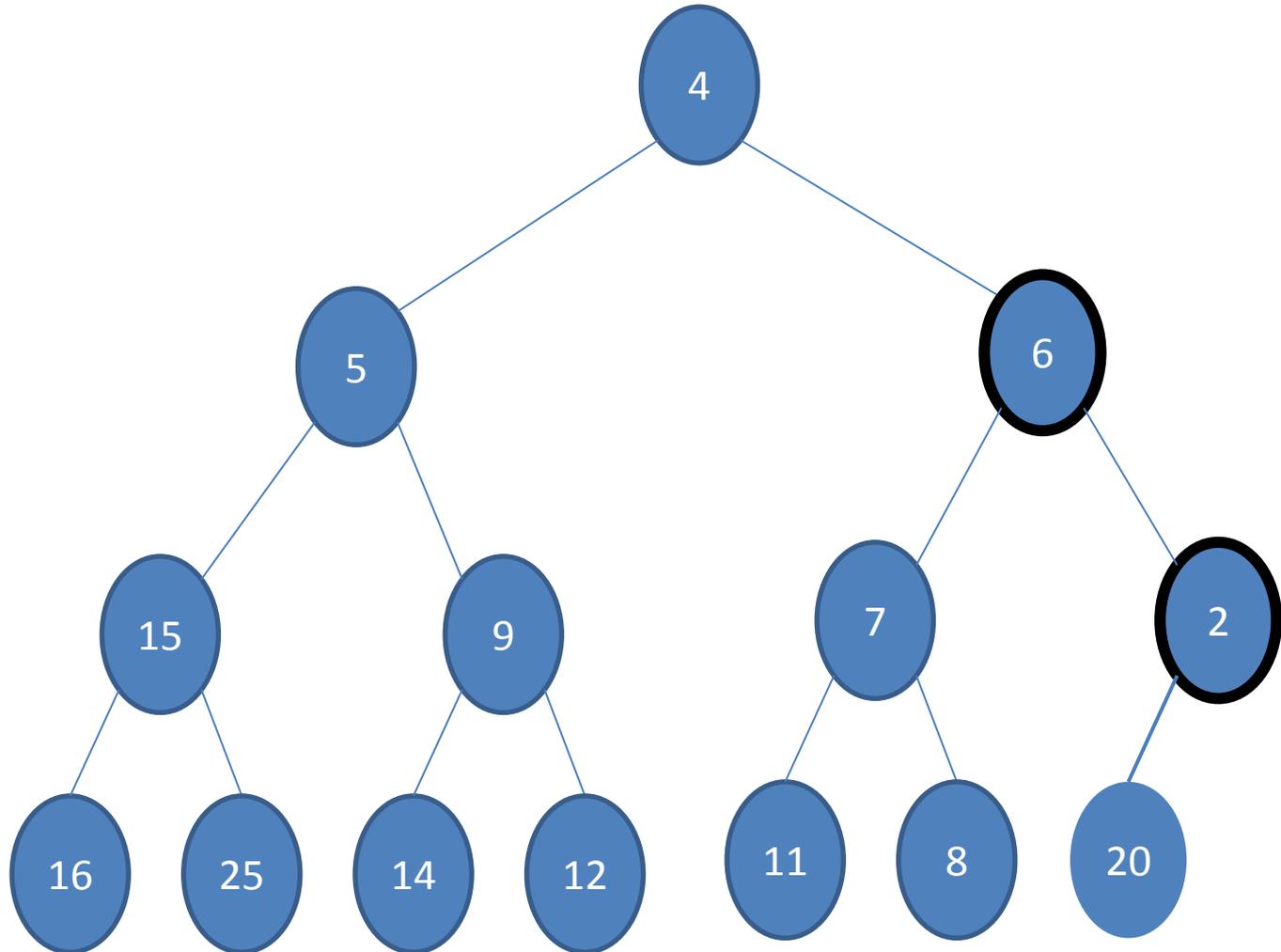
heap



Insert 2

Example: adding to a heap

heap

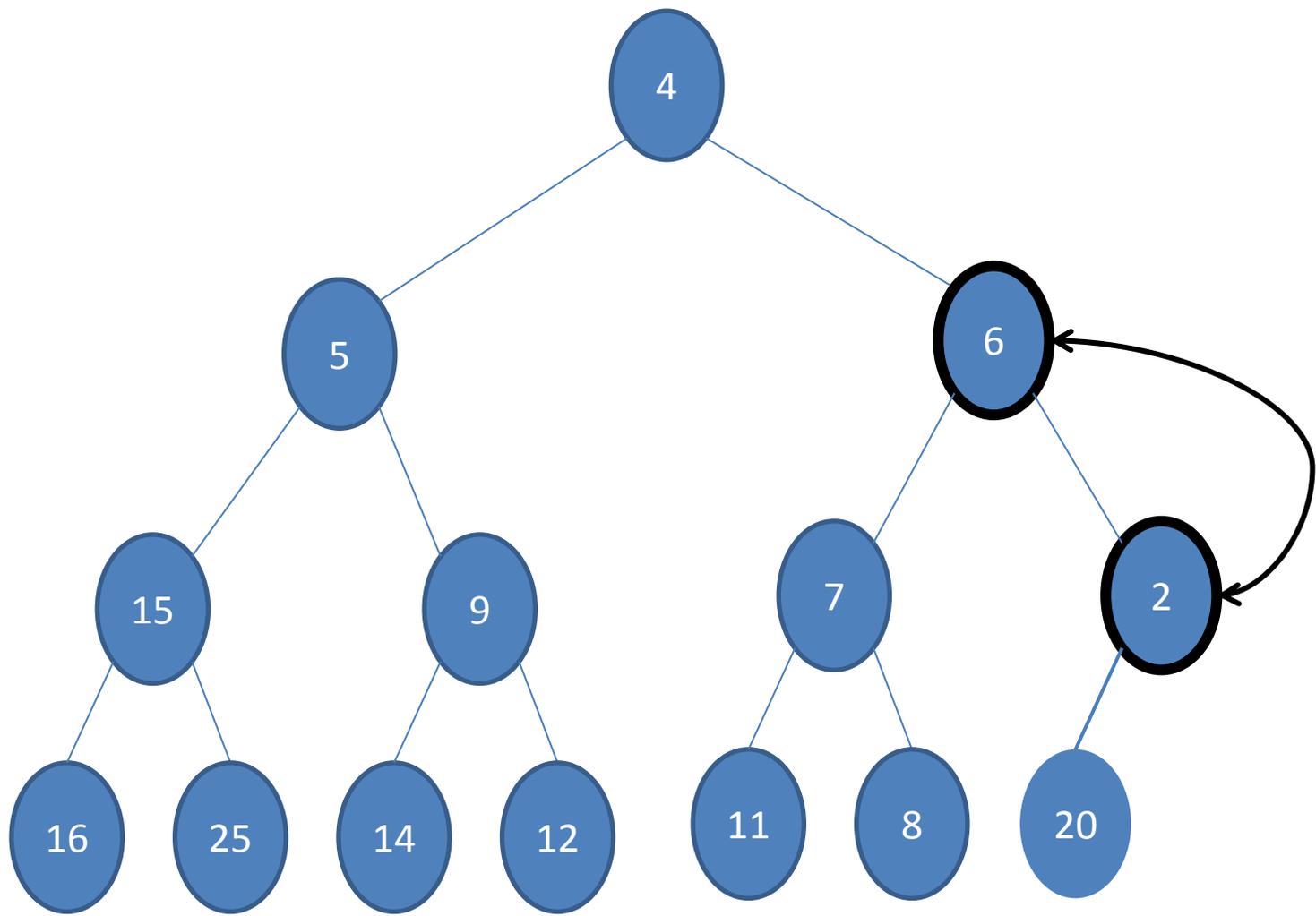


Insert 2

2 is less than parent 6

Example: adding to a heap

heap

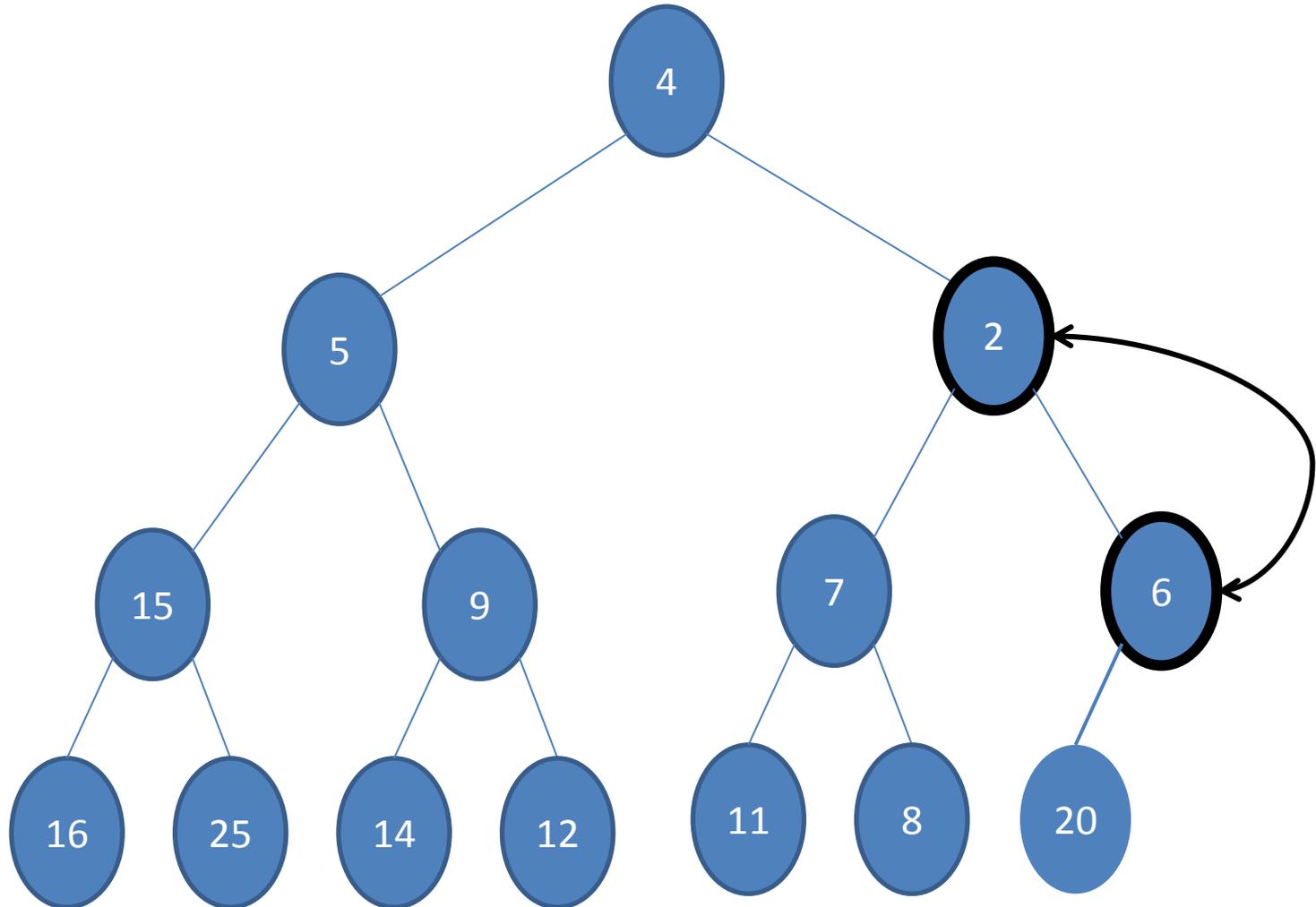


Insert 2

2 is less than parent 6 ... swap!

Example: adding to a heap

heap

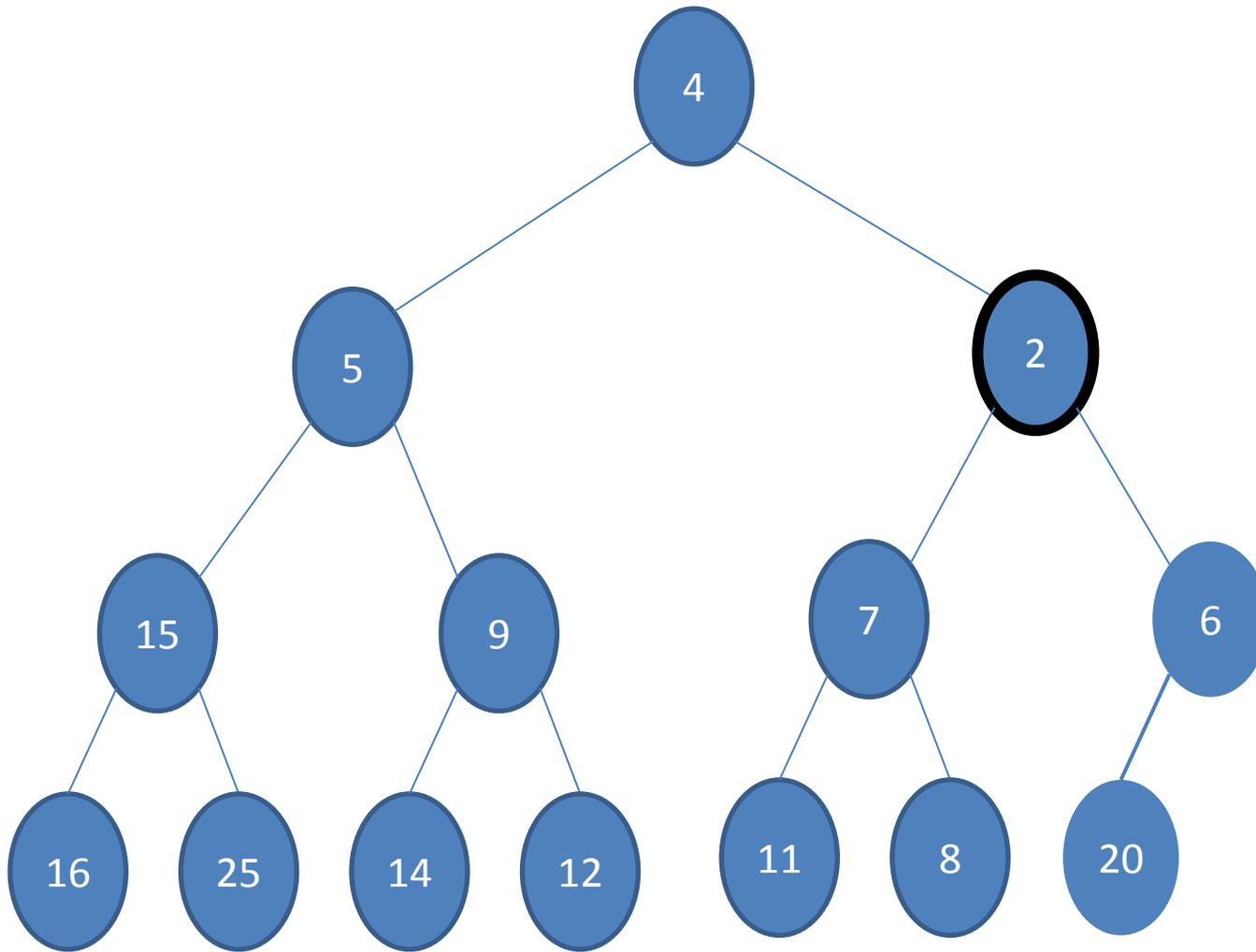


Insert 2

2 is less than parent 6 ... swap!

Example: adding to a heap

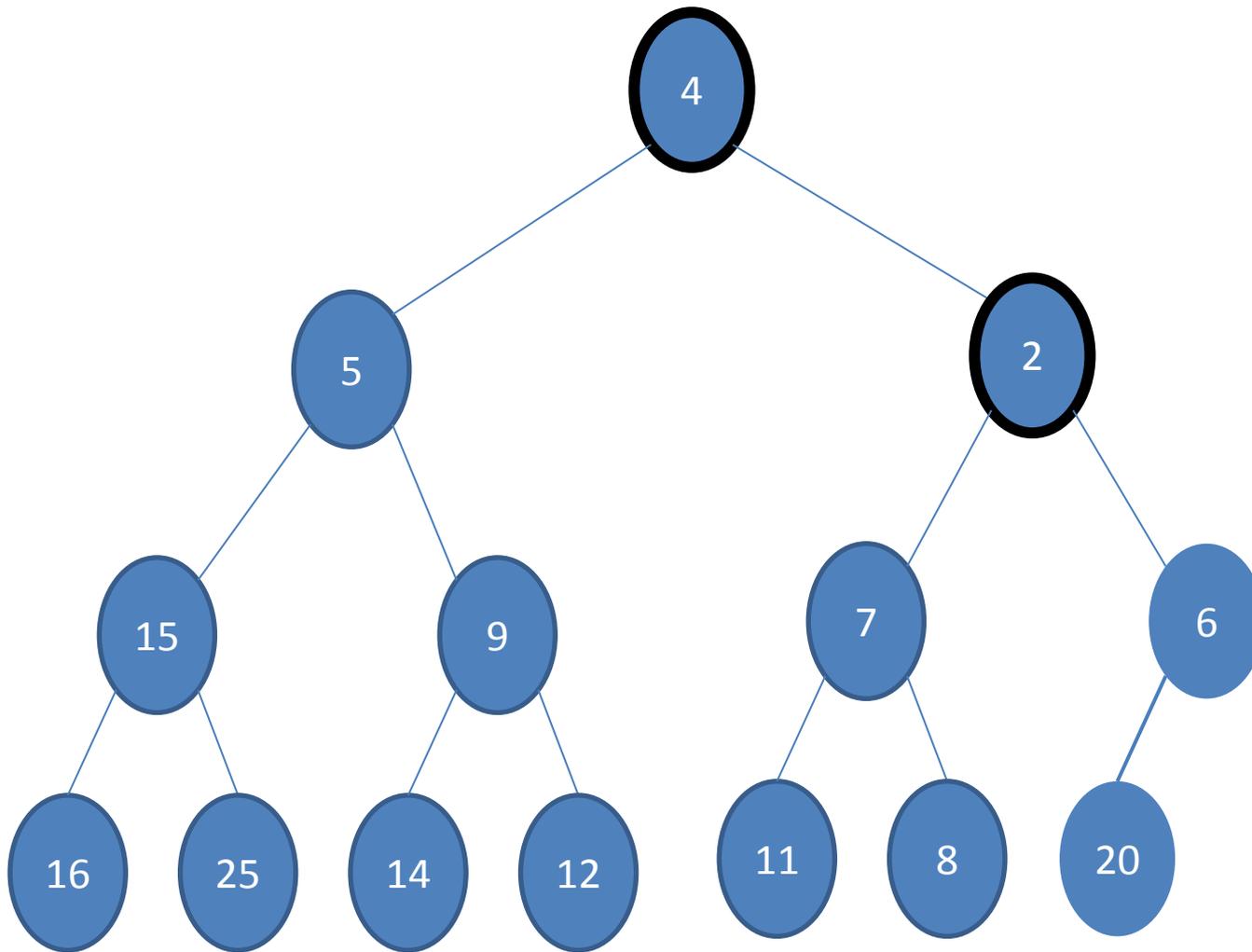
heap



Insert 2

Example: adding to a heap

heap

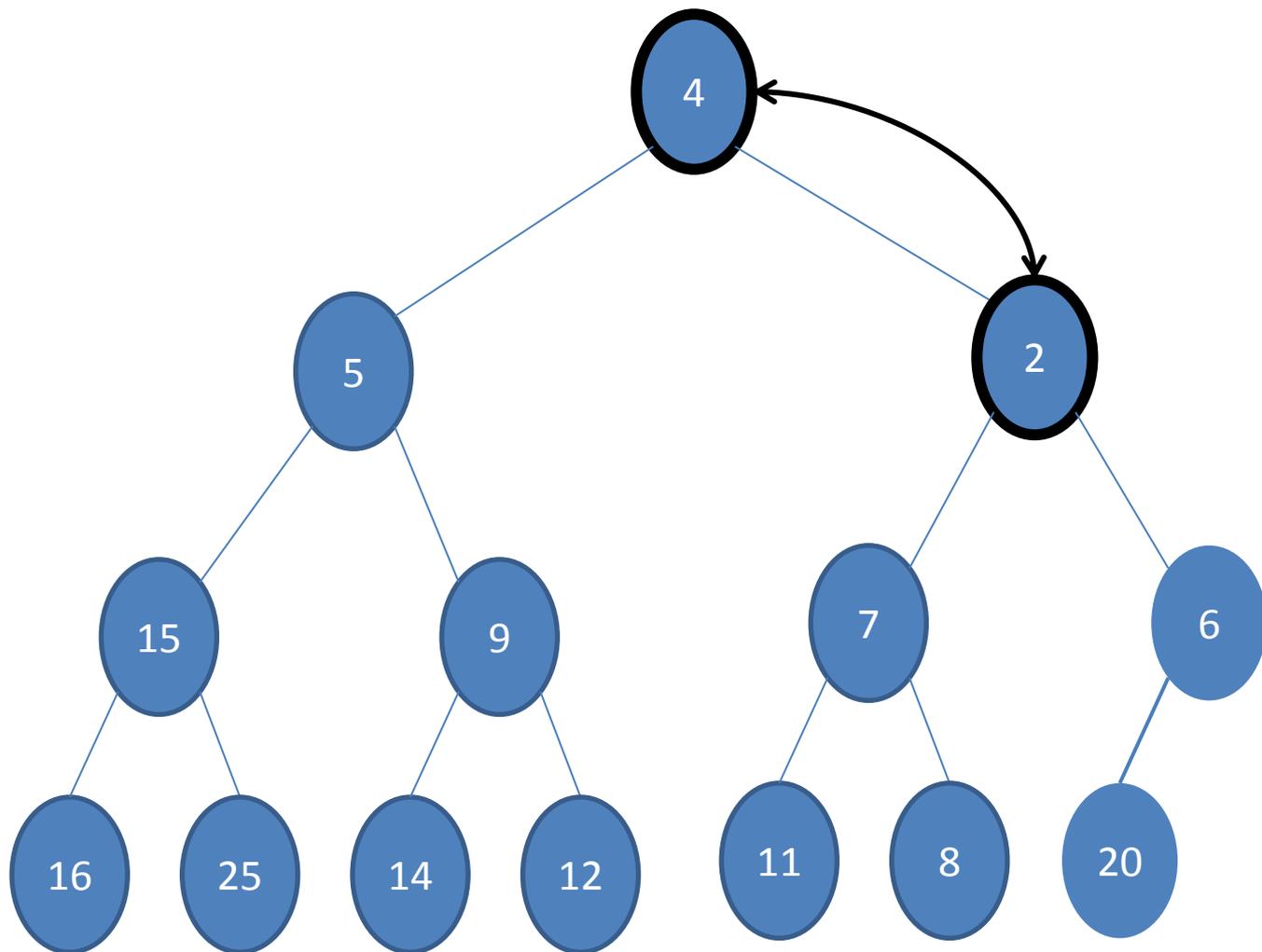


Insert 2

2 is less than parent 4

Example: adding to a heap

heap

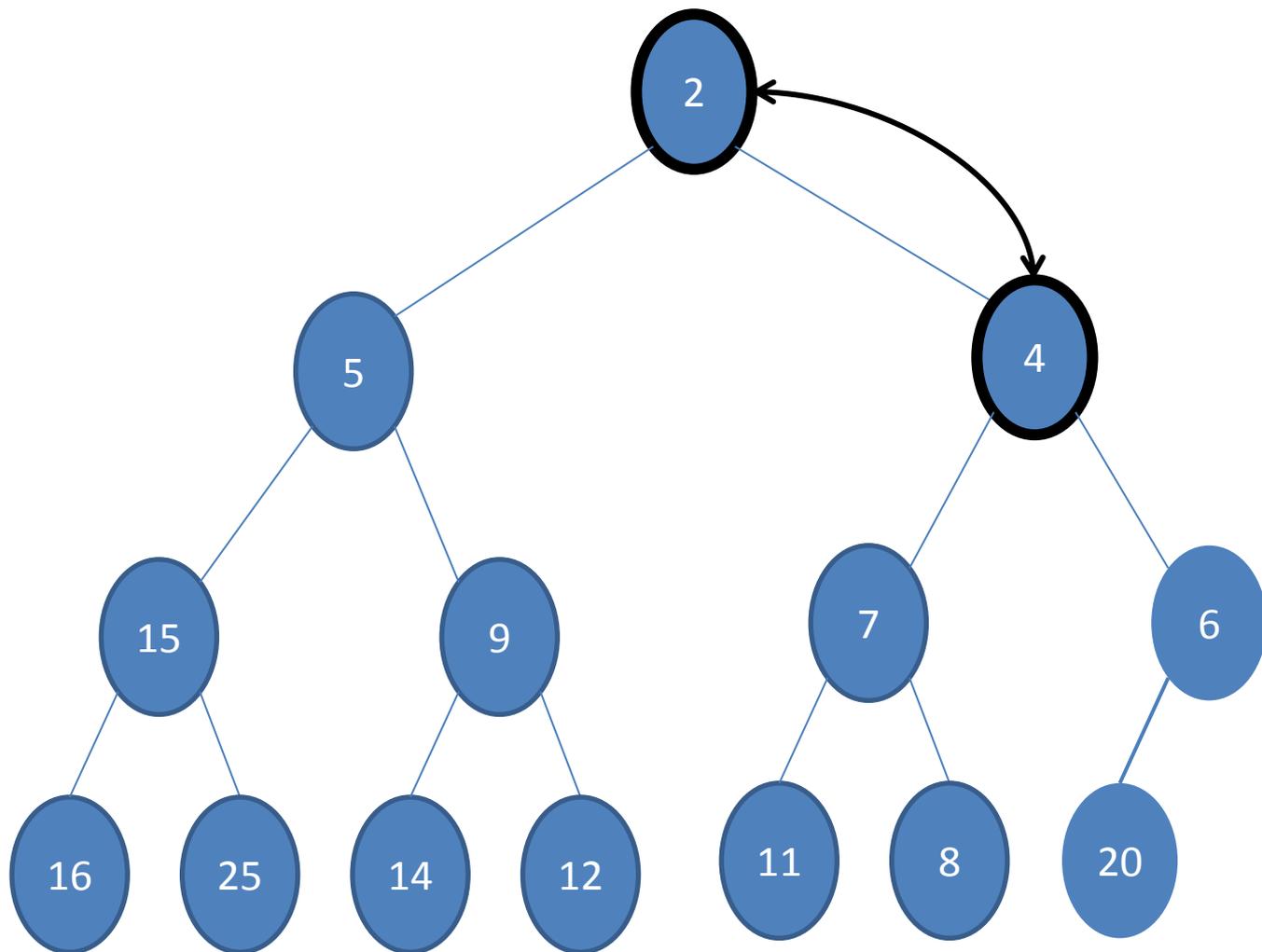


Insert 2

2 is less than parent 4 ... swap!

Example: adding to a heap

heap

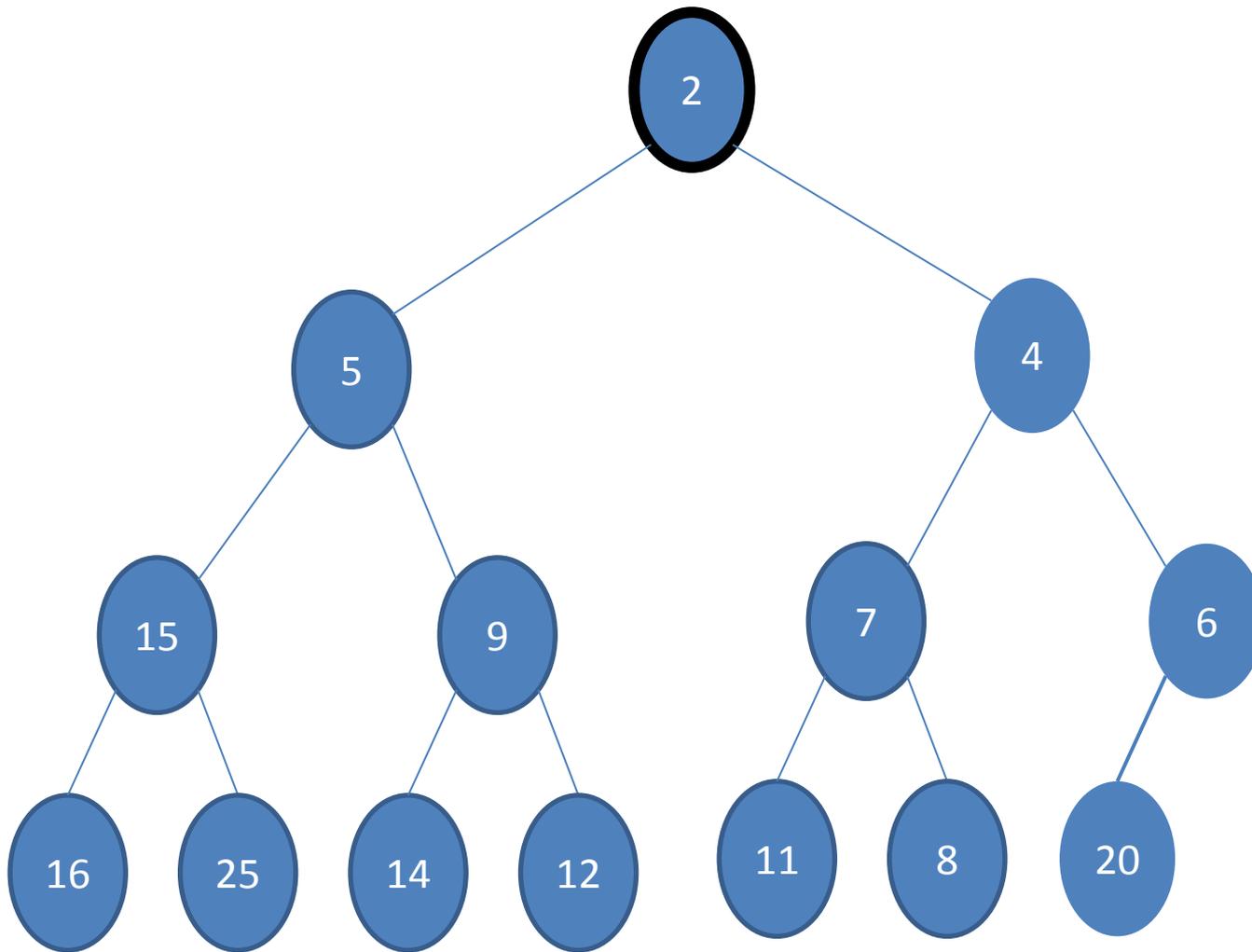


Insert 2

2 is less than parent 4 ... swap!

Example: adding to a heap

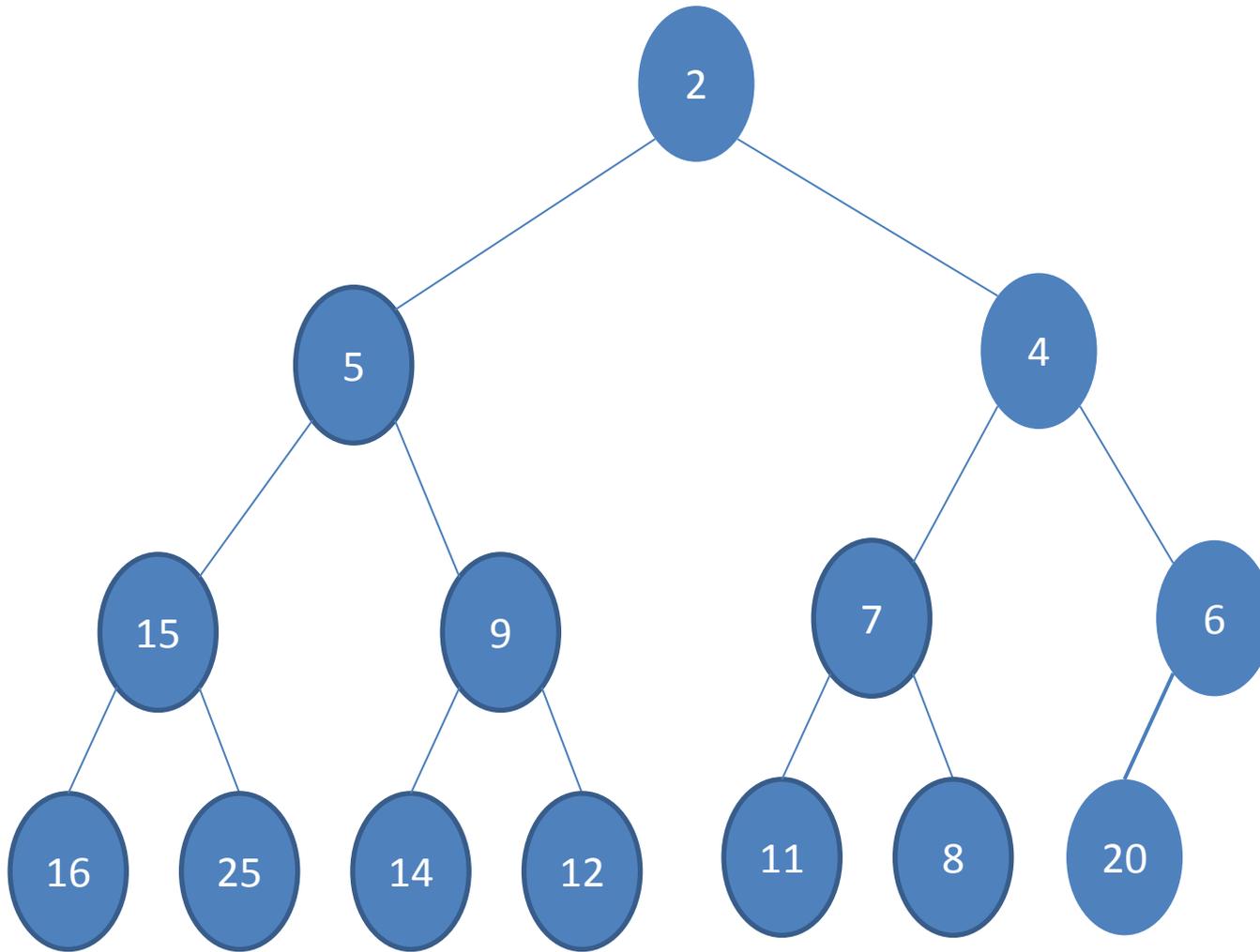
heap



Insert 2

Example: adding to a heap

heap

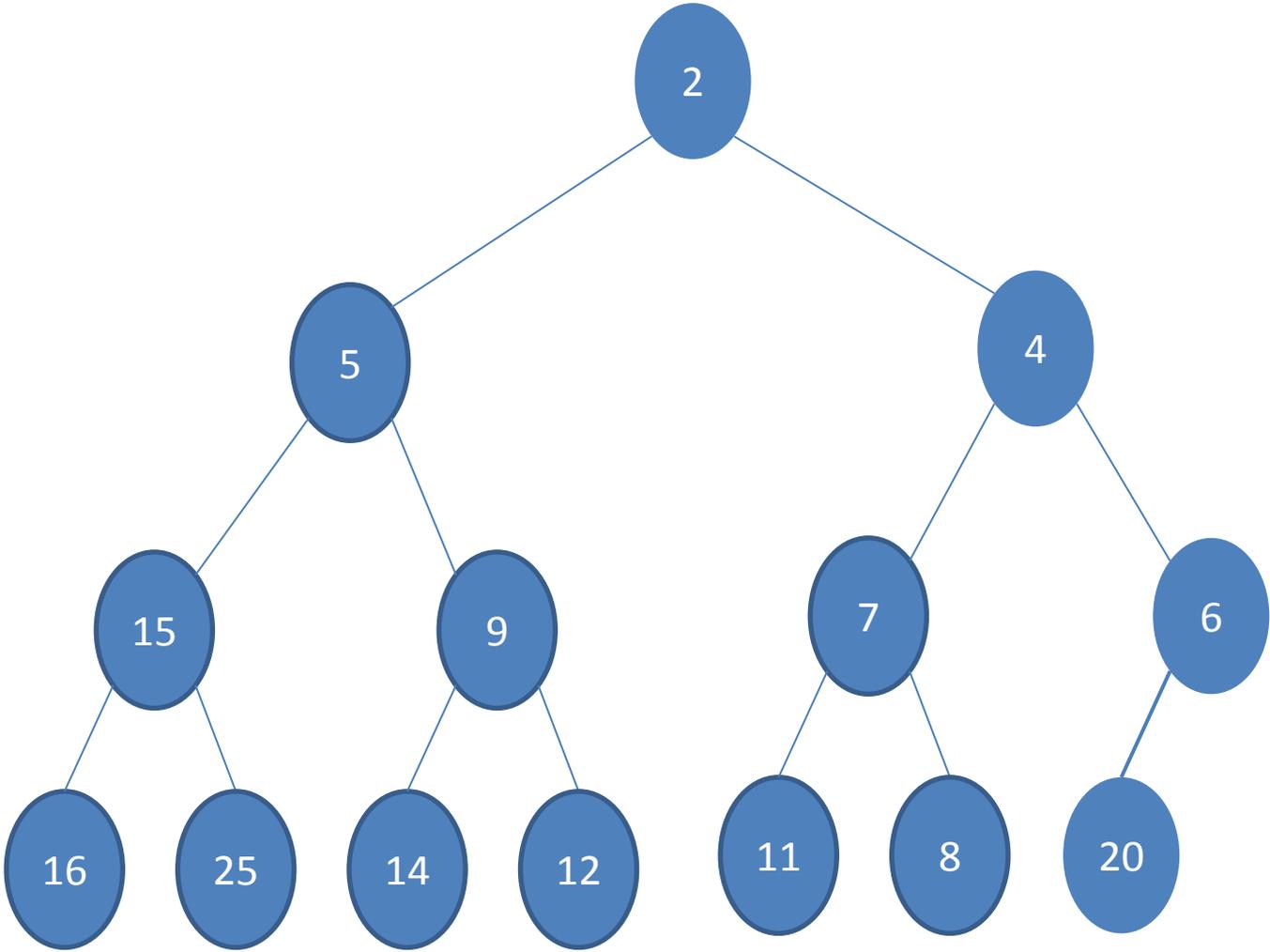


Insert 2

Done!

Example: adding to a heap

heap

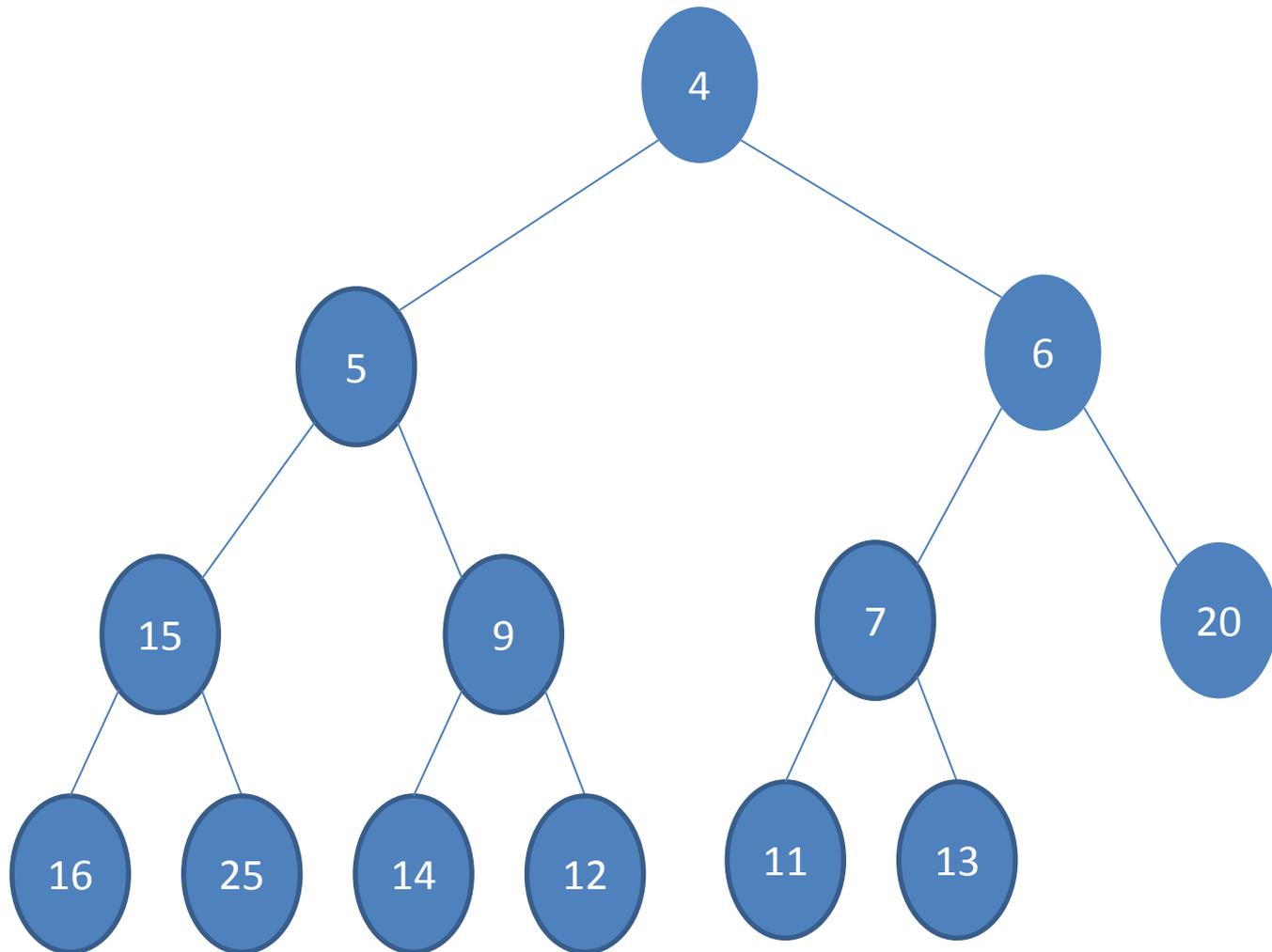


Example: removal from a heap

NOTE: it is a heap in a different state

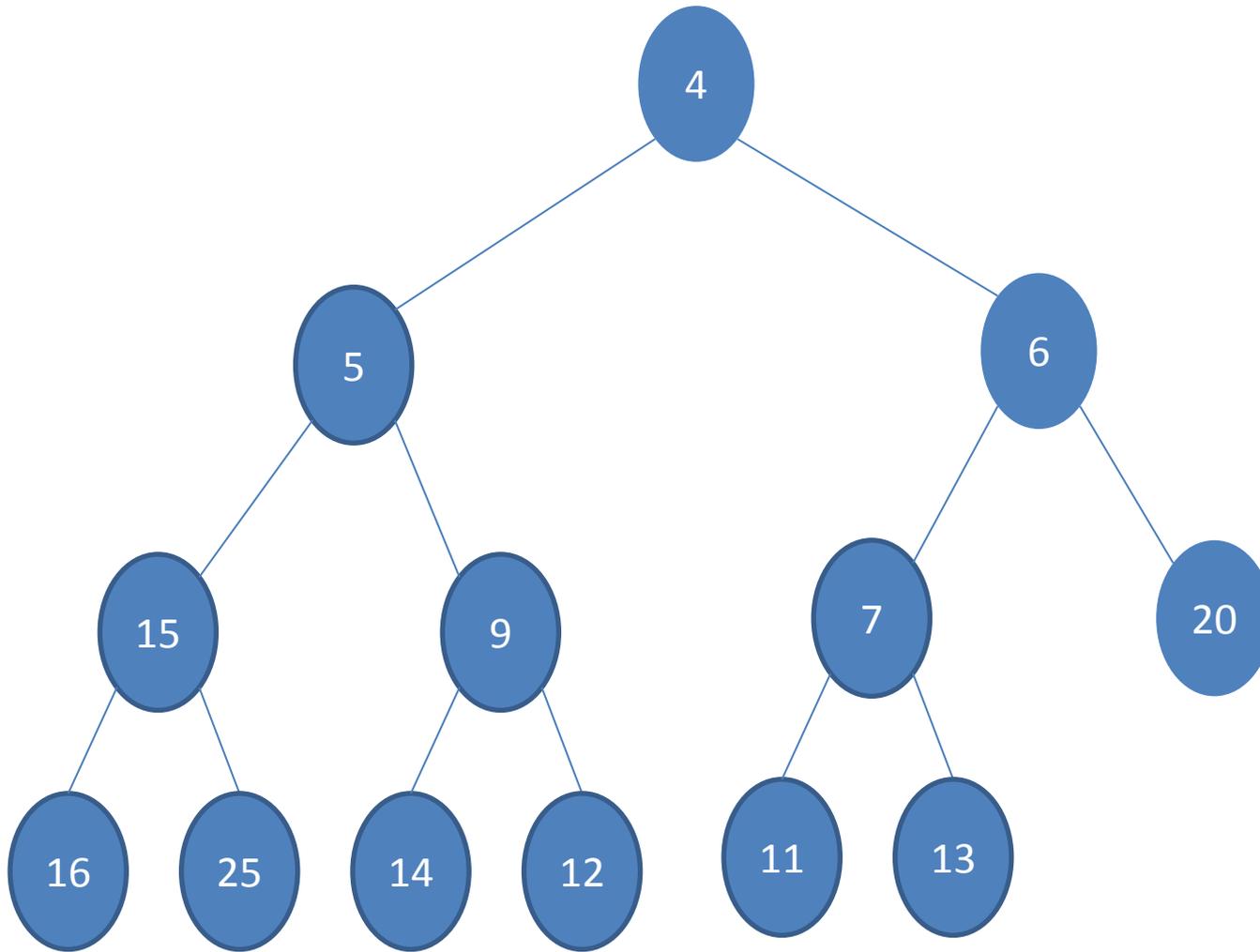
Example: removal from a heap

heap



Example: removal from a heap

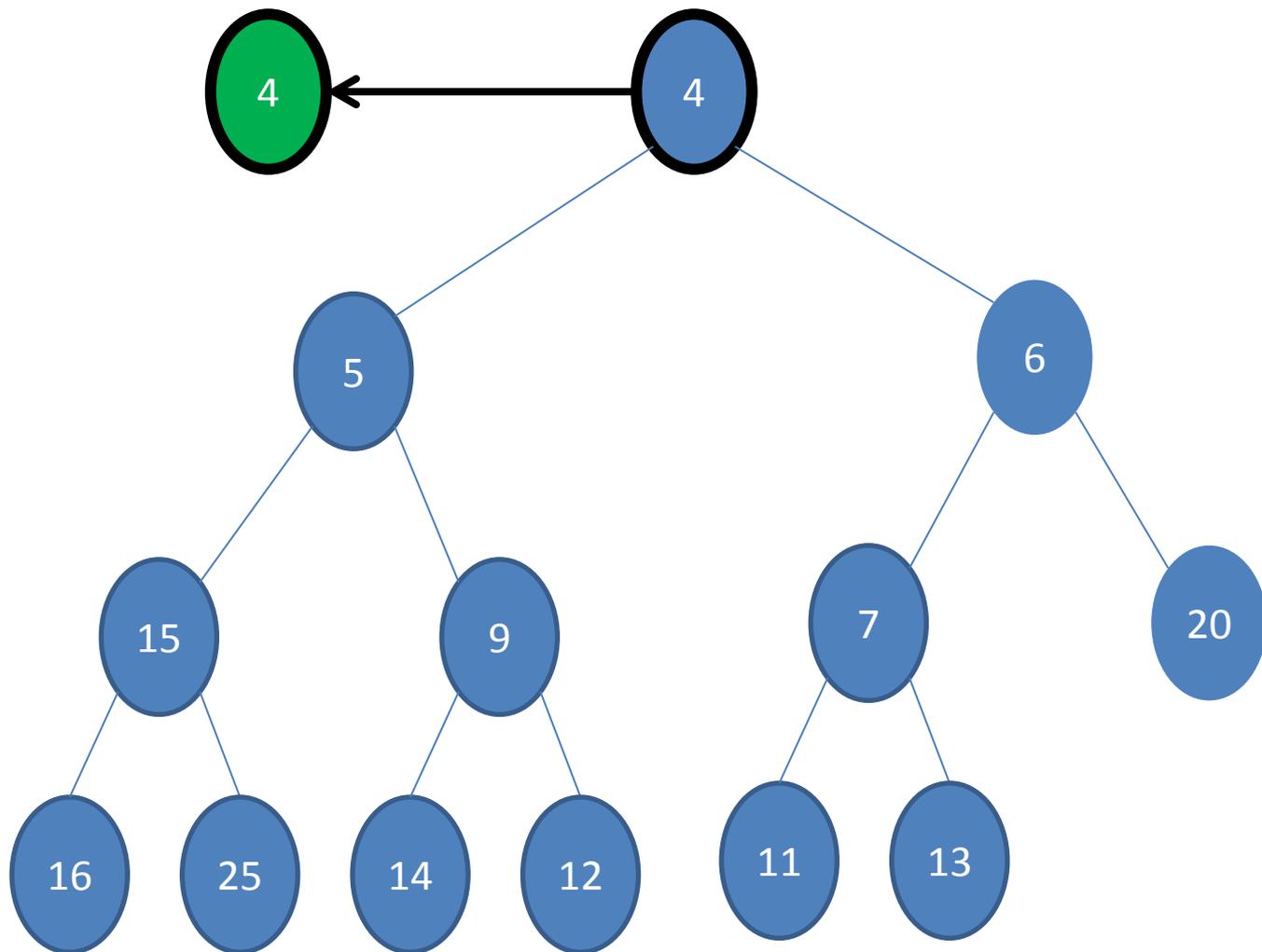
heap



Save off top of heap

Example: removal from a heap

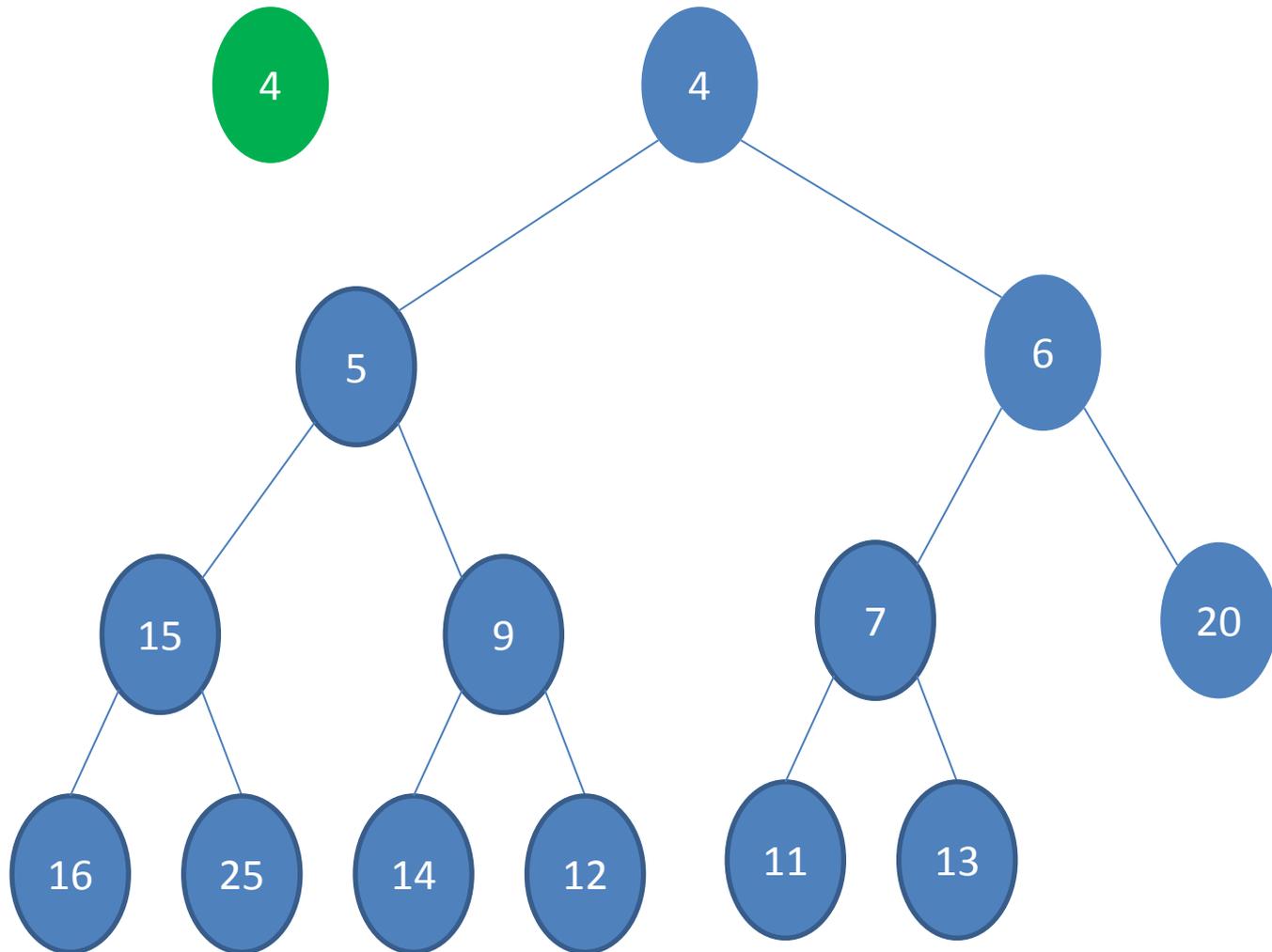
heap



Save off top of heap

Example: removal from a heap

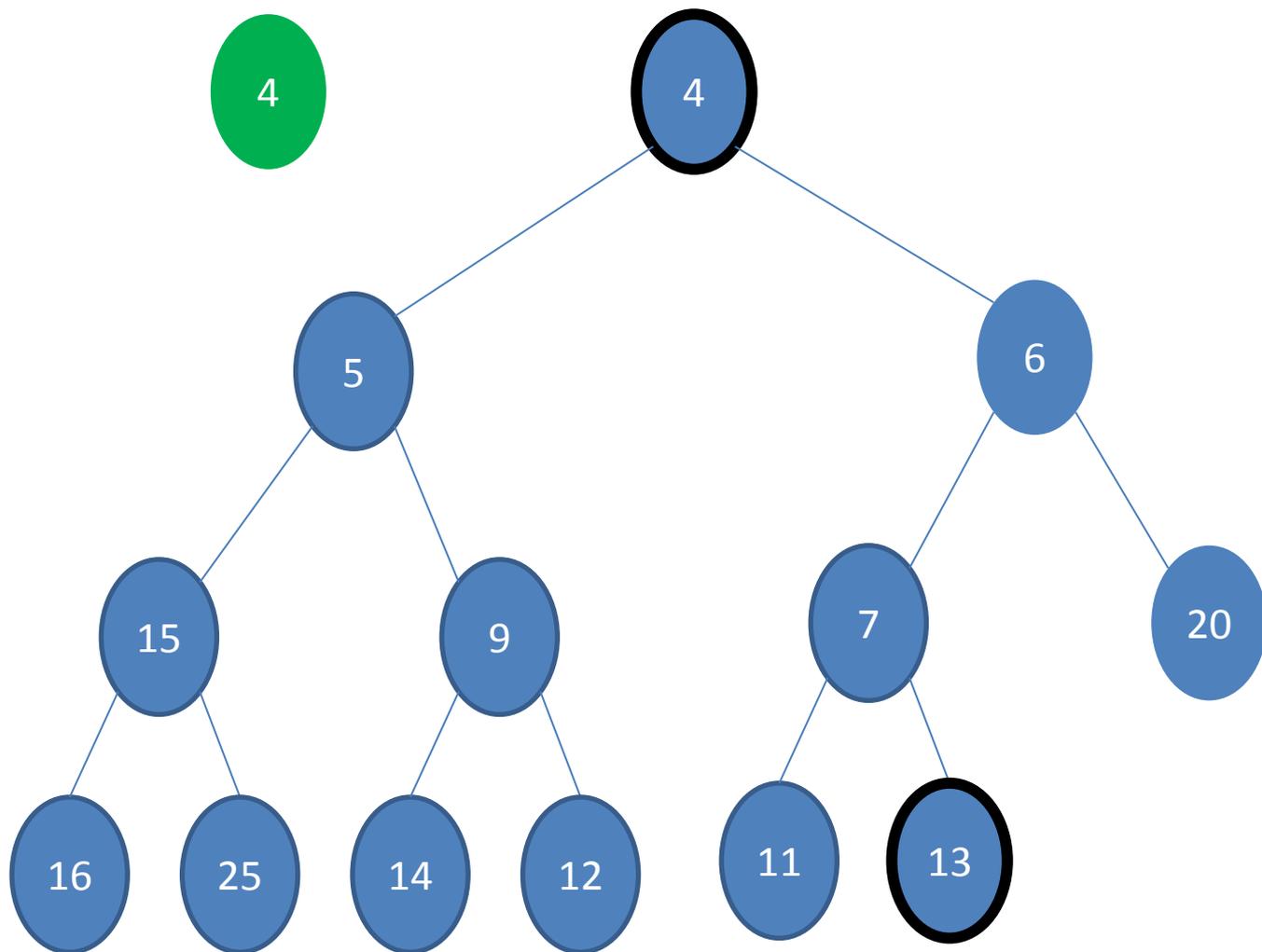
heap



Copy last item in heap to top of heap

Example: removal from a heap

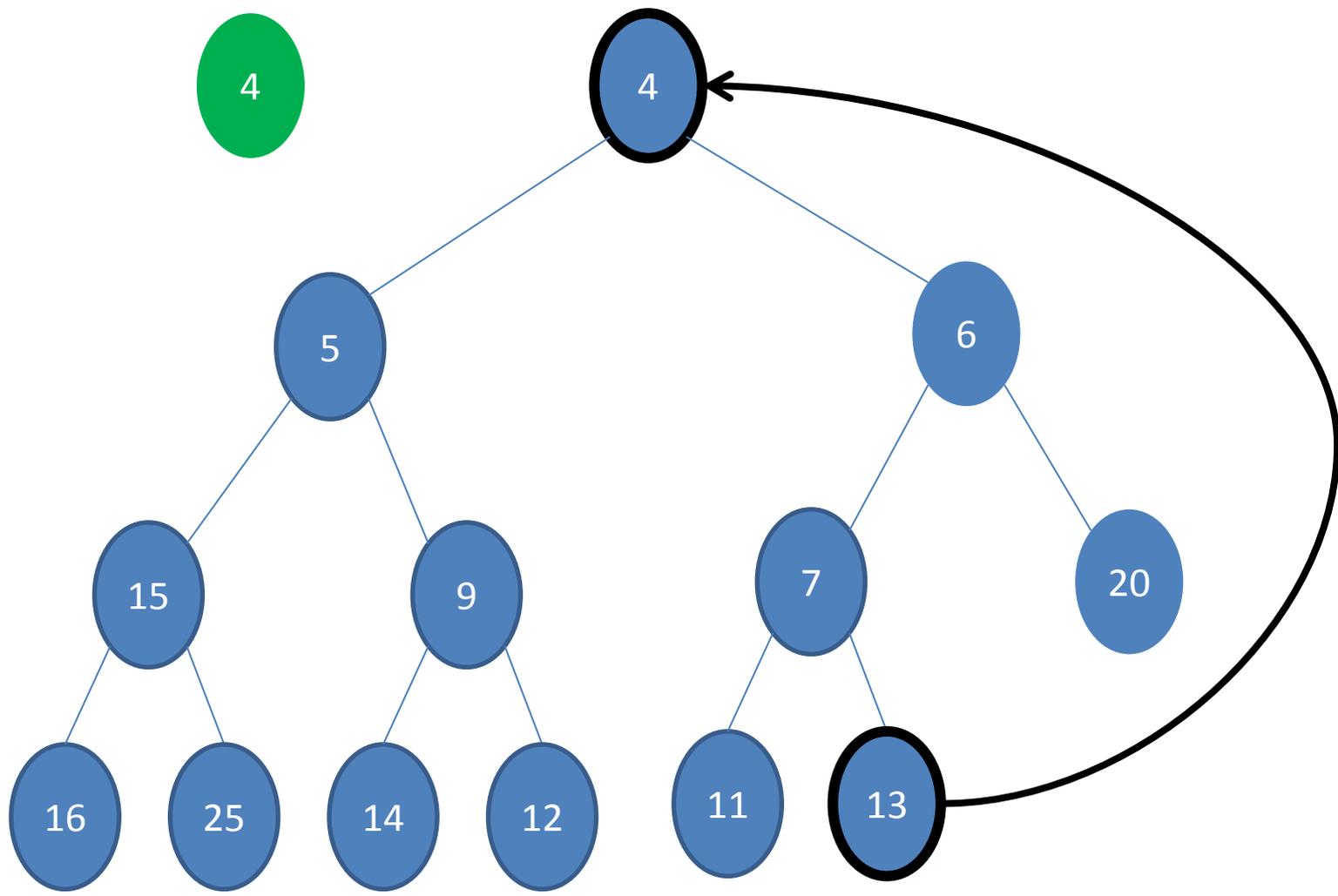
heap



Copy last item in heap to top of heap

Example: removal from a heap

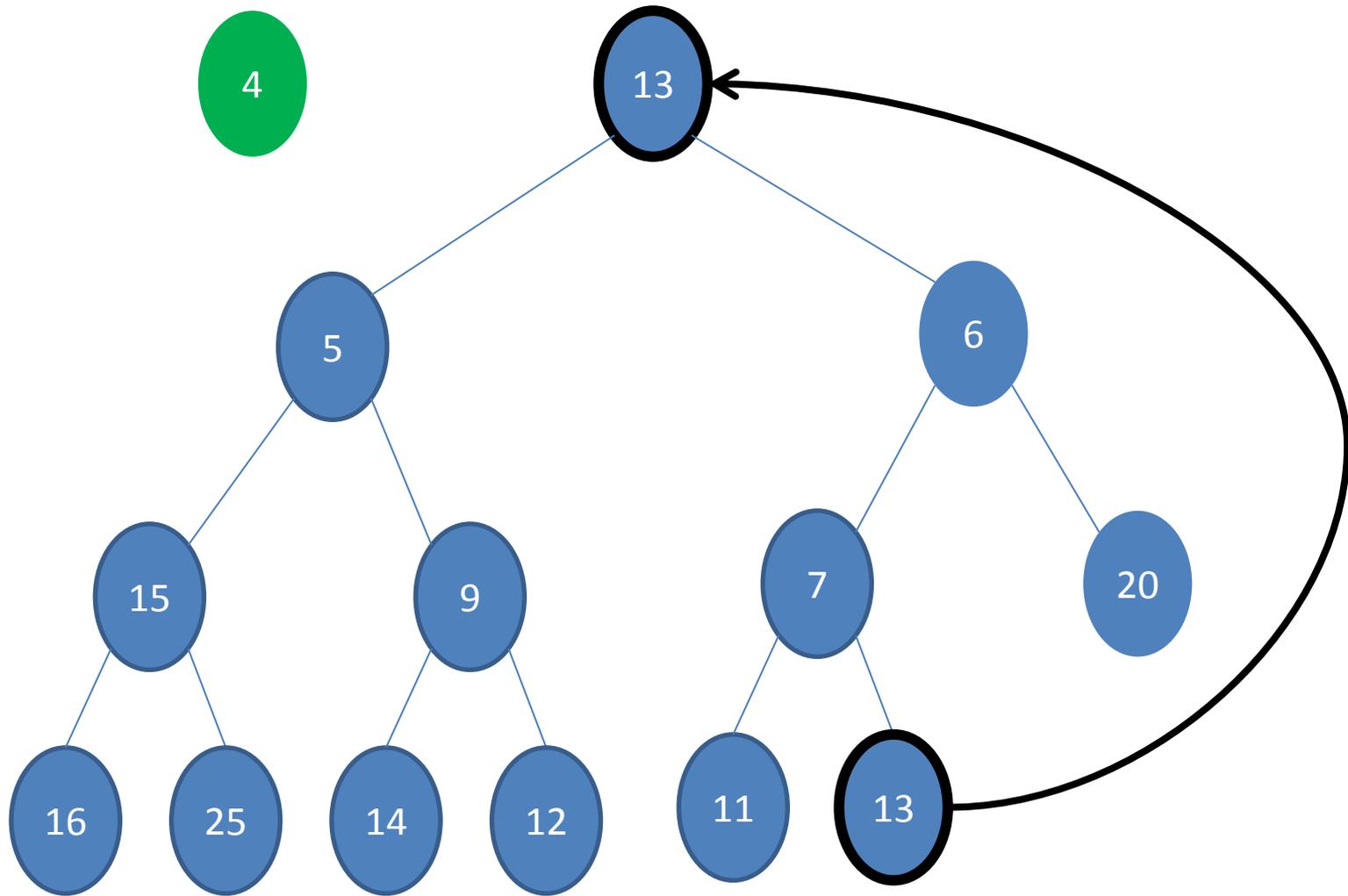
heap



Copy last item in heap to top of heap

Example: removal from a heap

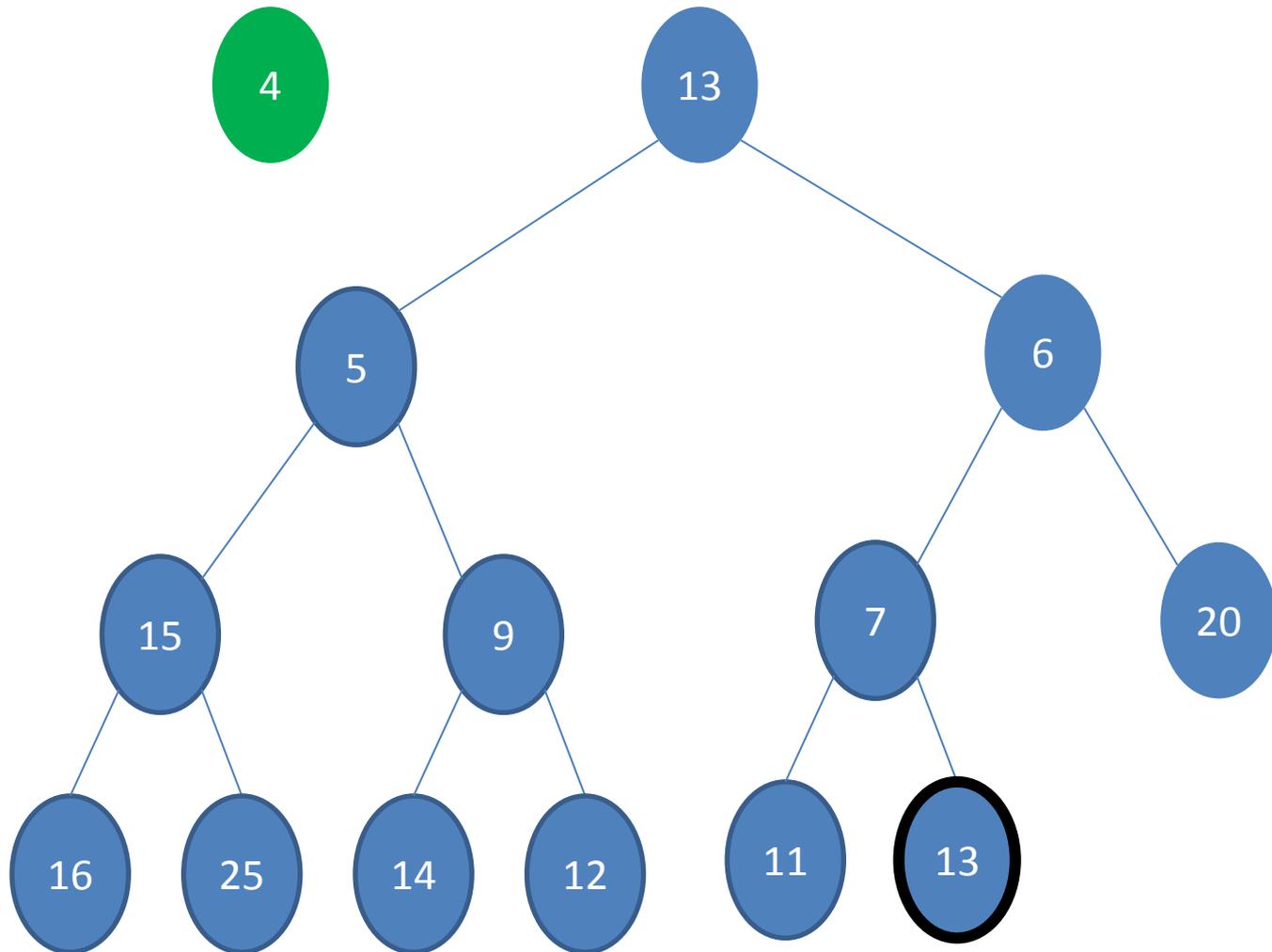
heap



Copy last item in heap to top of heap

Example: removal from a heap

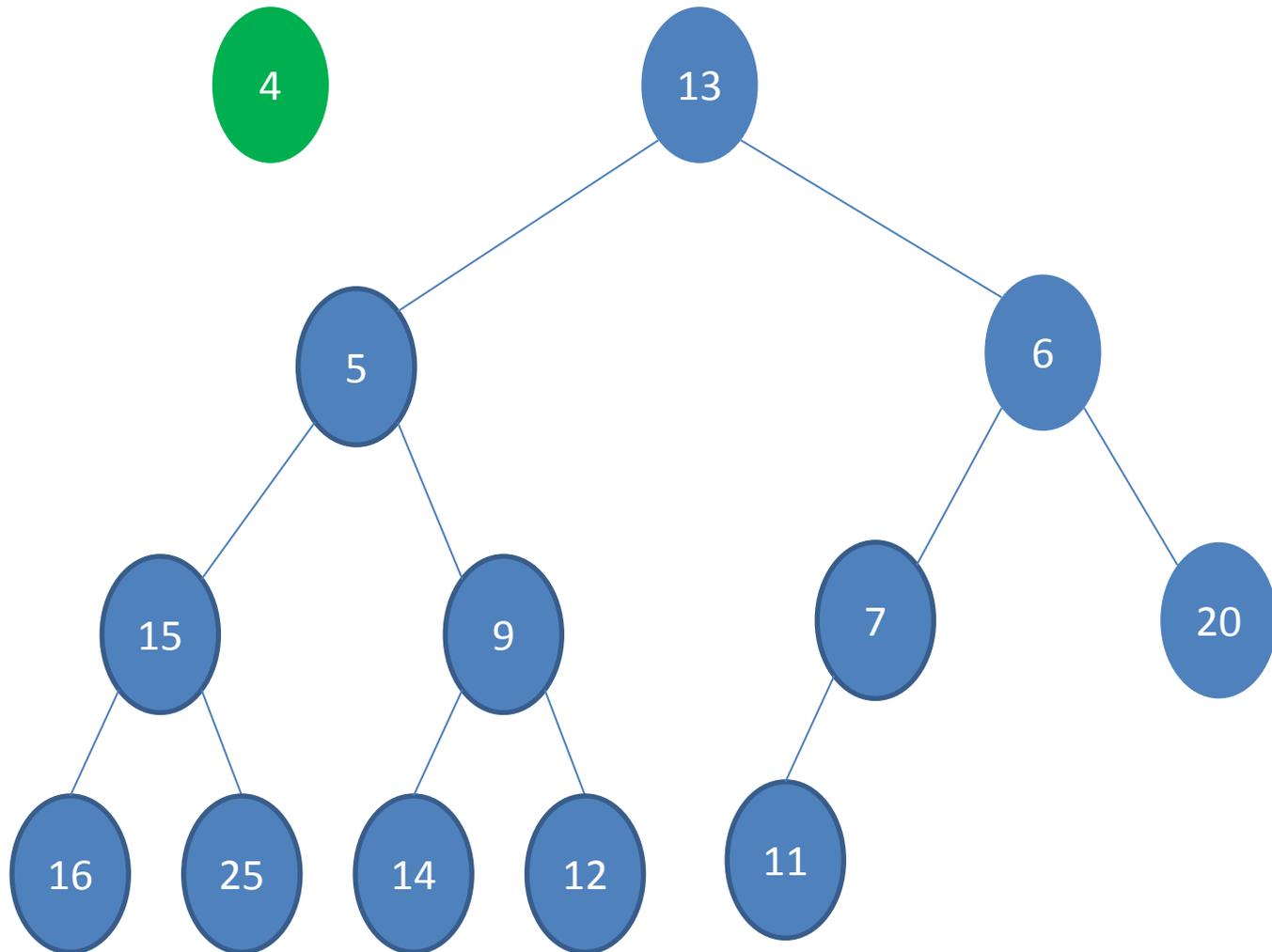
heap



Delete last item in heap

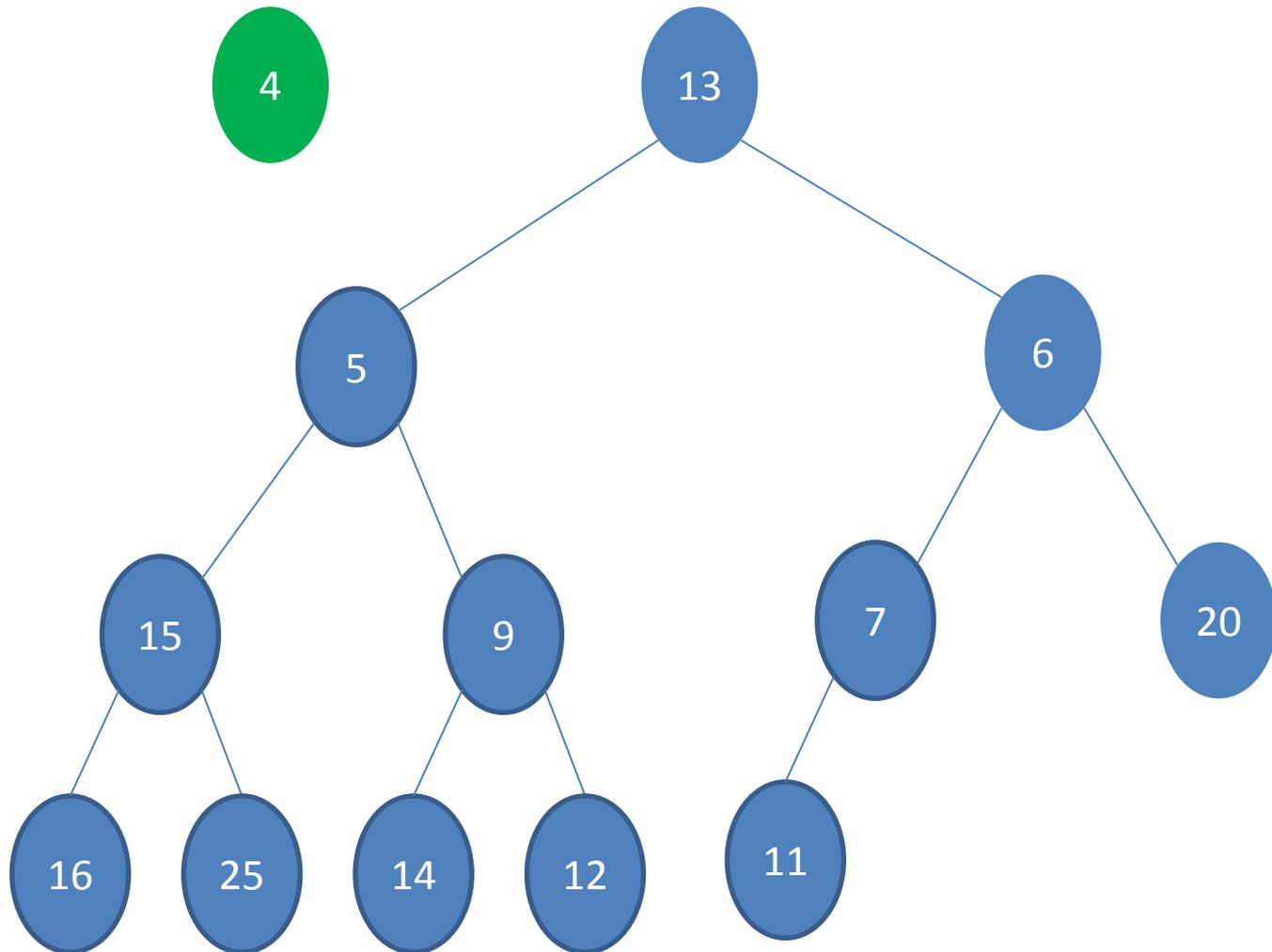
Example: removal from a heap

heap



Example: removal from a heap

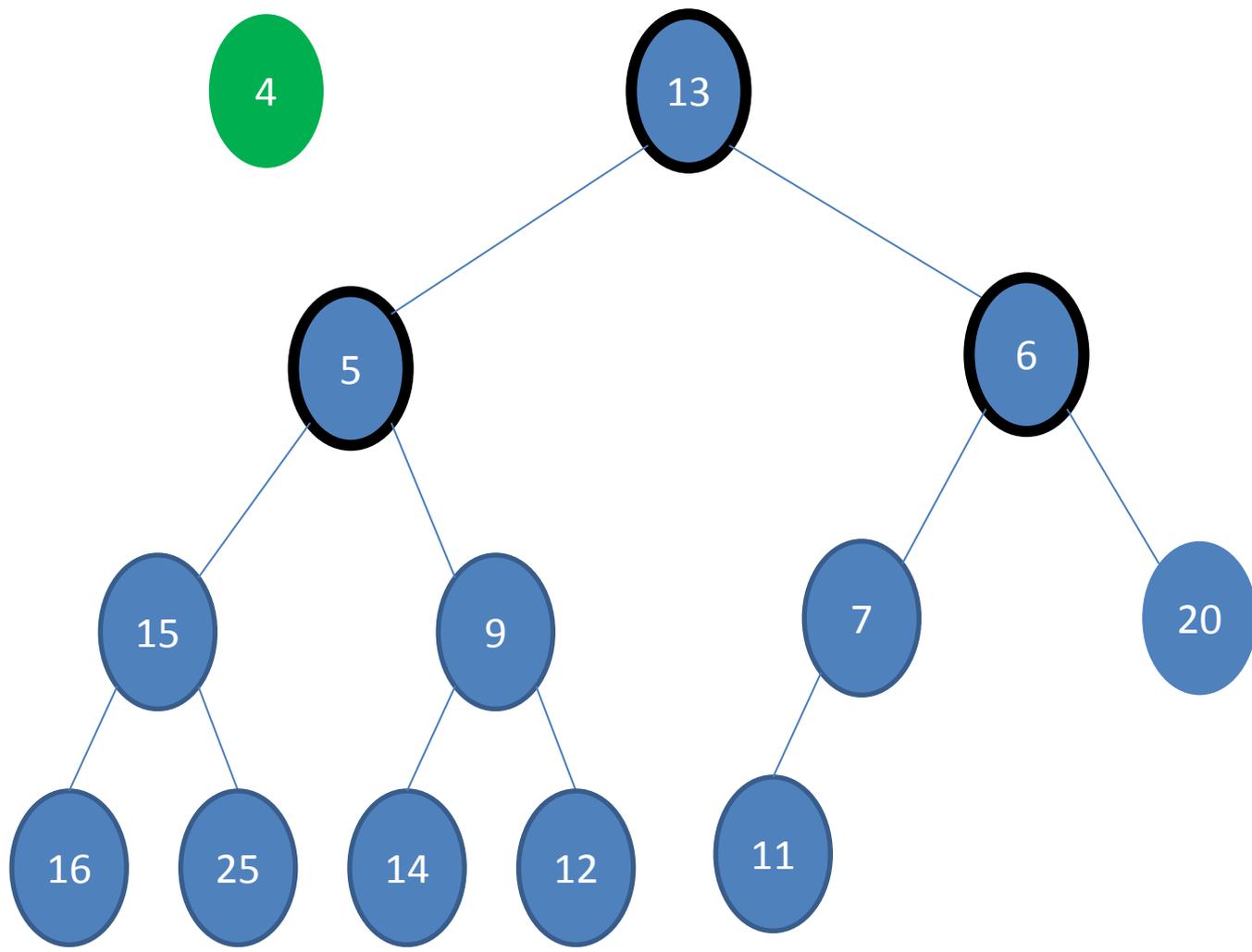
heap



Compare current node with its children

Example: removal from a heap

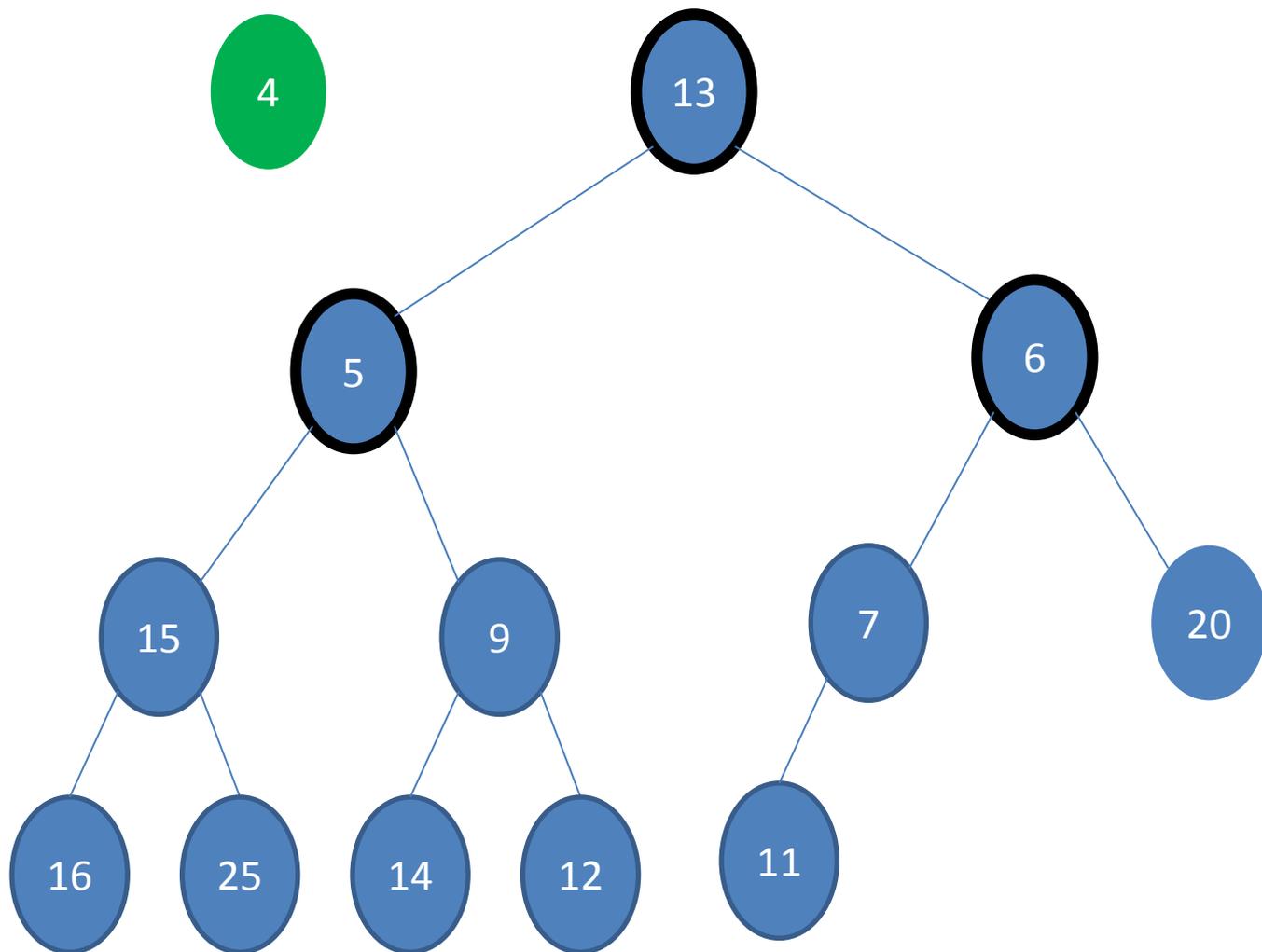
heap



Compare current node with its children

Example: removal from a heap

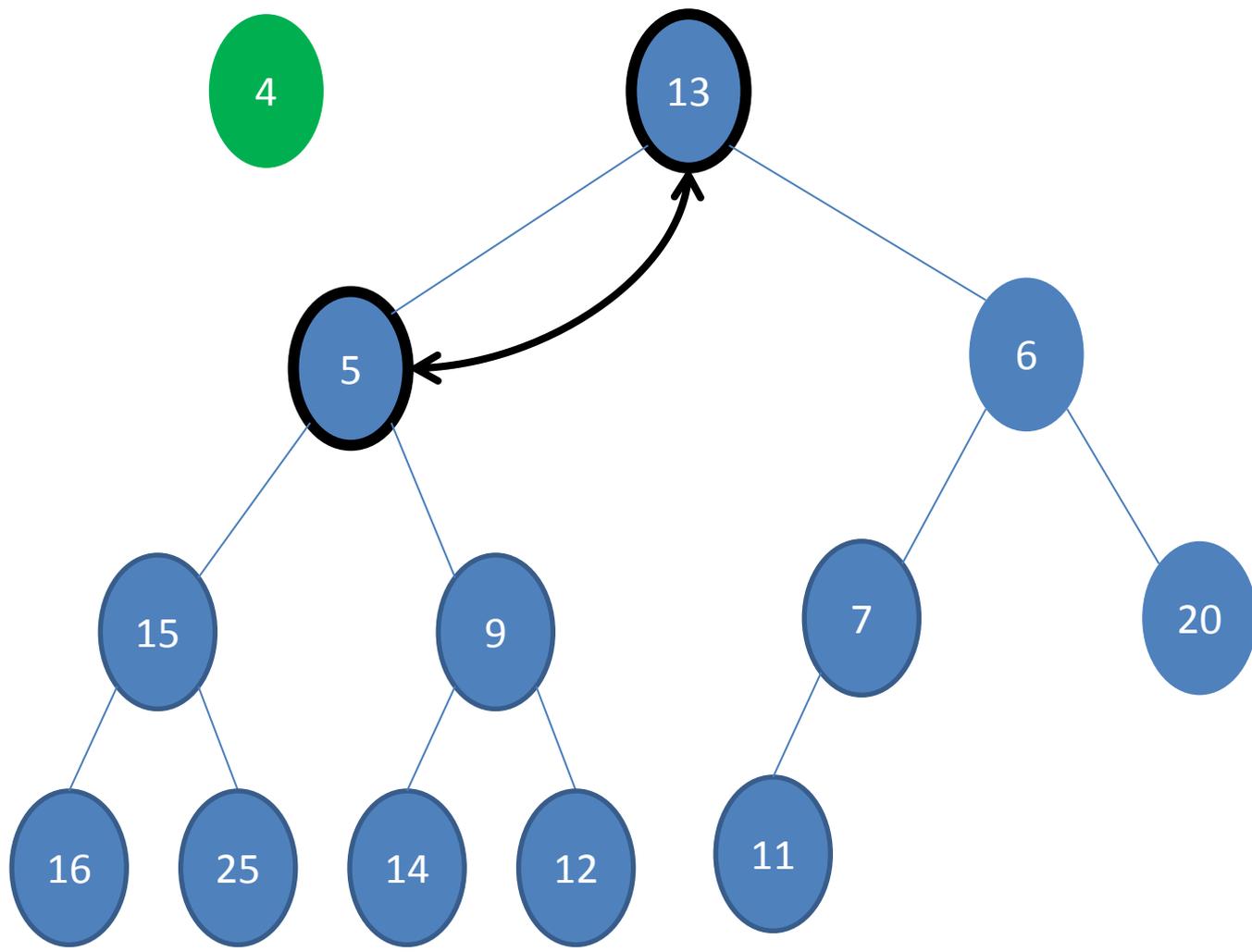
heap



If greater then swap with smallest child

Example: removal from a heap

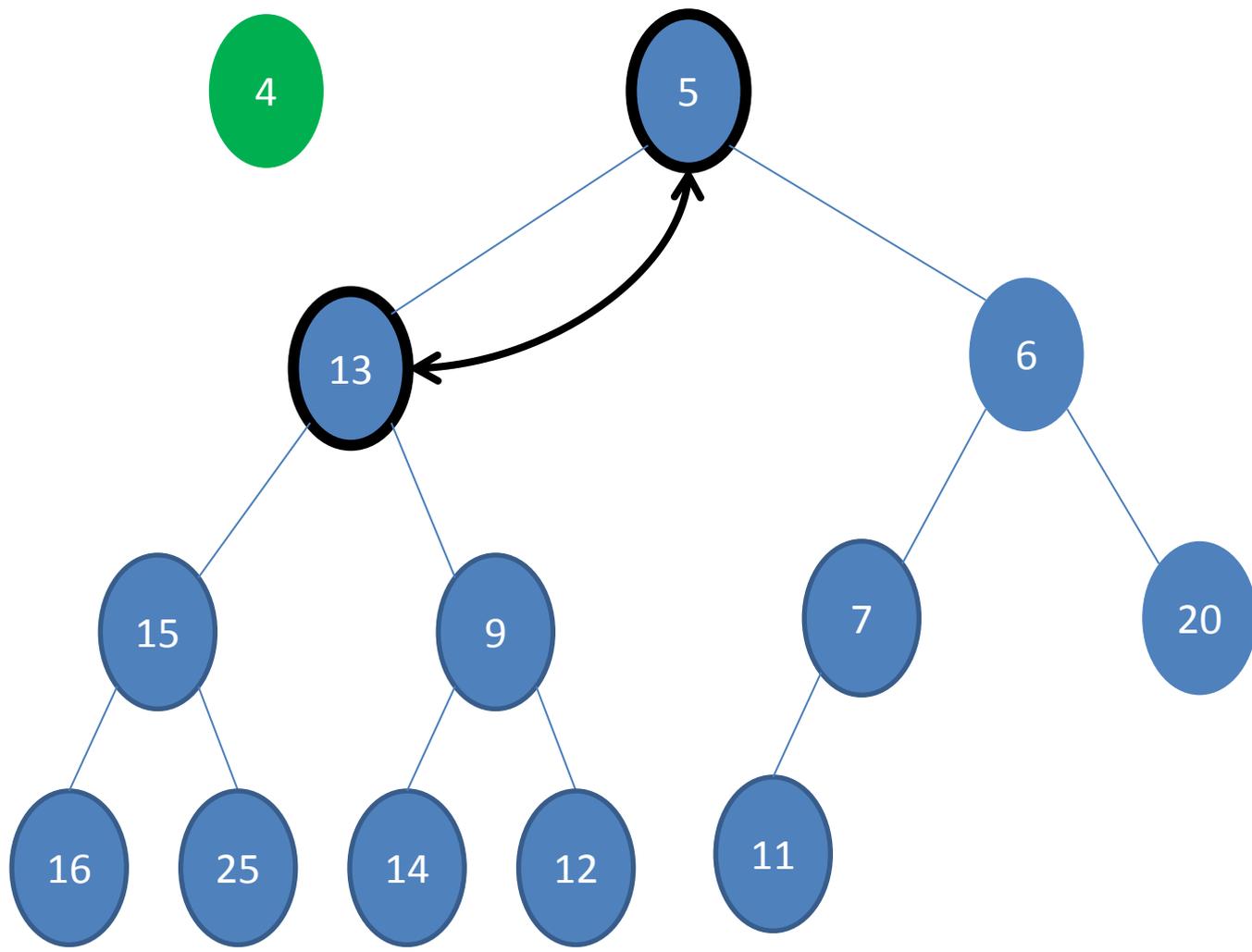
heap



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Example: removal from a heap

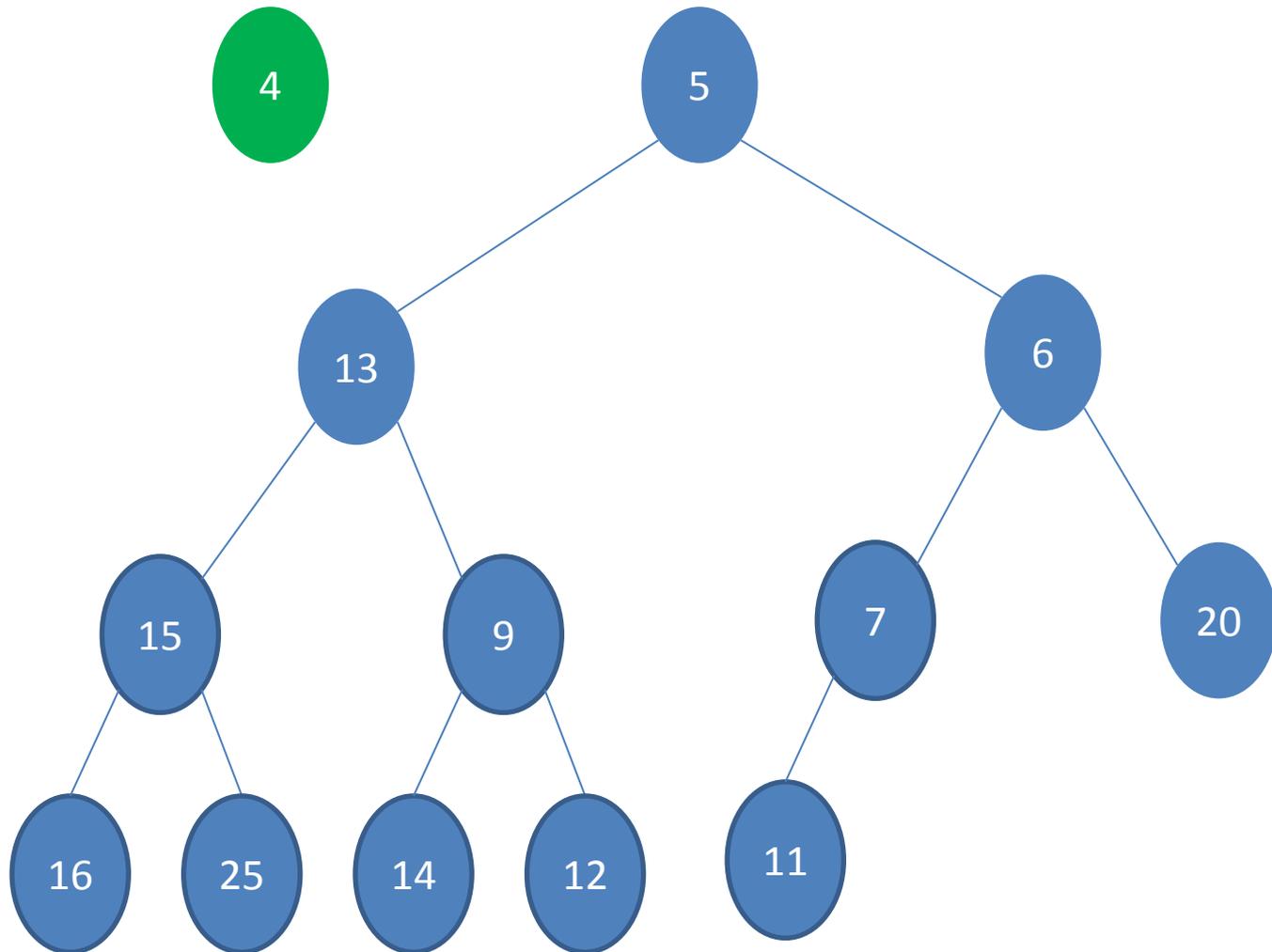
heap



If greater then swap with smallest child

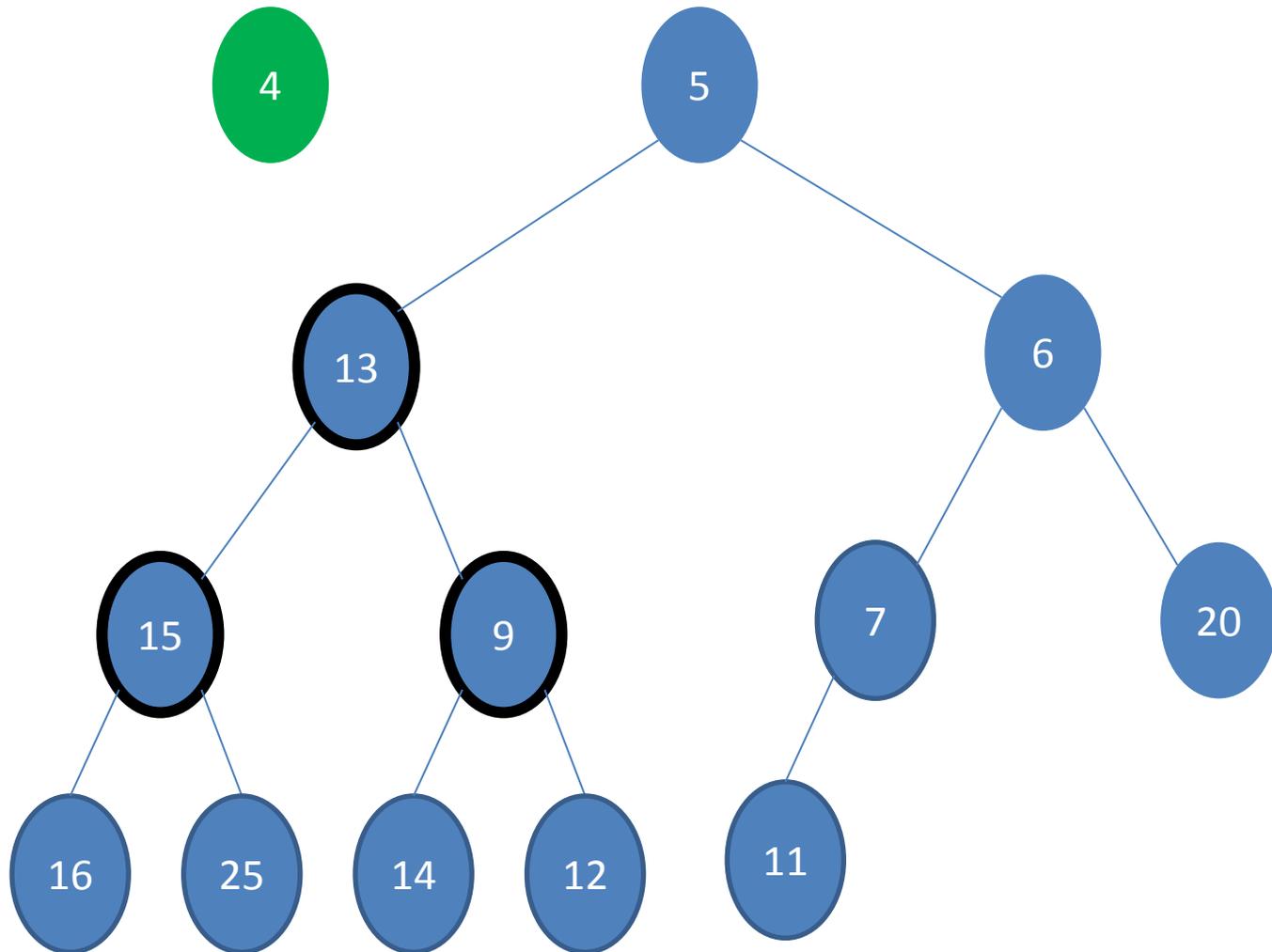
Example: removal from a heap

heap



Example: removal from a heap

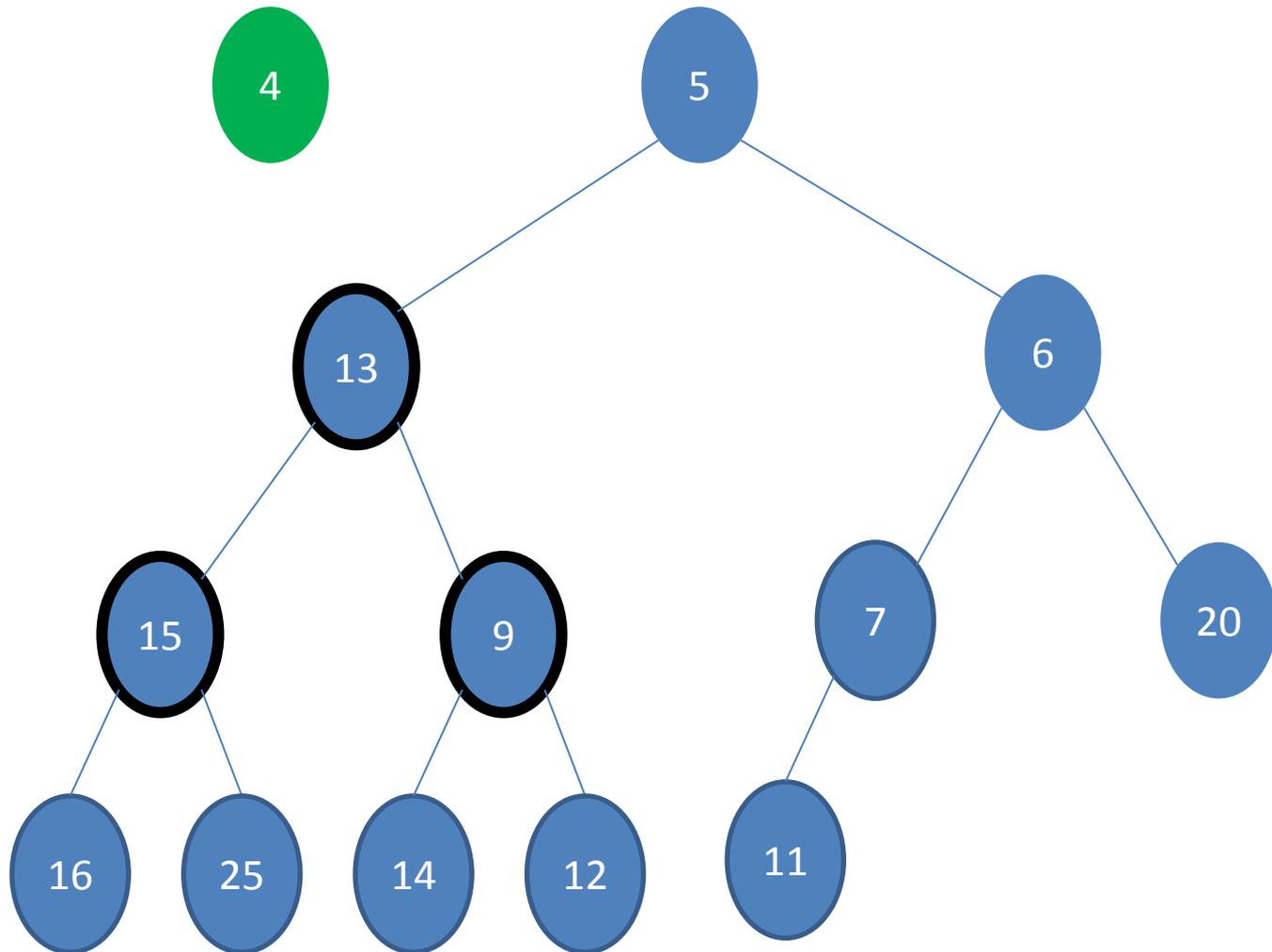
heap



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Example: removal from a heap

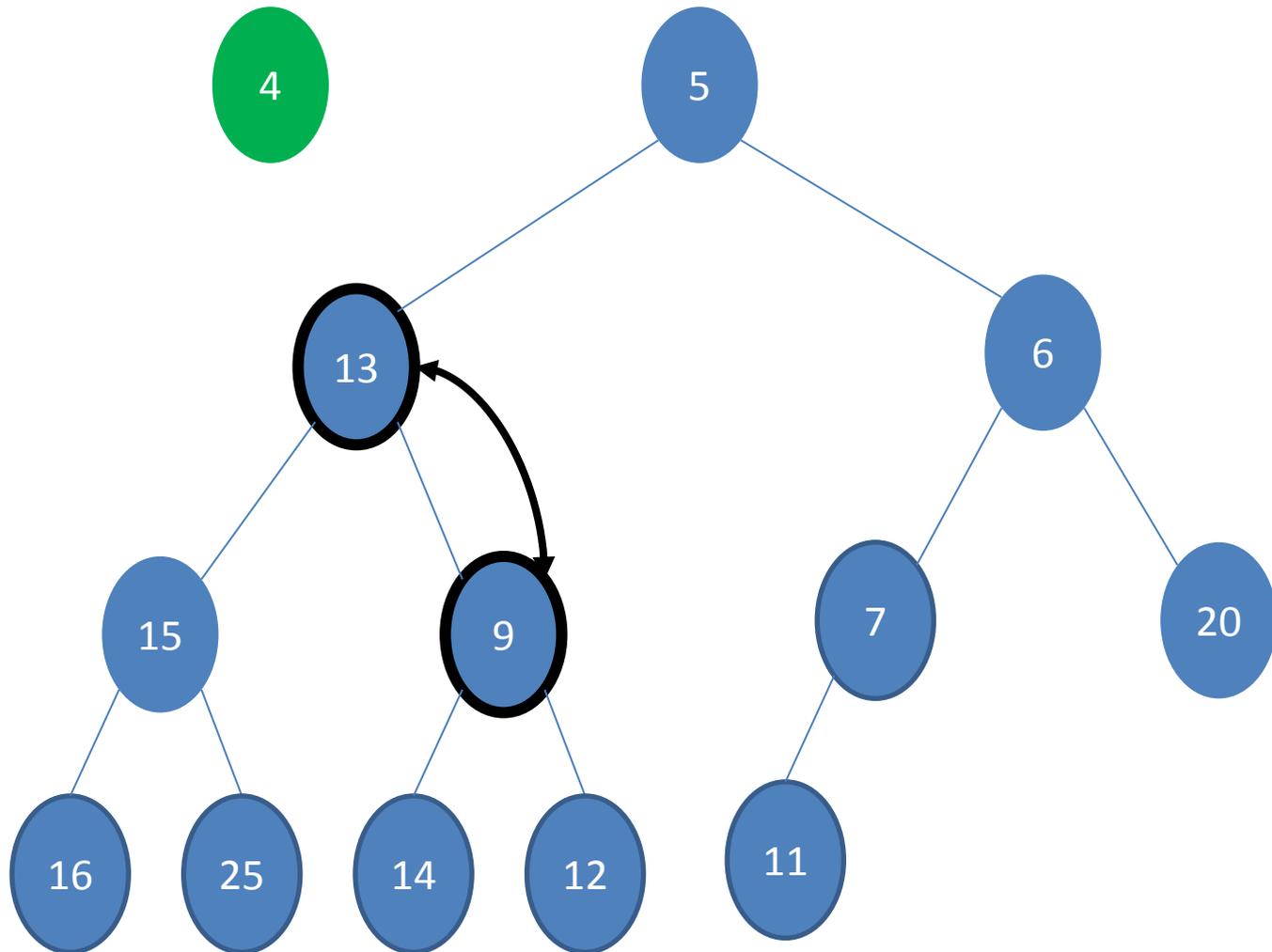
heap



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Example: removal from a heap

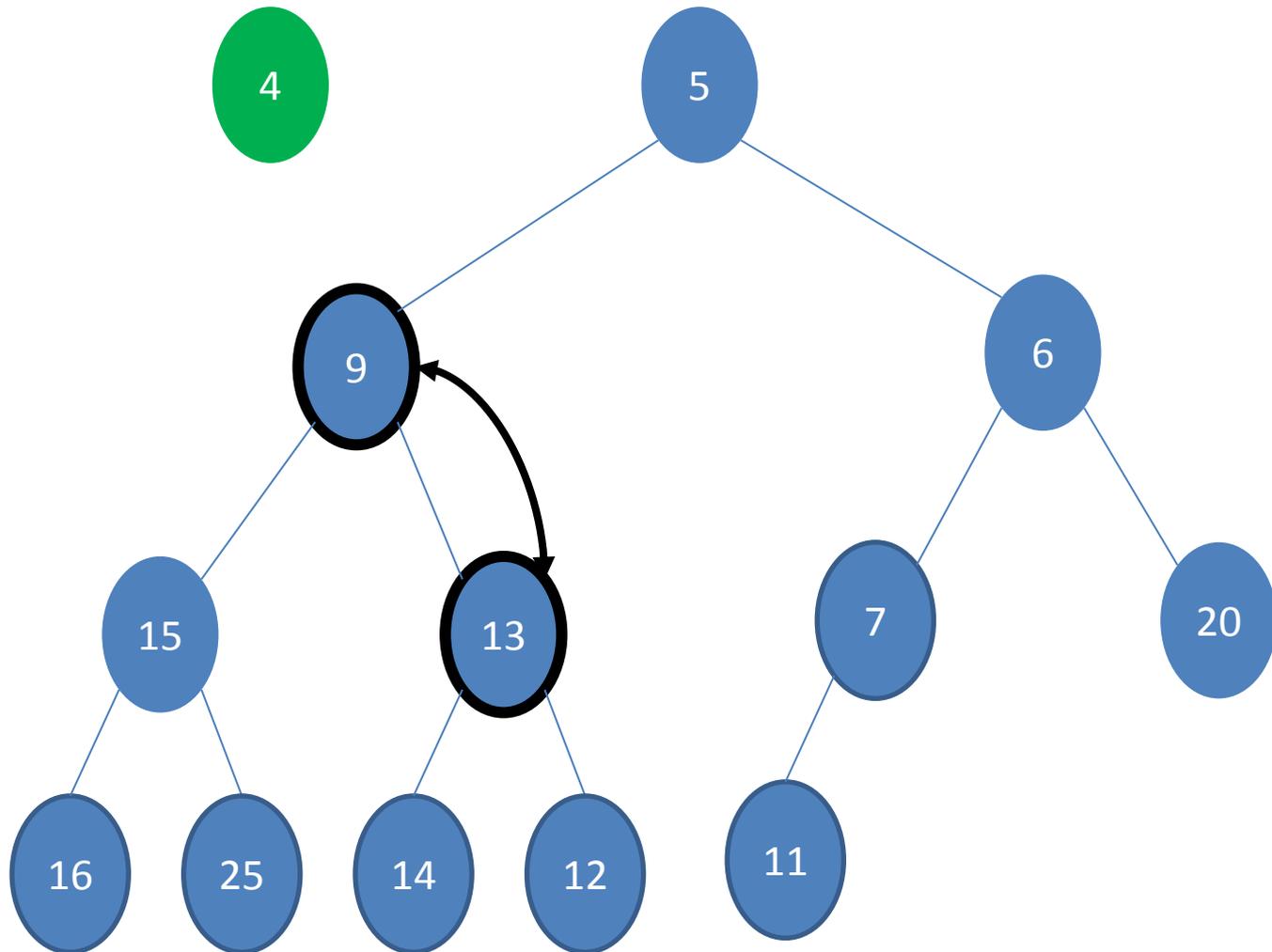
heap



If greater then swap with smallest child

Example: removal from a heap

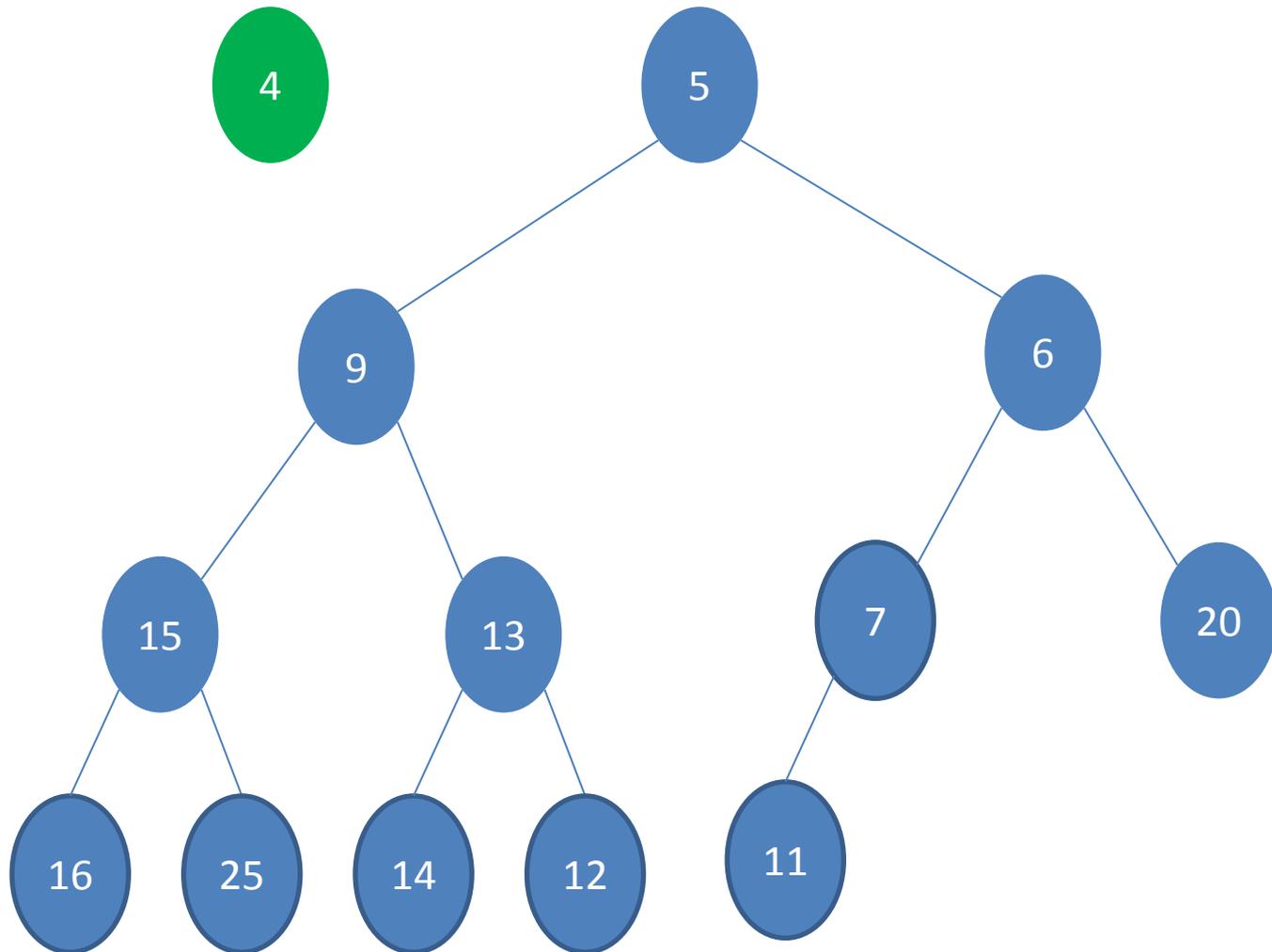
heap



If greater then swap with smallest child

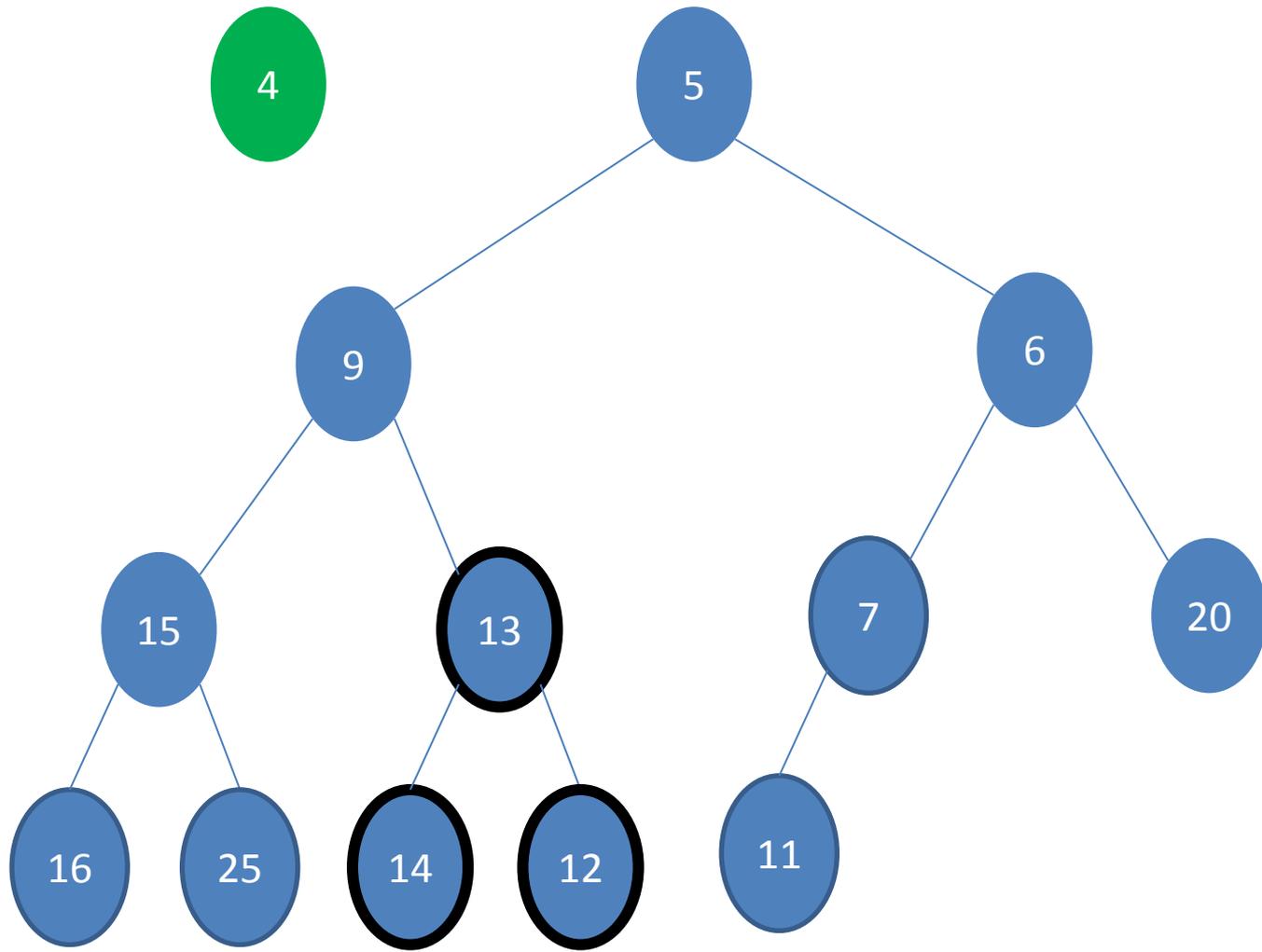
Example: removal from a heap

heap



Example: removal from a heap

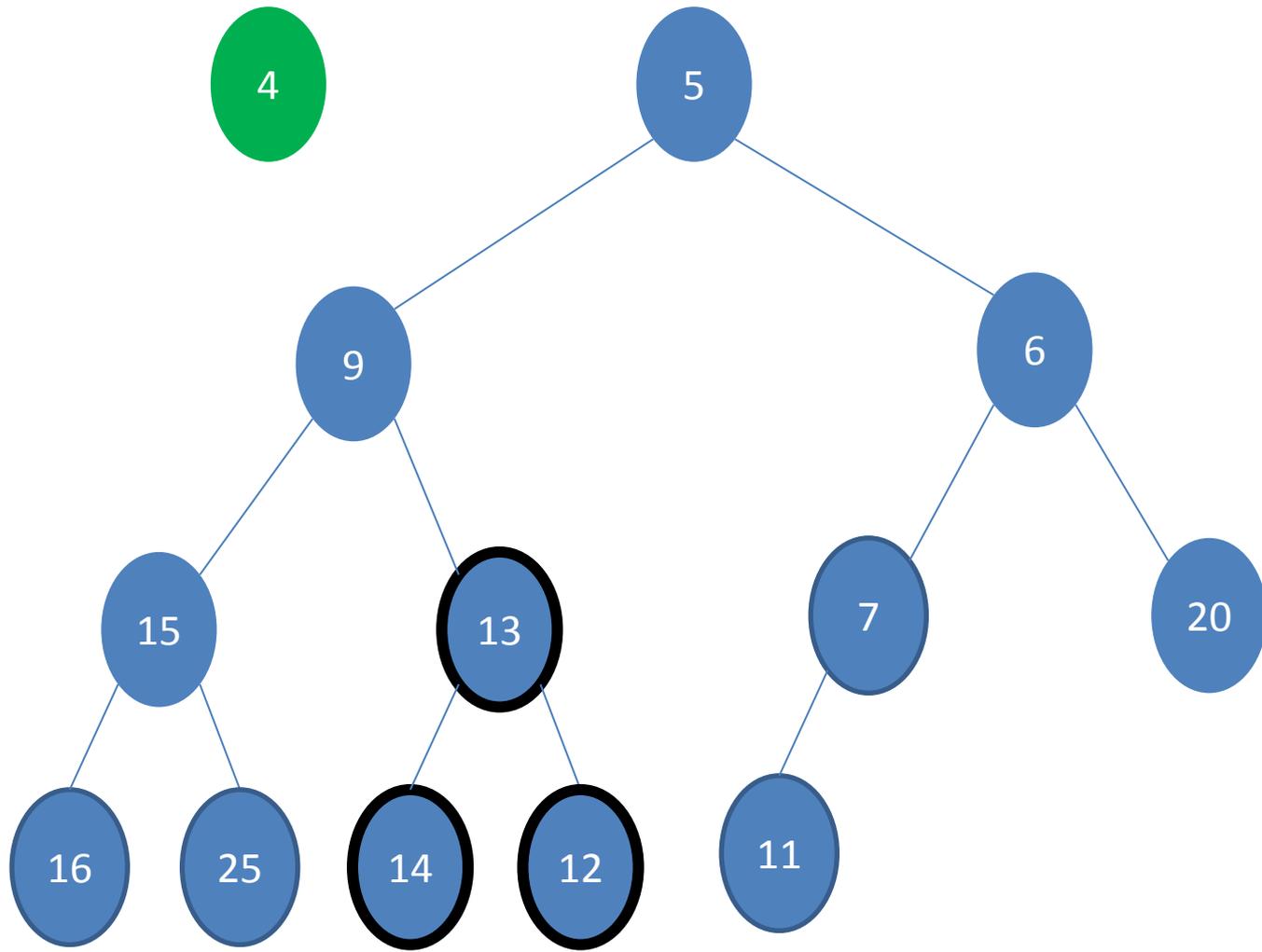
heap



Compare current node with its children

Example: removal from a heap

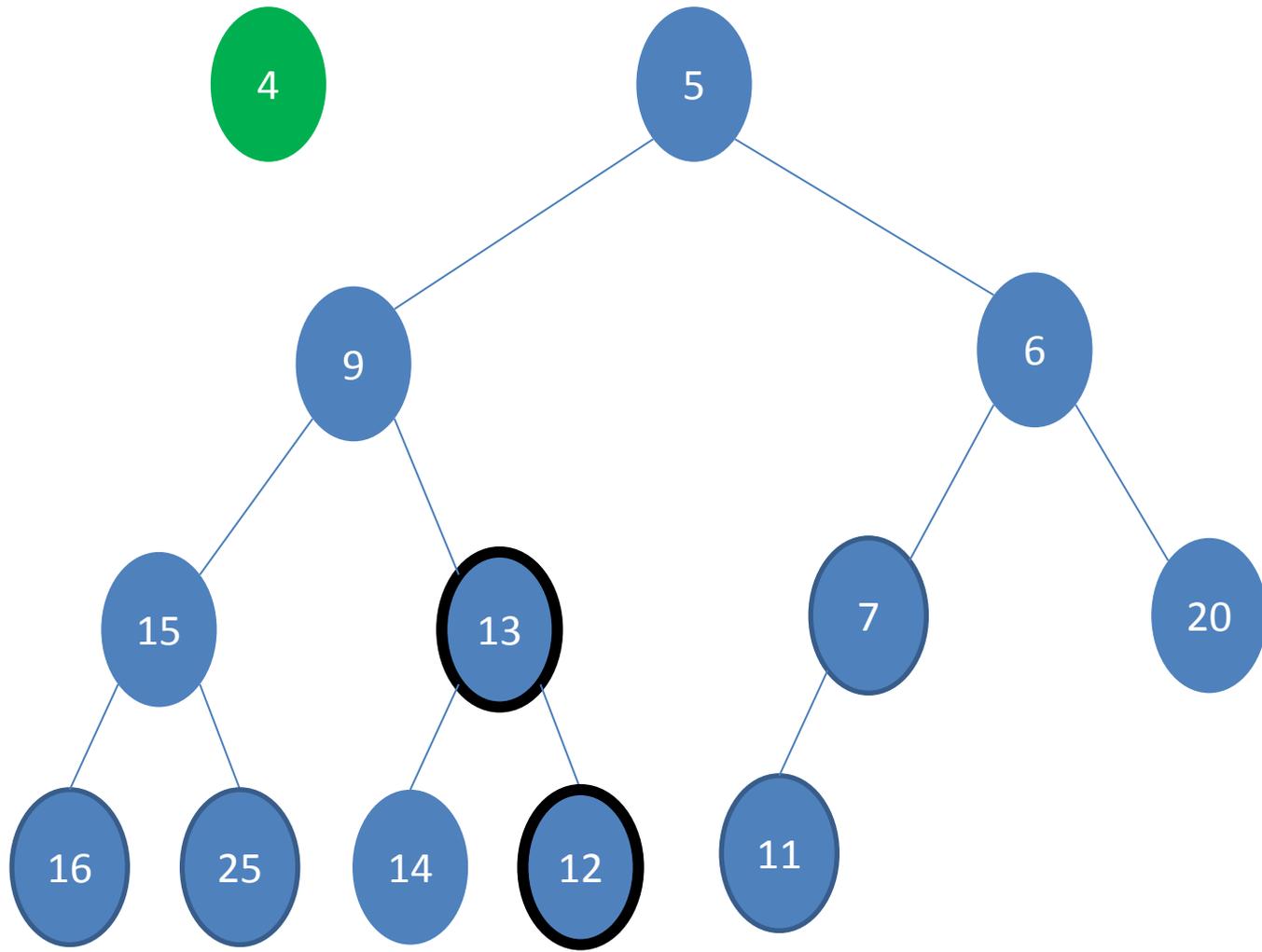
heap



If greater then swap with smallest child

Example: removal from a heap

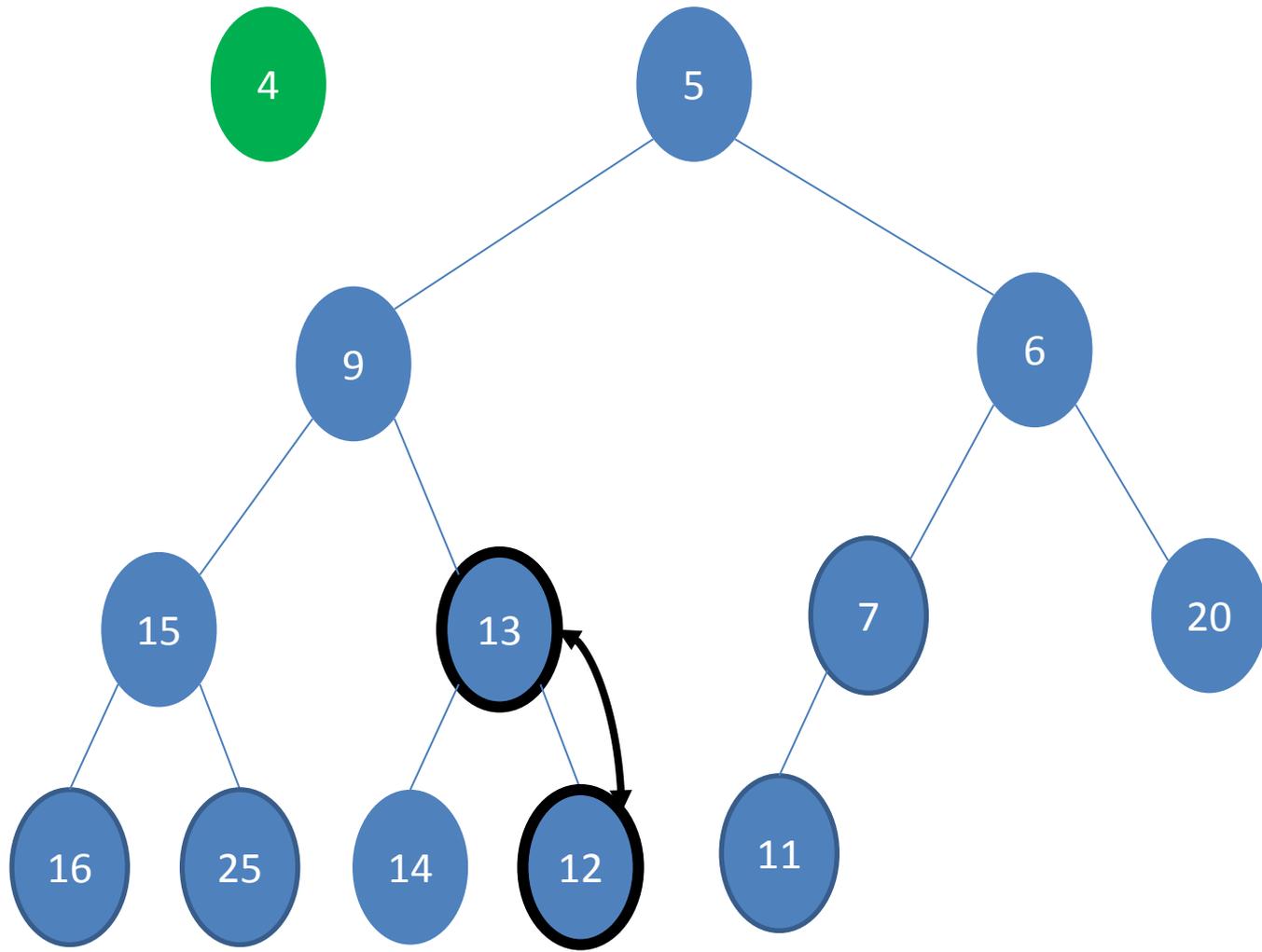
heap



If greater then swap with smallest child

Example: removal from a heap

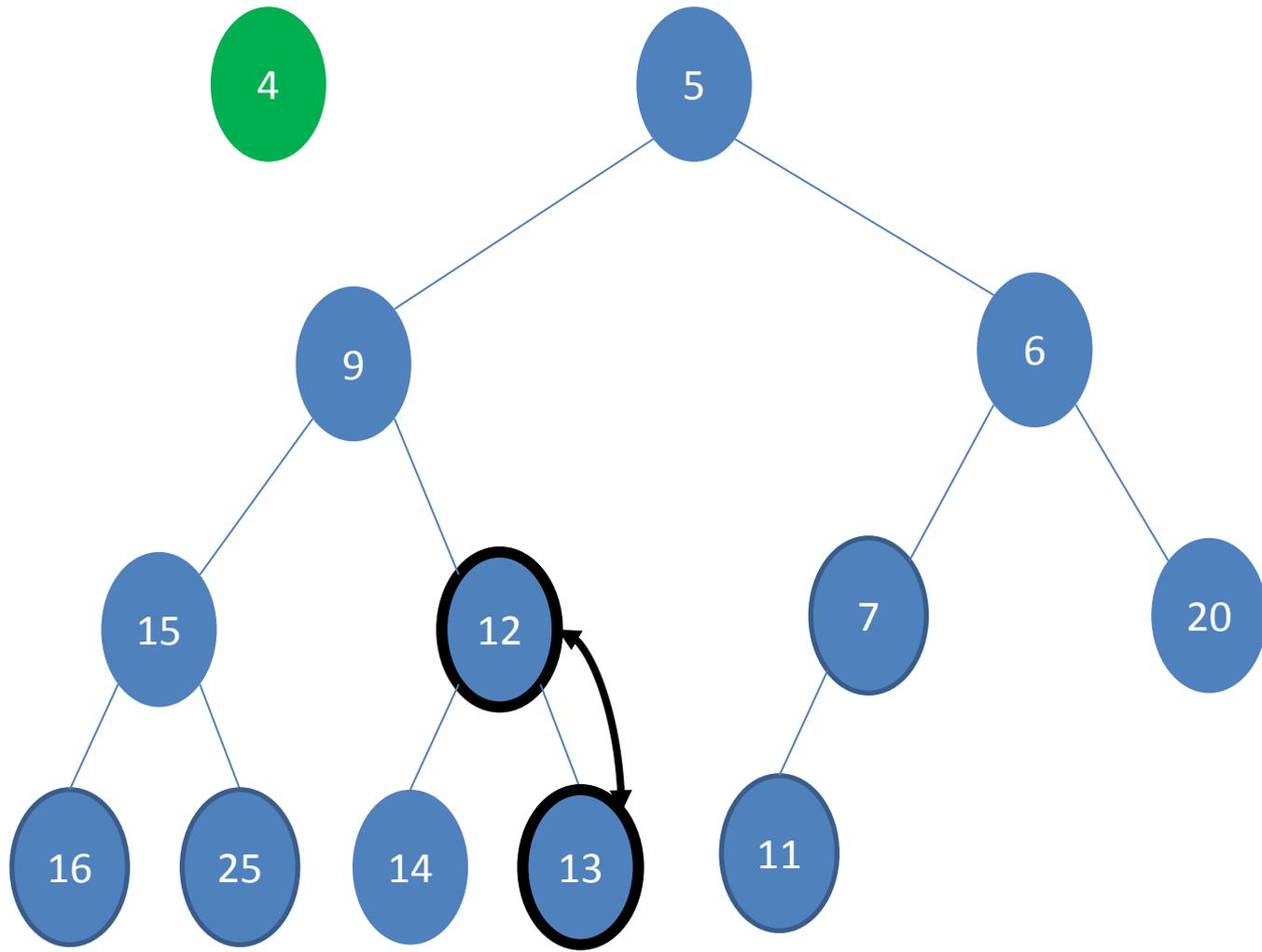
heap



If greater then swap with smallest child

Example: removal from a heap

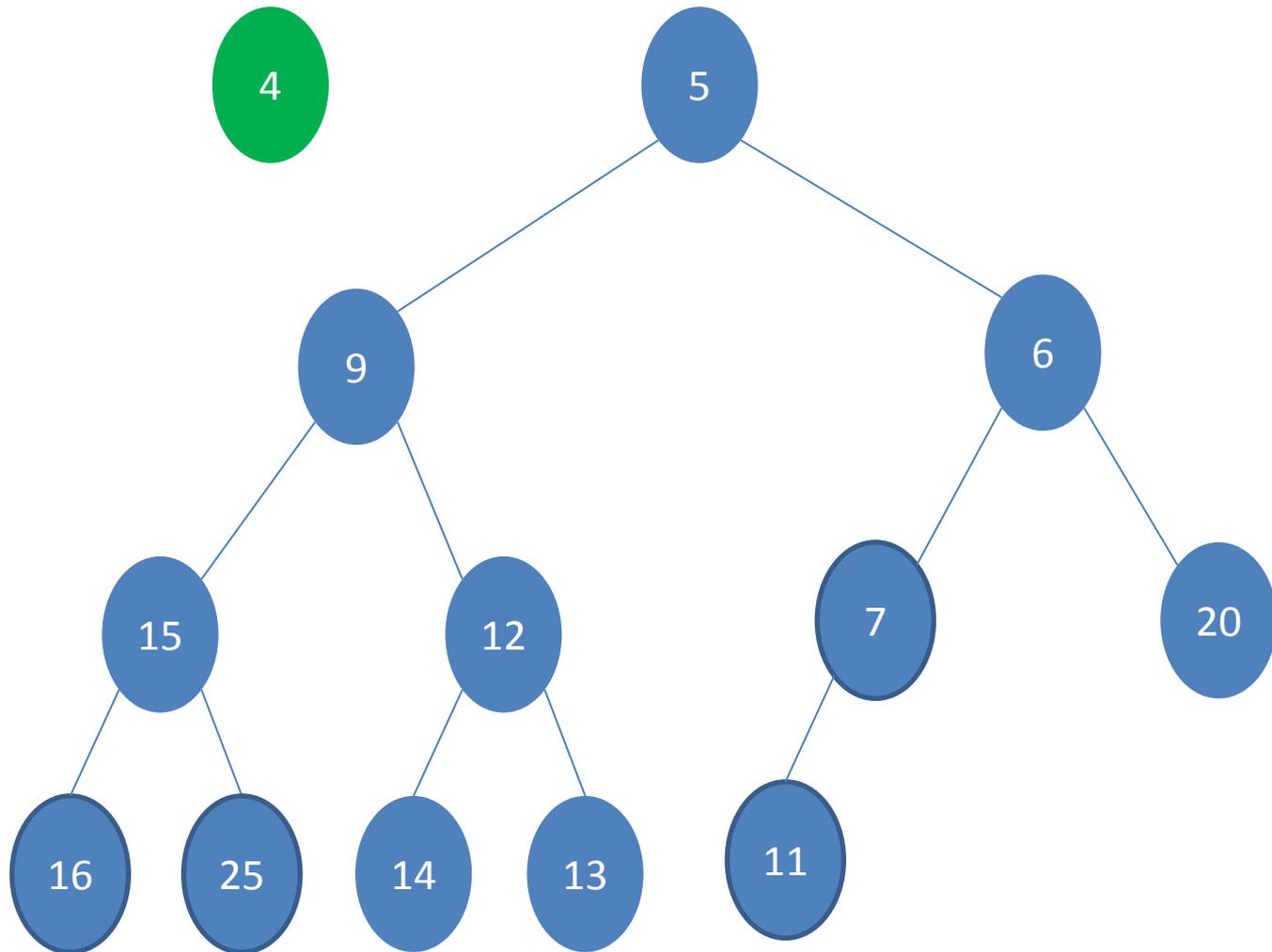
heap



If greater then swap with smallest child

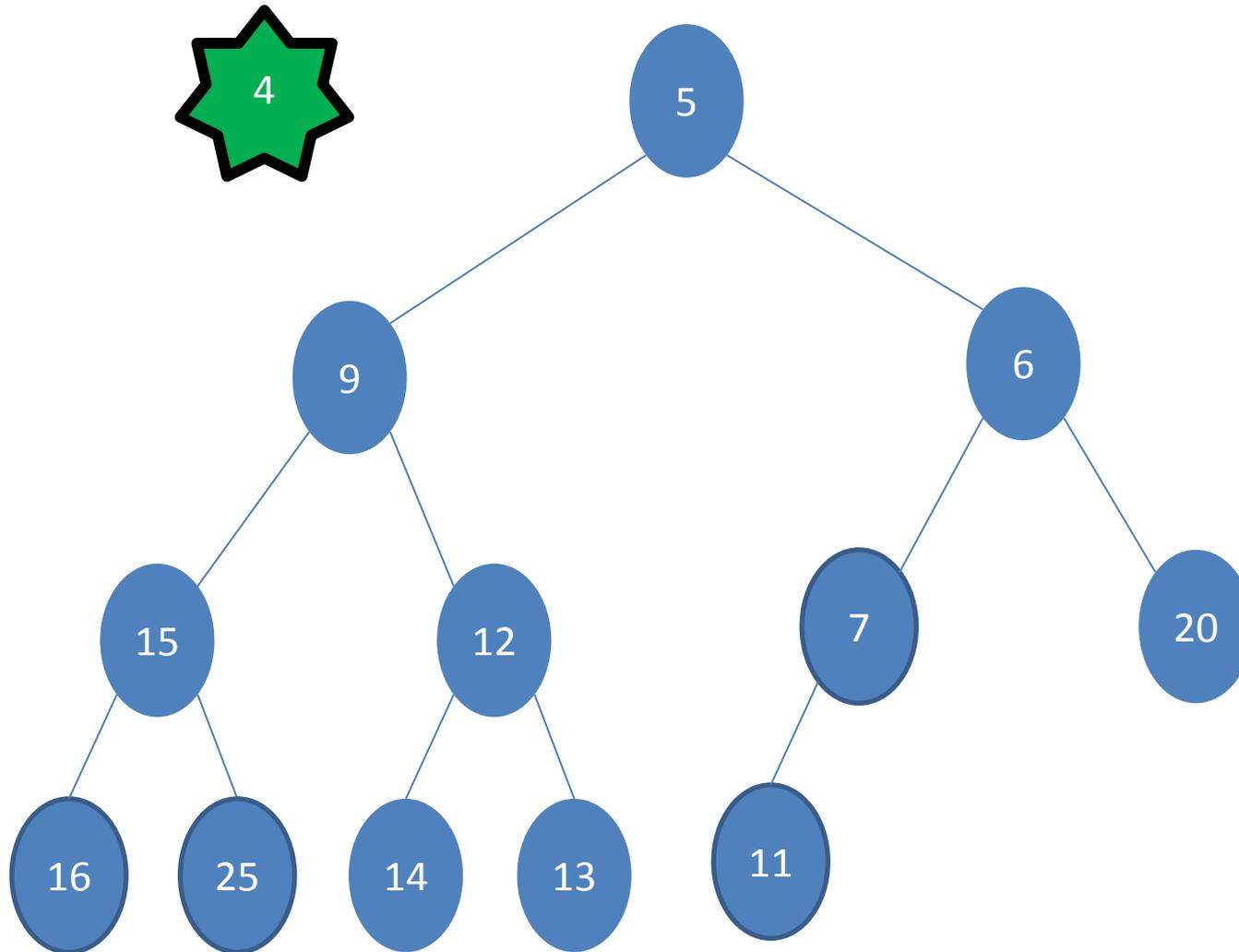
Example: removal from a heap

heap



Example: removal from a heap

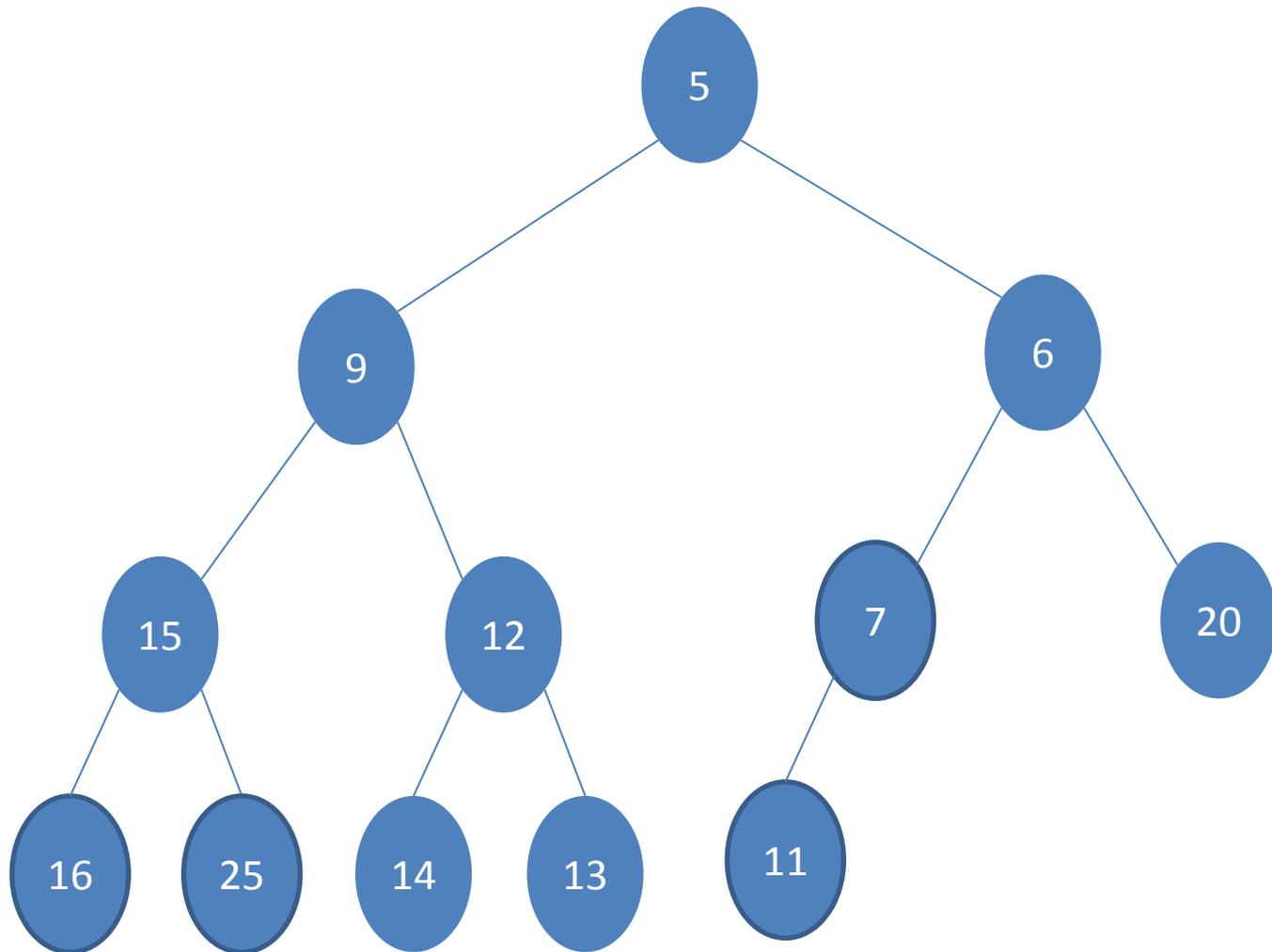
heap



Return result

Example: removal from a heap

heap



Done 😊

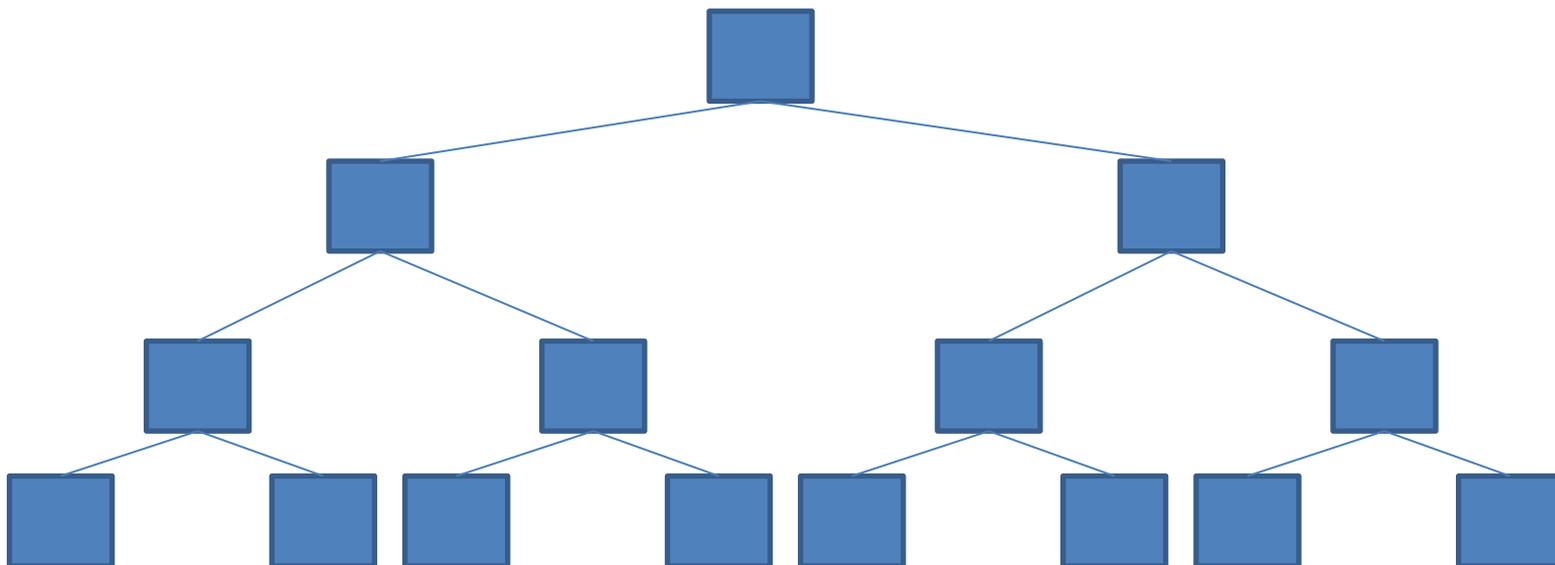
upheap bubbling: when we add to the heap

downheap bubbling: when we remove from the heap



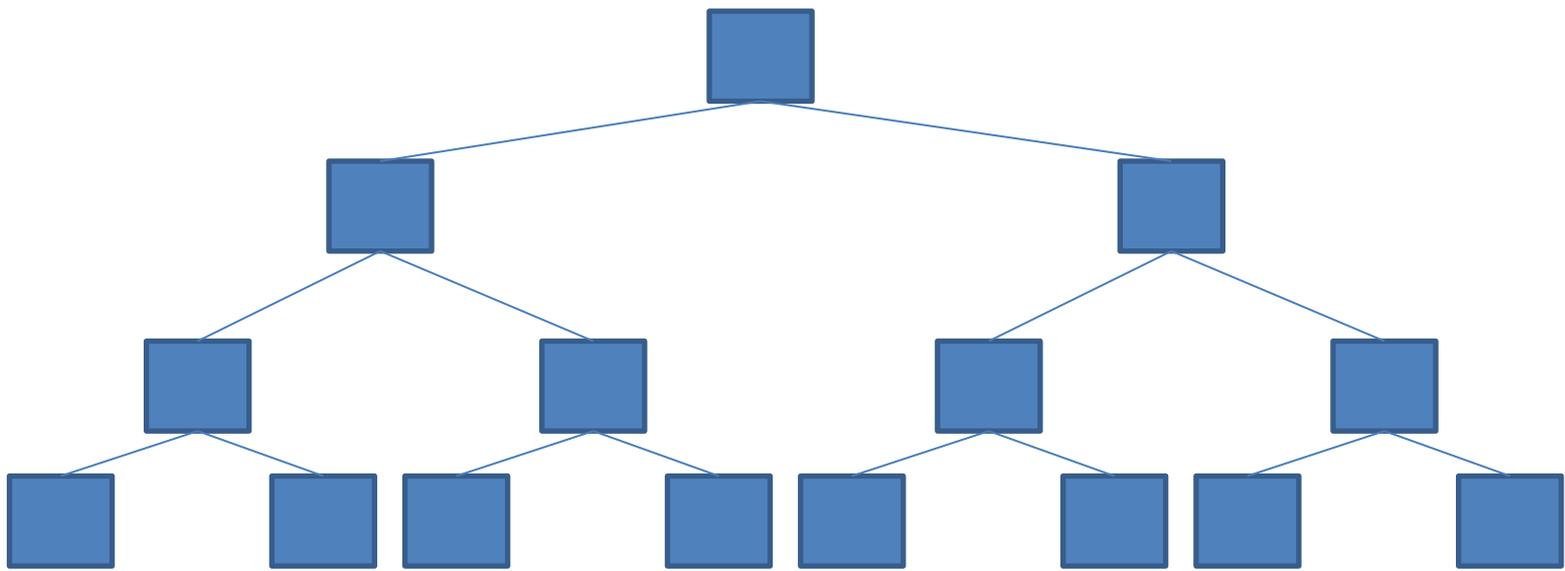
An implementation of a Heap data structure

An implementation of a Heap data structure



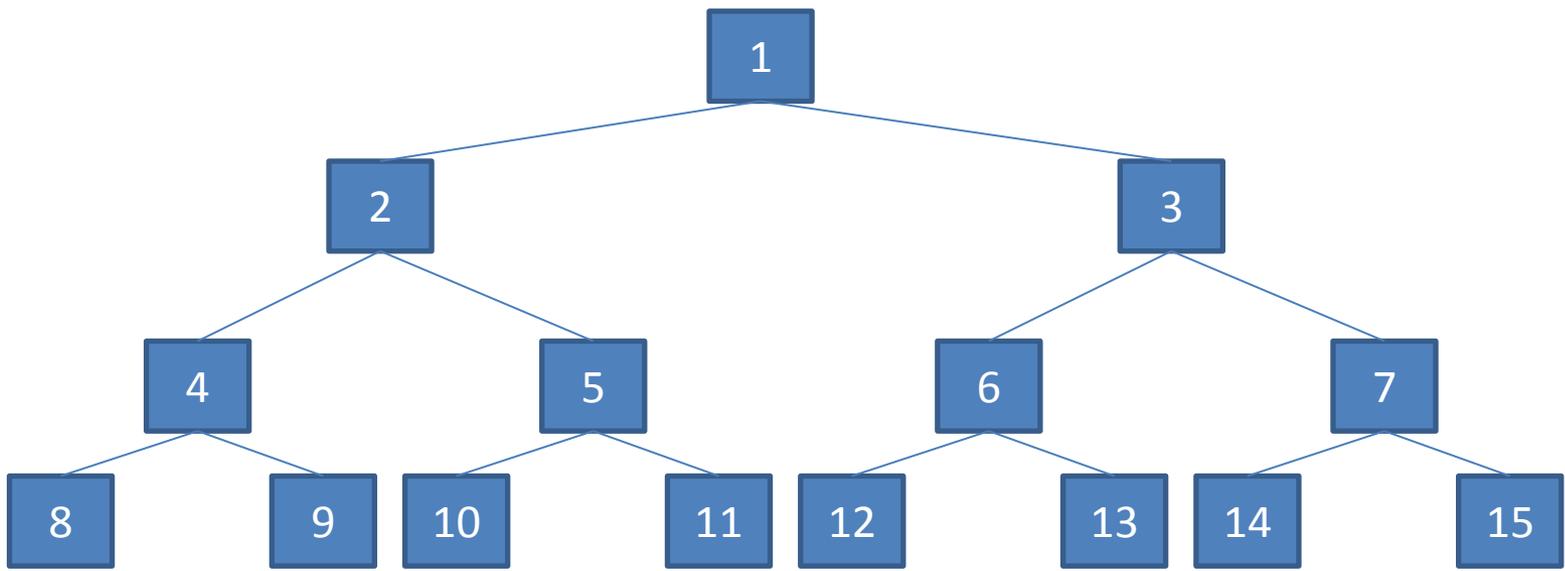
Number the vertices as follows

An implementation of a Heap data structure



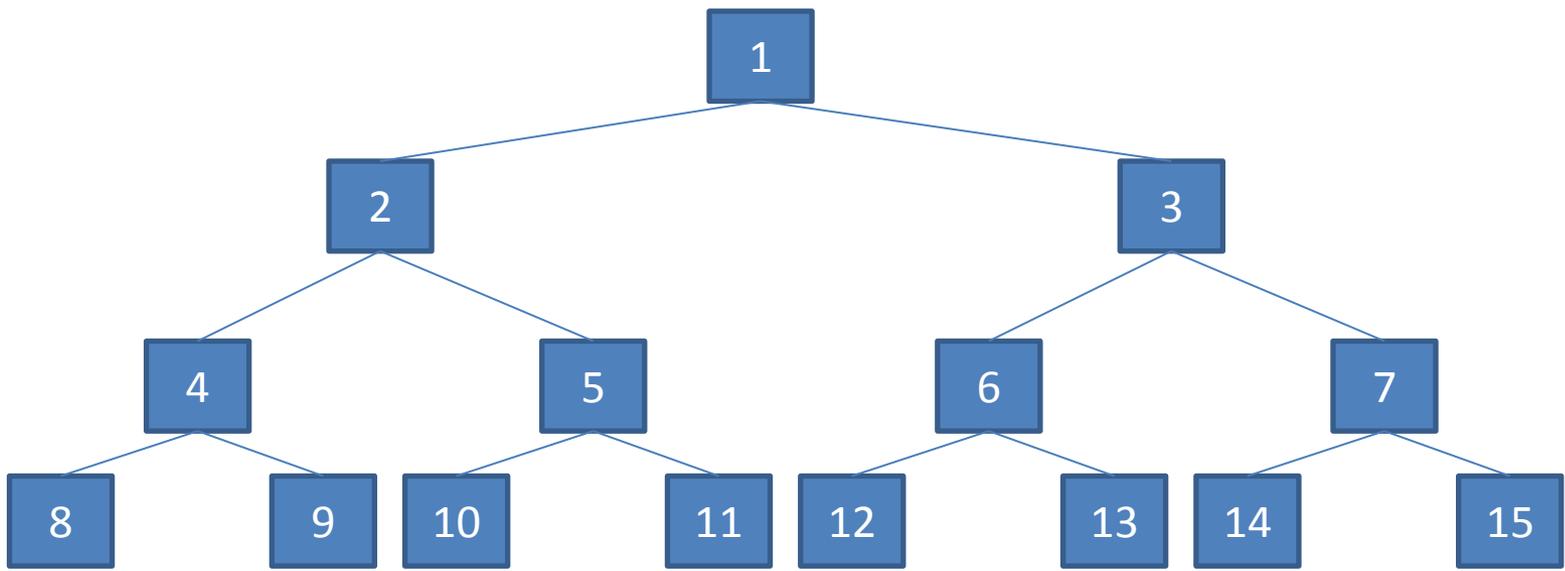
Number the vertices as follows

An implementation of a Heap data structure



Number the vertices as follows

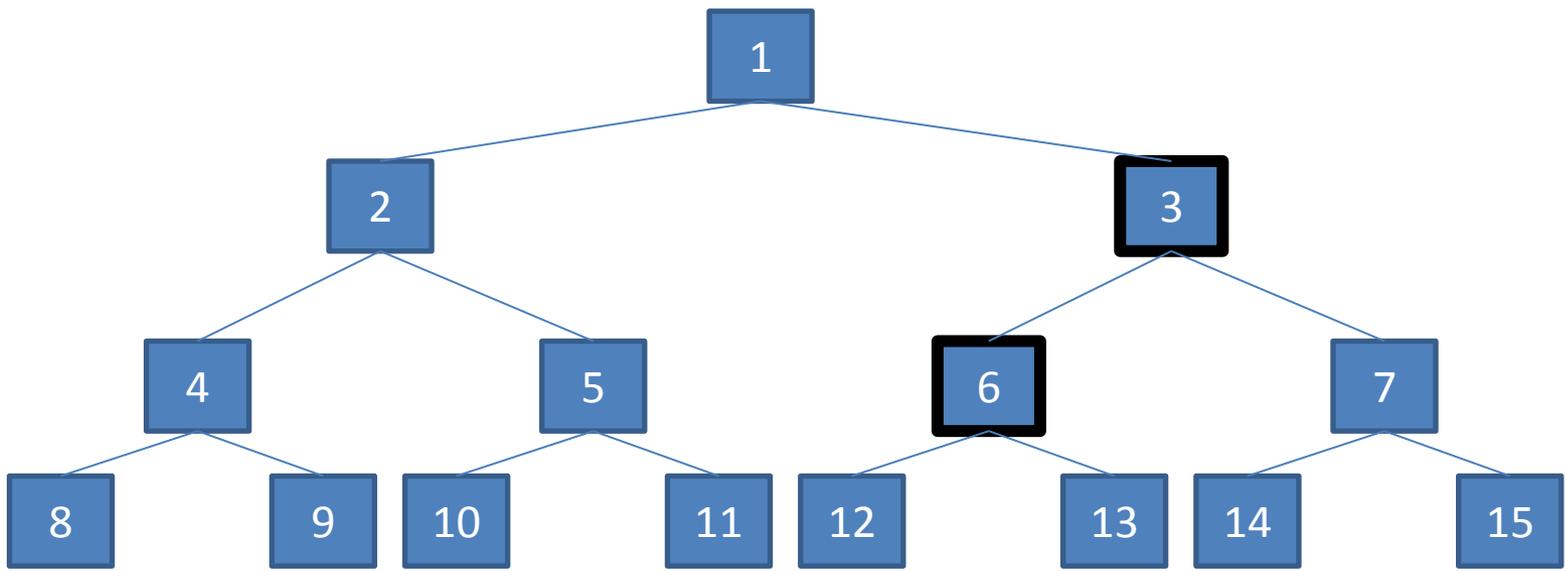
An implementation of a Heap data structure



Note: **parent** of node i is $i/2$

Number the vertices as follows

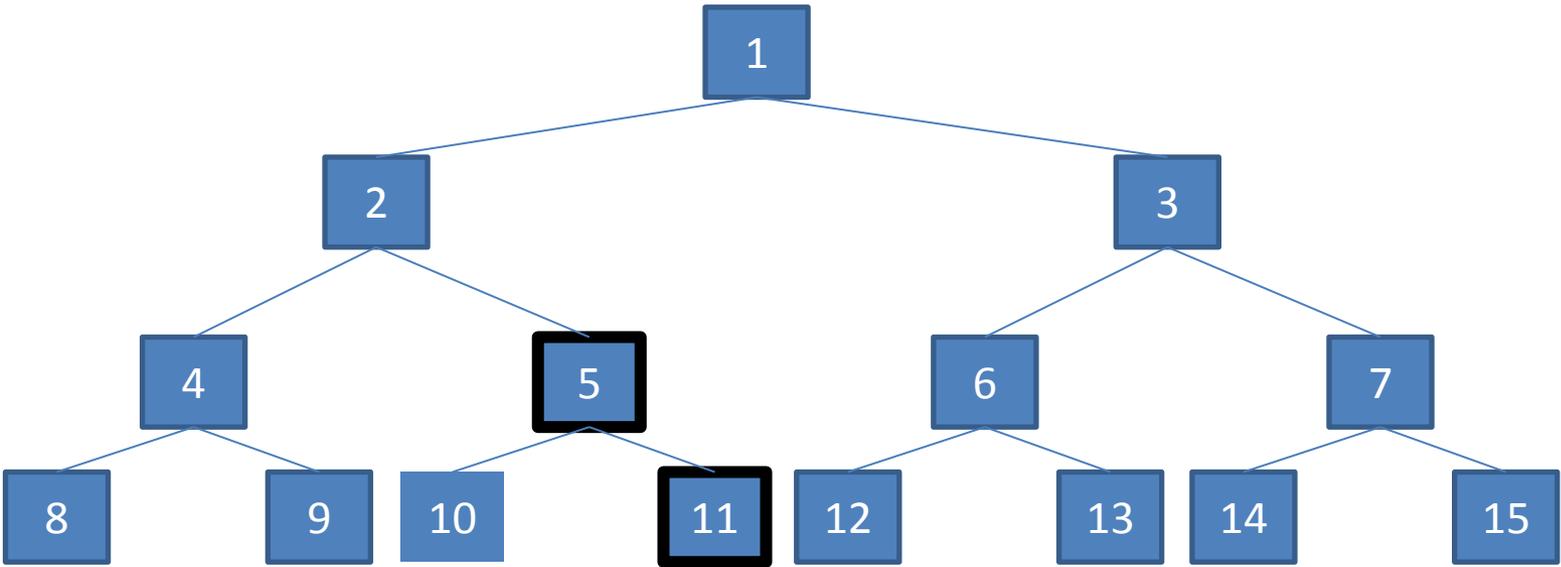
An implementation of a Heap data structure



Note: **parent** of node i is $i/2$

Number the vertices as follows

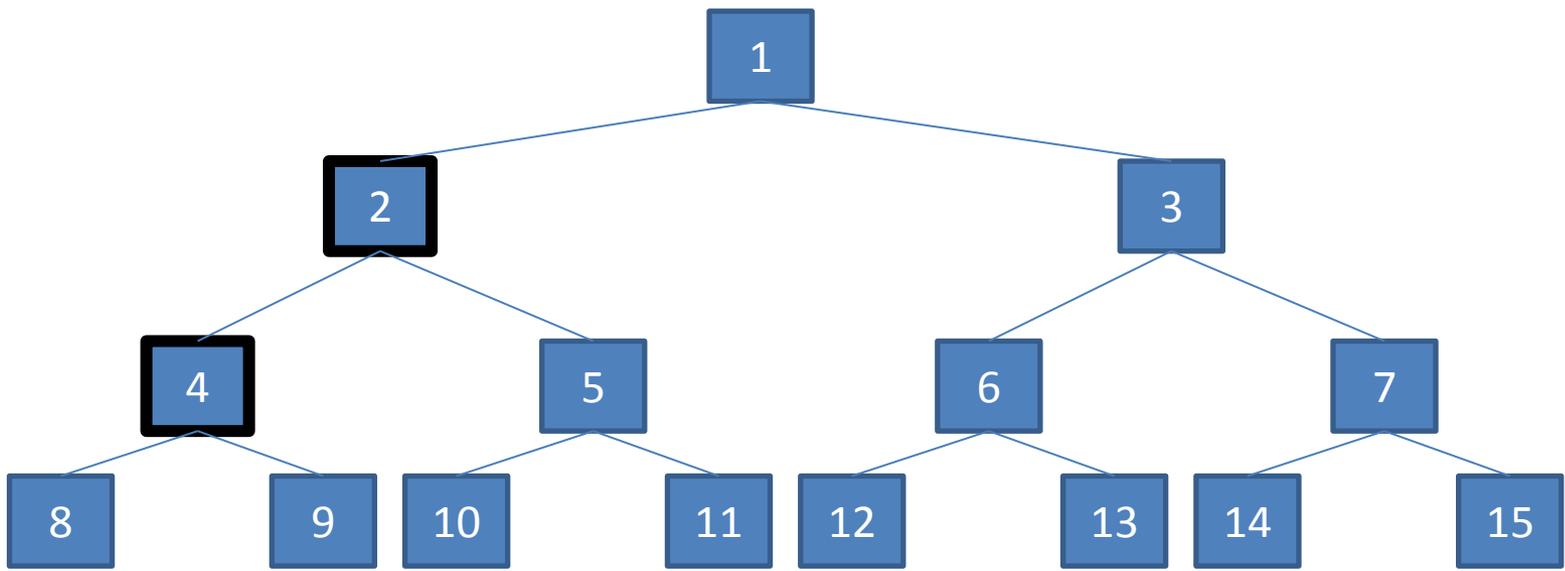
An implementation of a Heap data structure



Note: *parent* of node i is $i/2$

Number the vertices as follows

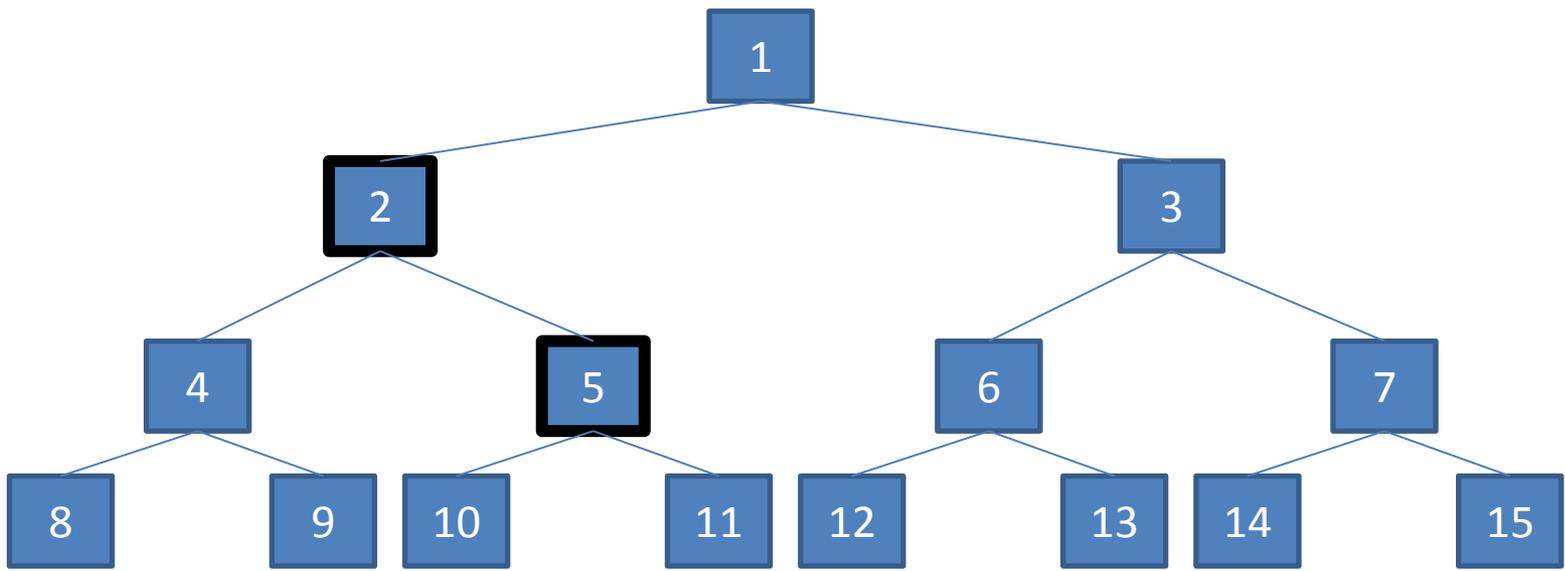
An implementation of a Heap data structure



Note: *left child* of i is $i \times 2$

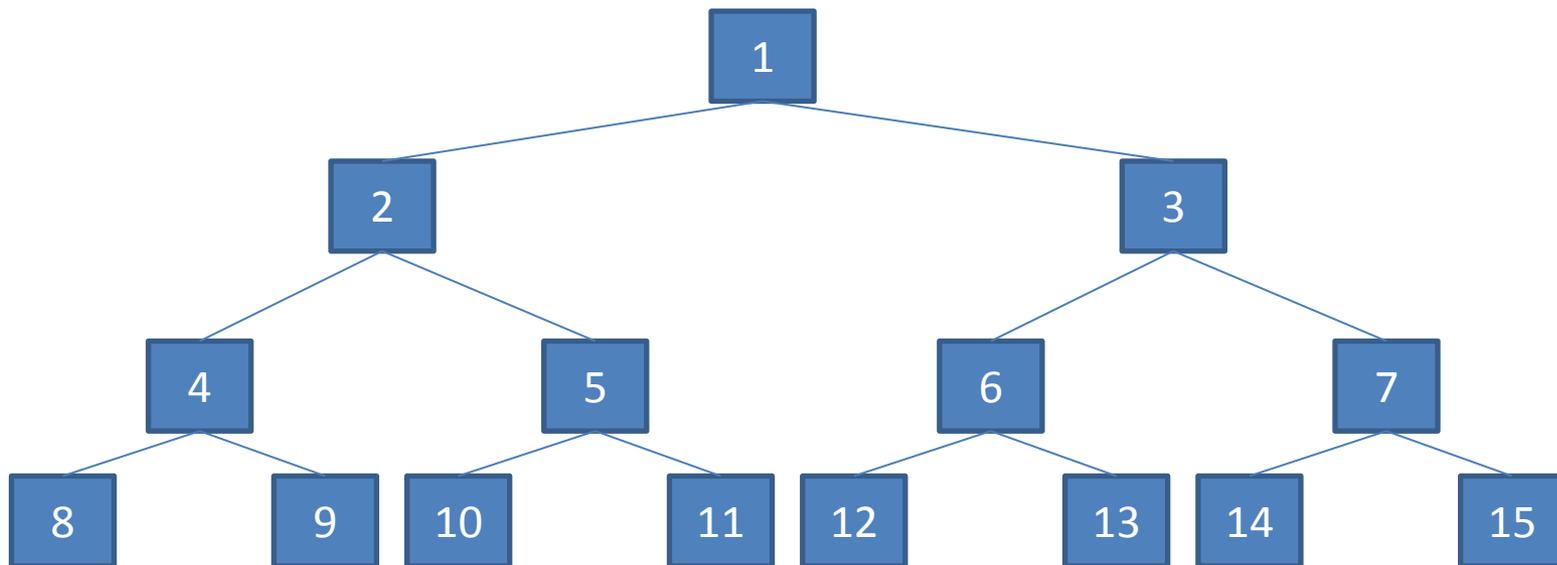
Number the vertices as follows

An implementation of a Heap data structure



Note: **right child** of i is $(i \times 2) + 1$

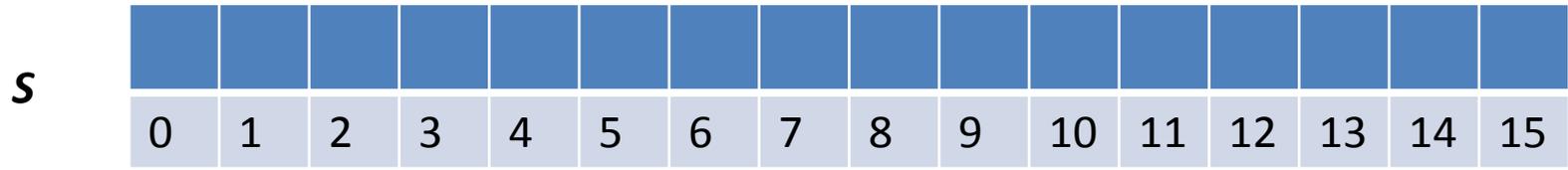
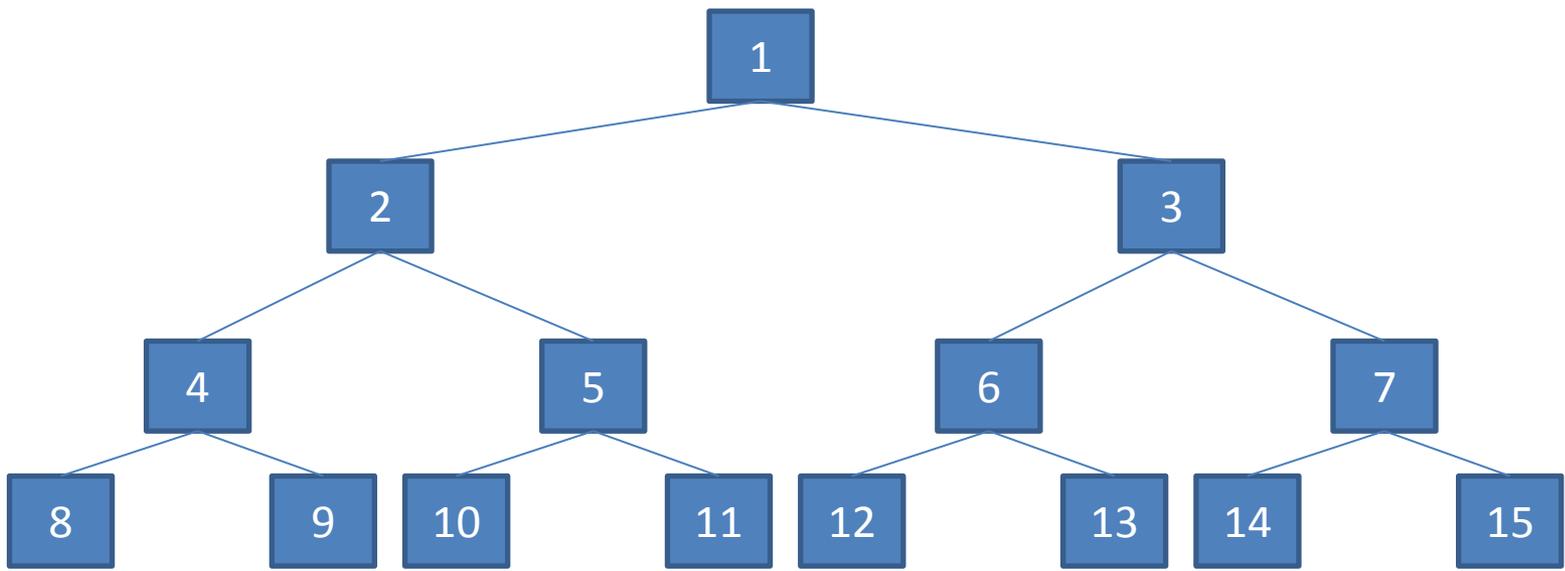
An implementation of a Heap data structure



Represent as a one *dimensional* array

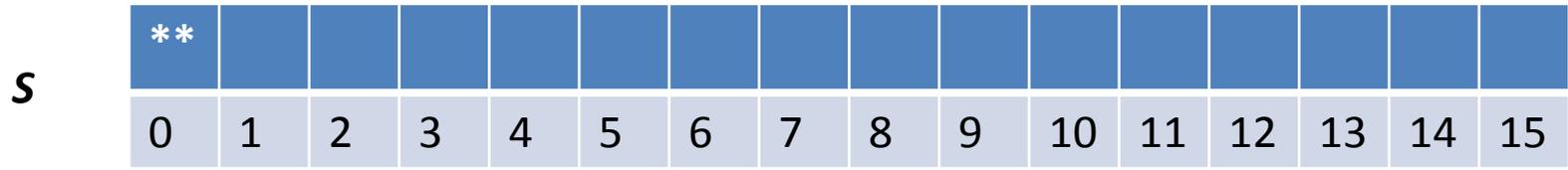
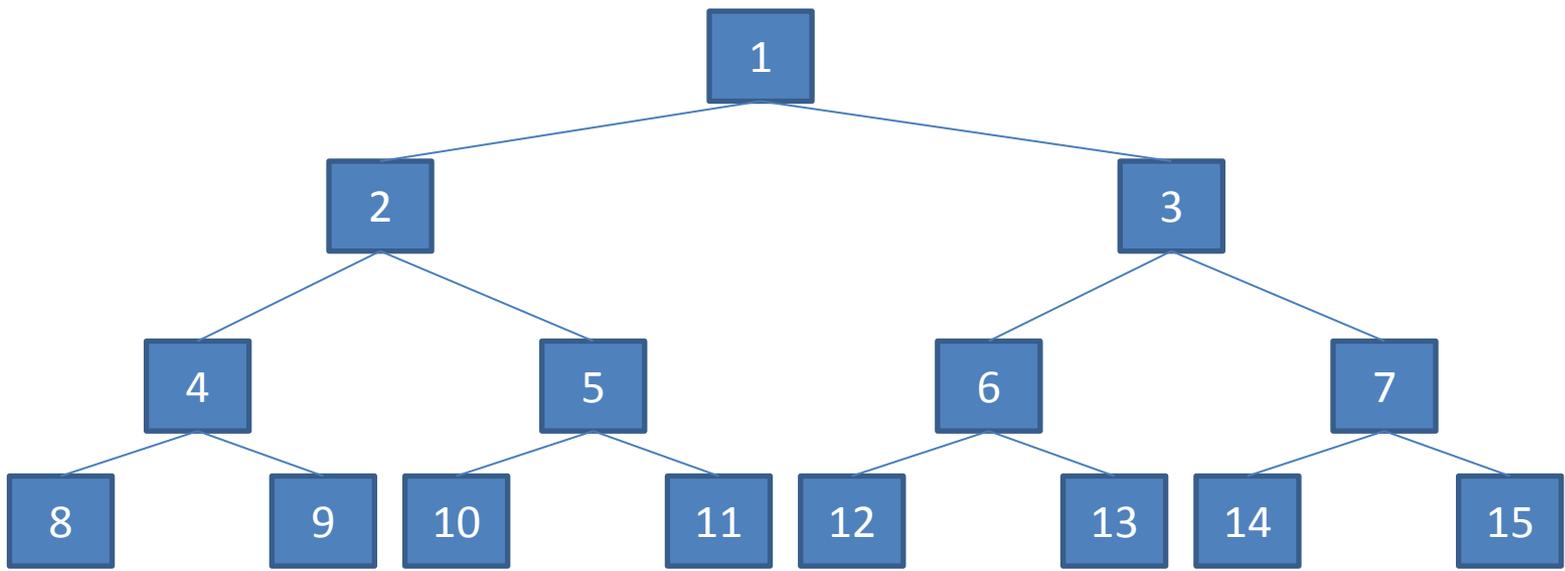
Represent as a one *dimensional* array **S**

An implementation of a Heap data structure



Represent as a one *dimensional* array S

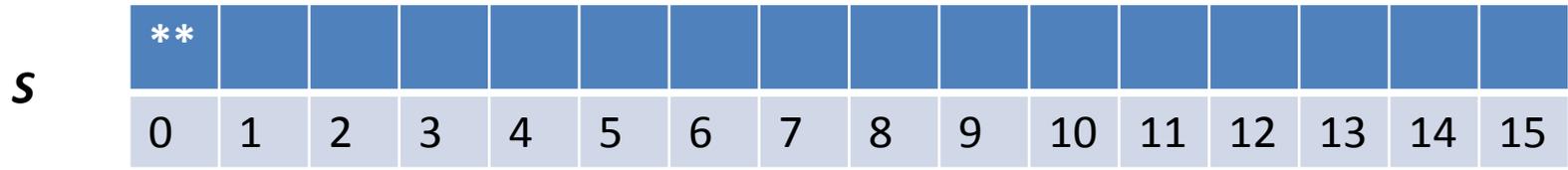
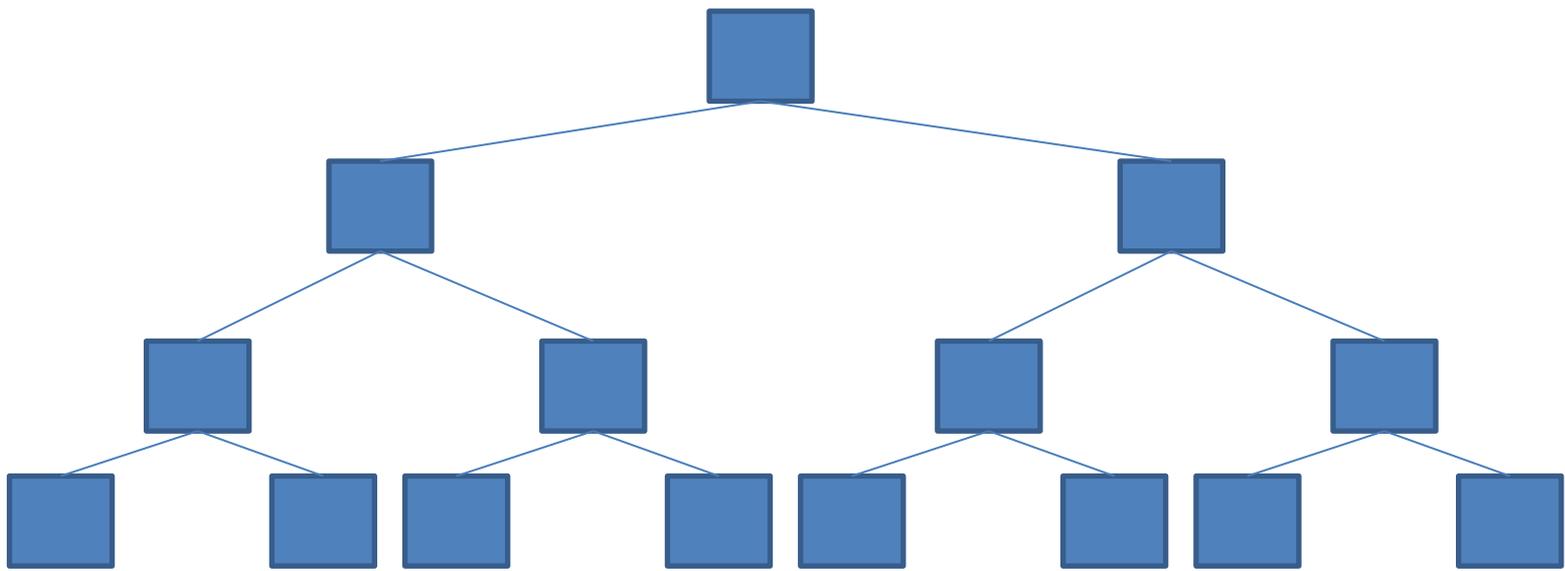
An implementation of a Heap data structure



To simplify implementation we **do not** use $S[0]$

Represent as a one *dimensional* array *S*

An implementation of a Heap data structure



last: 0

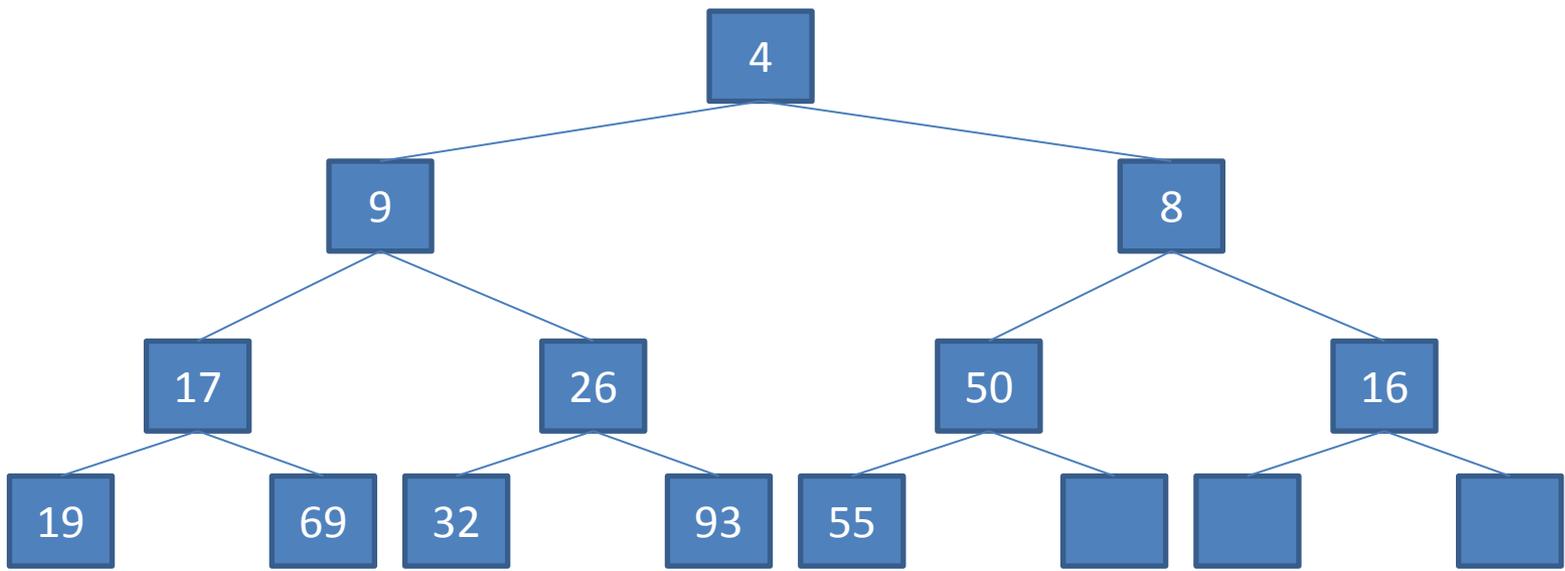
capacity: 15

An implementation of a Heap data structure

Consider the following heap H

Consider the following heap H

An implementation of a Heap data structure



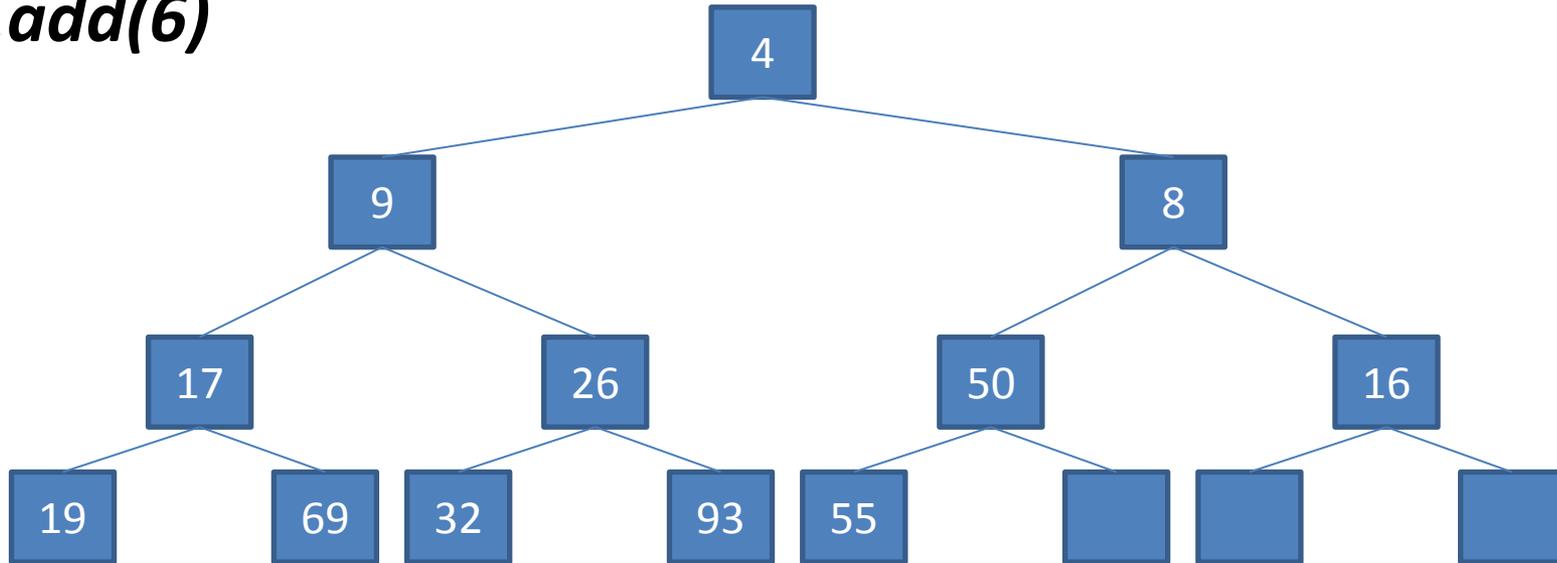
S

**	4	9	8	17	26	50	16	19	69	32	93	55			
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 12

capacity: 15

H.add(6)



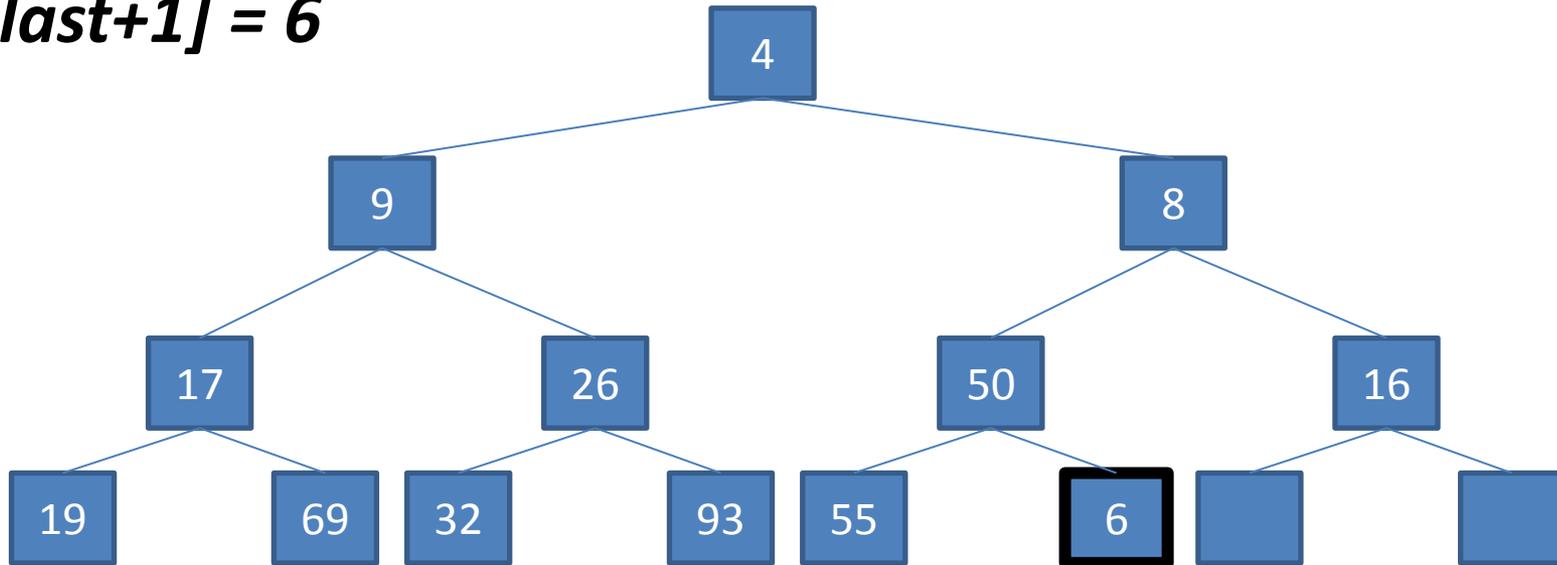
S

**	4	9	8	17	26	50	16	19	69	32	93	55			
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 12

capacity: 15

$S[\text{last}+1] = 6$

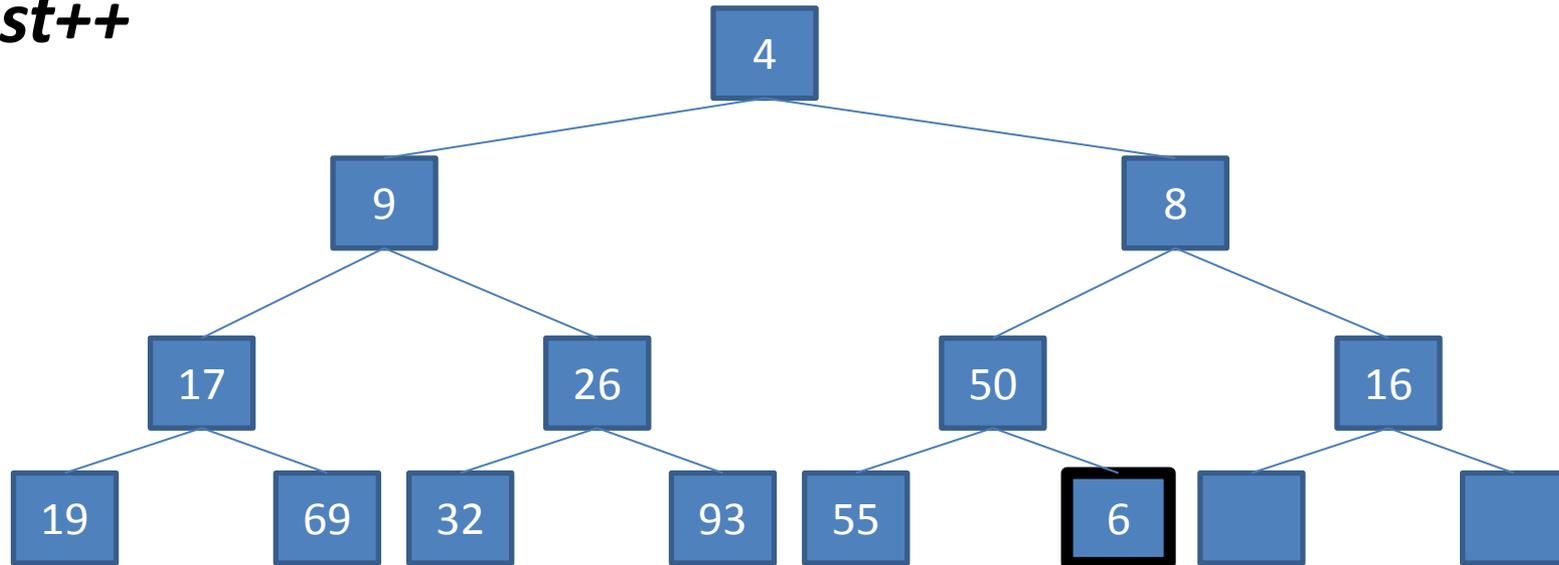


S	**	4	9	8	17	26	50	16	19	69	32	93	55	6		
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 12

capacity: 15

last++



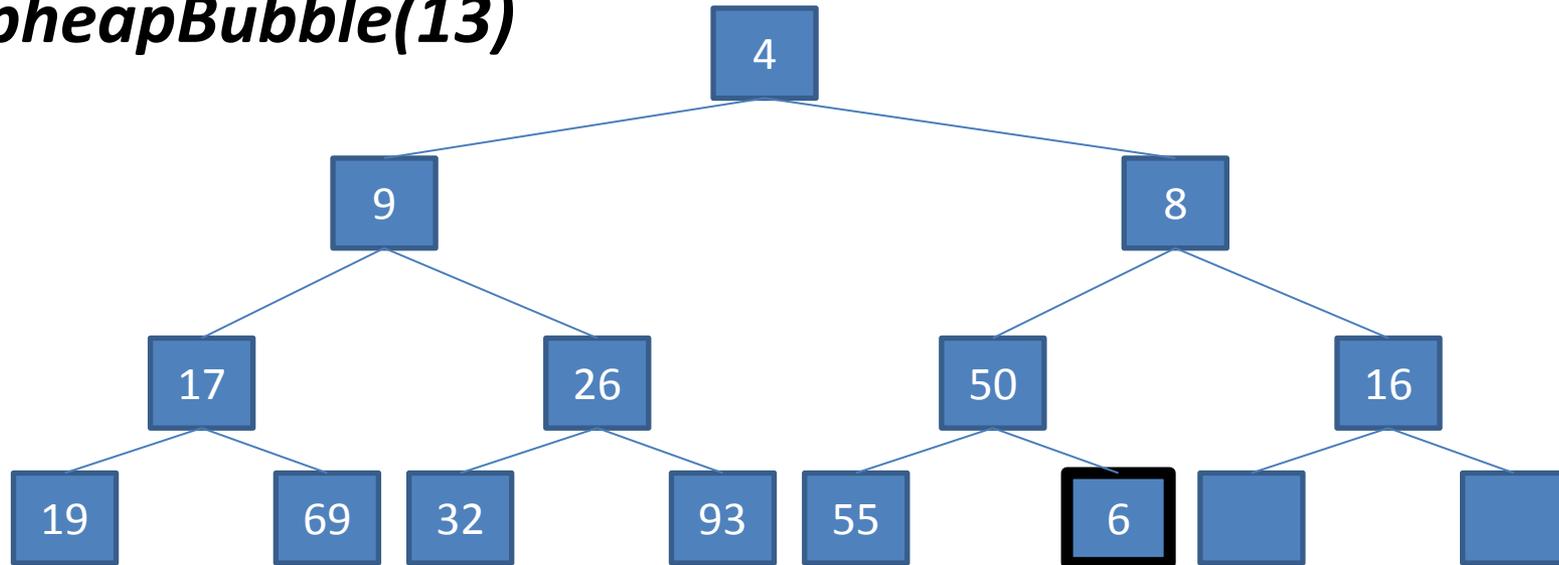
S

**	4	9	8	17	26	50	16	19	69	32	93	55	6		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

upheapBubble(13)



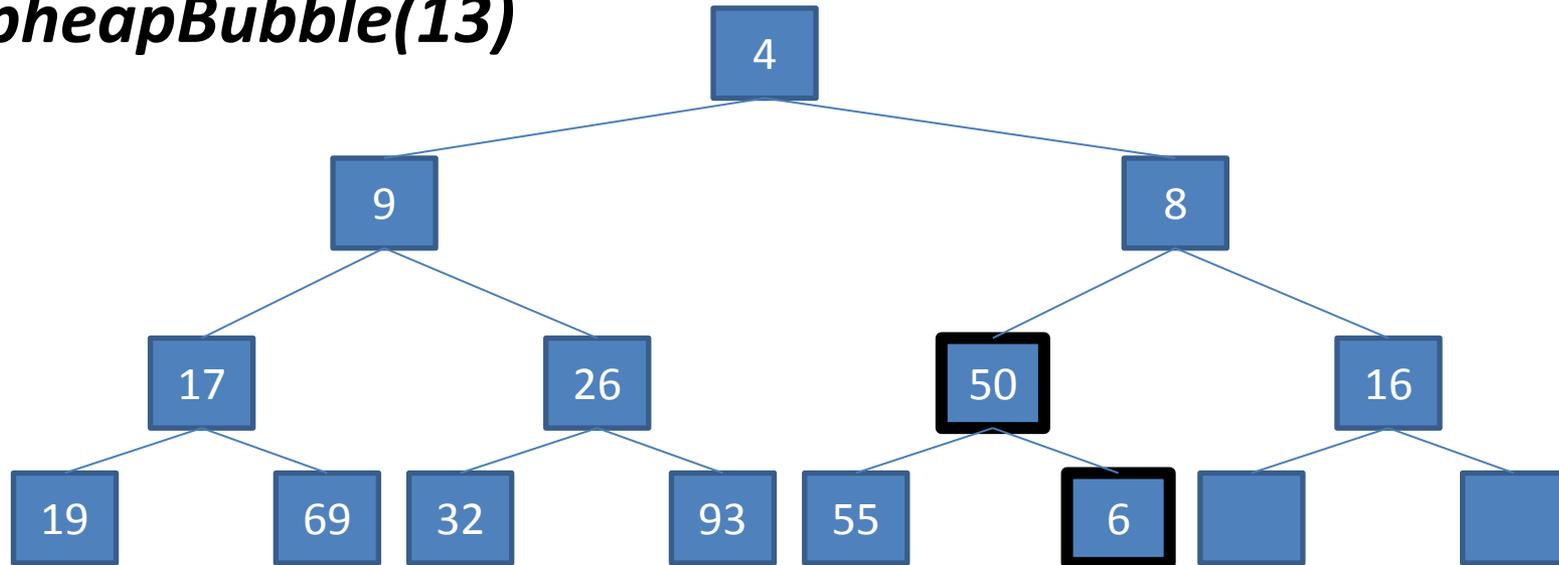
S

**	4	9	8	17	26	50	16	19	69	32	93	55	6		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

upheapBubble(13)



S

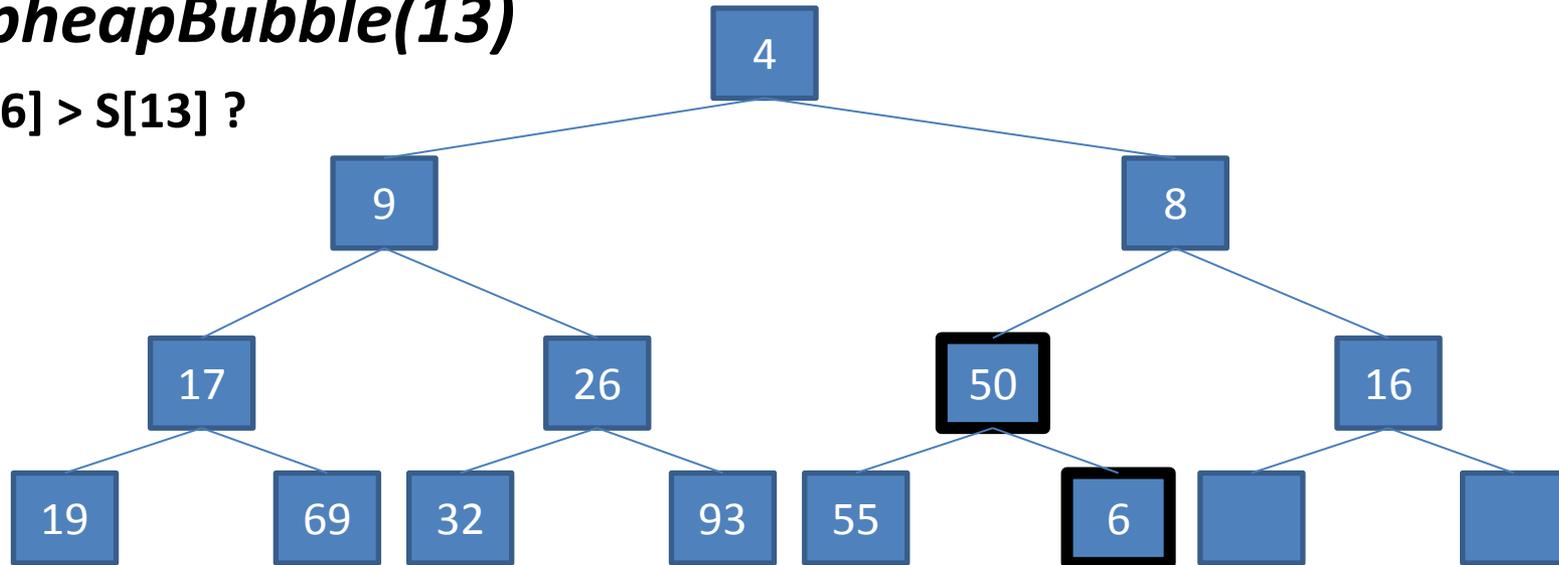
**	4	9	8	17	26	50	16	19	69	32	93	55	6		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

upheapBubble(13)

$S[6] > S[13]$?



S

**	4	9	8	17	26	50	16	19	69	32	93	55	6		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

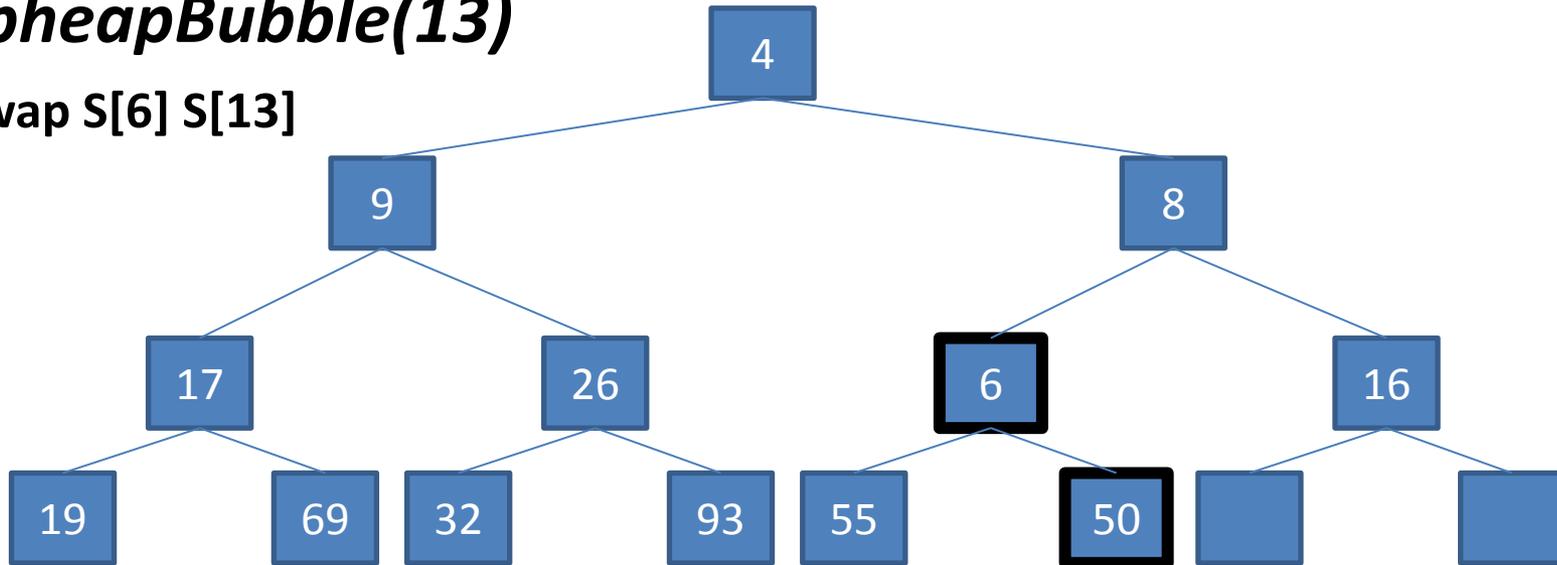


last: 13

capacity: 15

upheapBubble(13)

swap S[6] S[13]



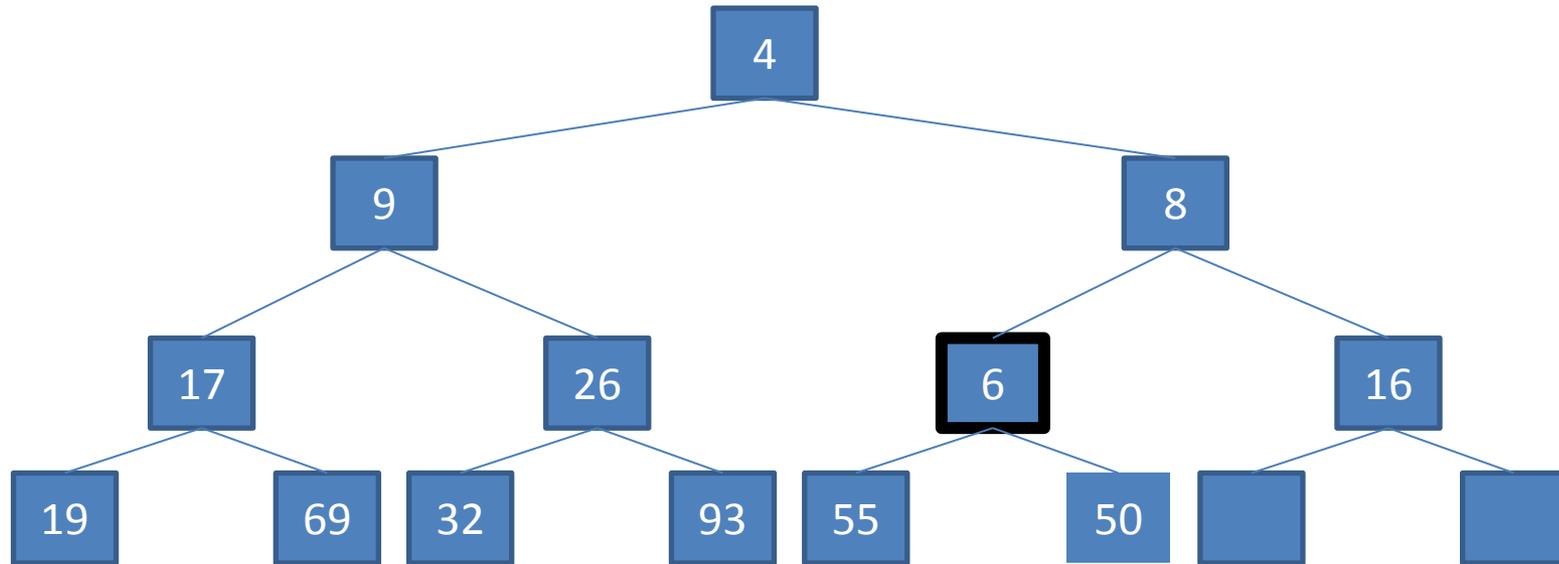
S	**	4	9	8	17	26	6	16	19	69	32	93	55	50		
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



last: 13

capacity: 15

An implementation of a Heap data structure



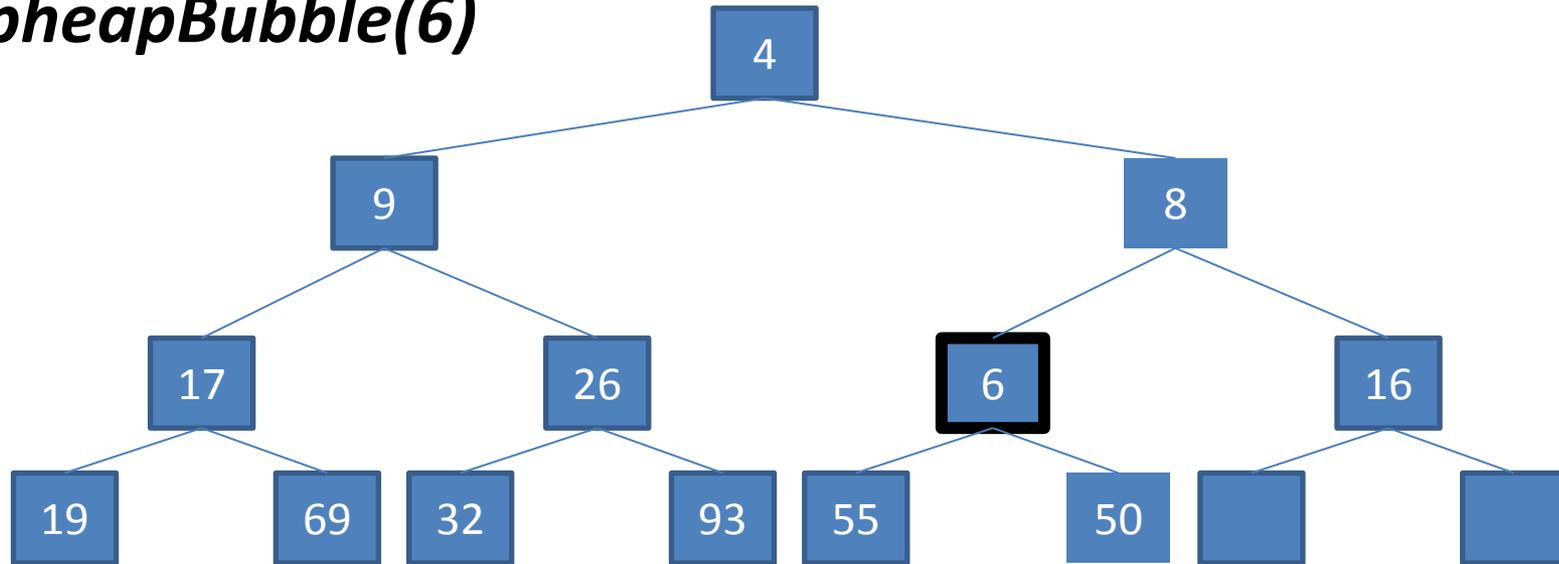
S

**	4	9	8	17	26	6	16	19	69	32	93	55	50		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

upheapBubble(6)



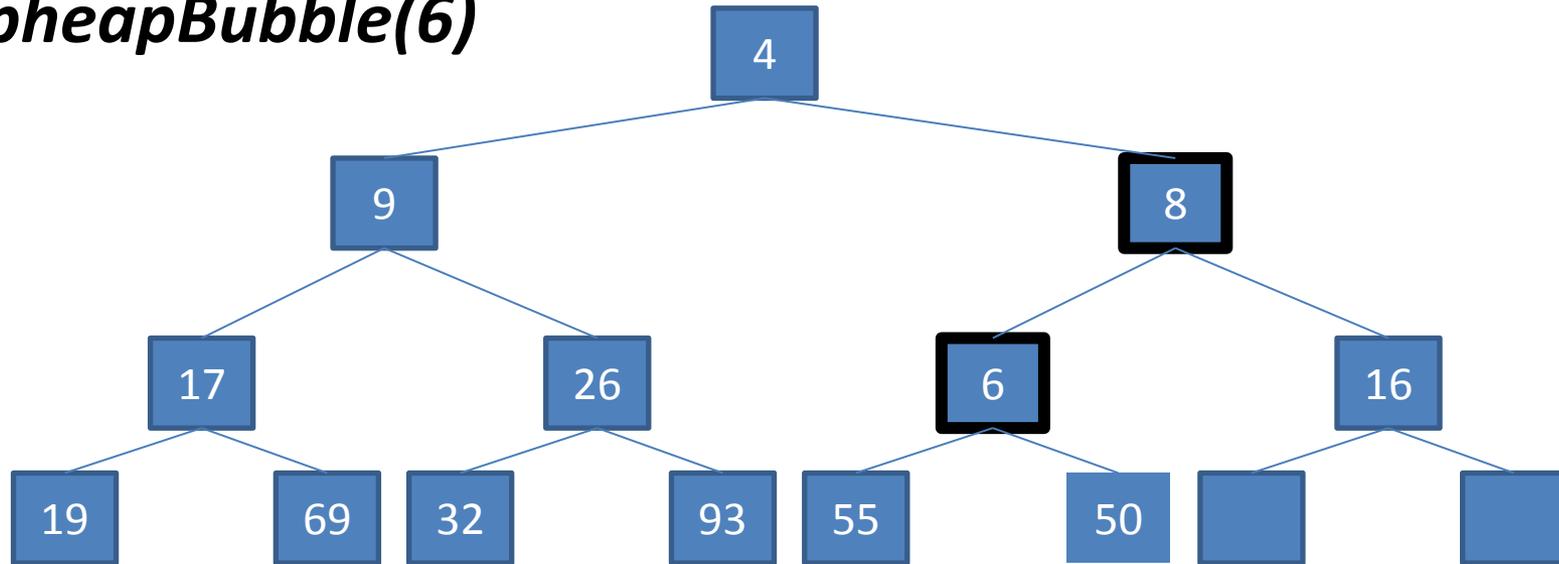
S

**	4	9	8	17	26	6	16	19	69	32	93	55	50		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

upheapBubble(6)



S

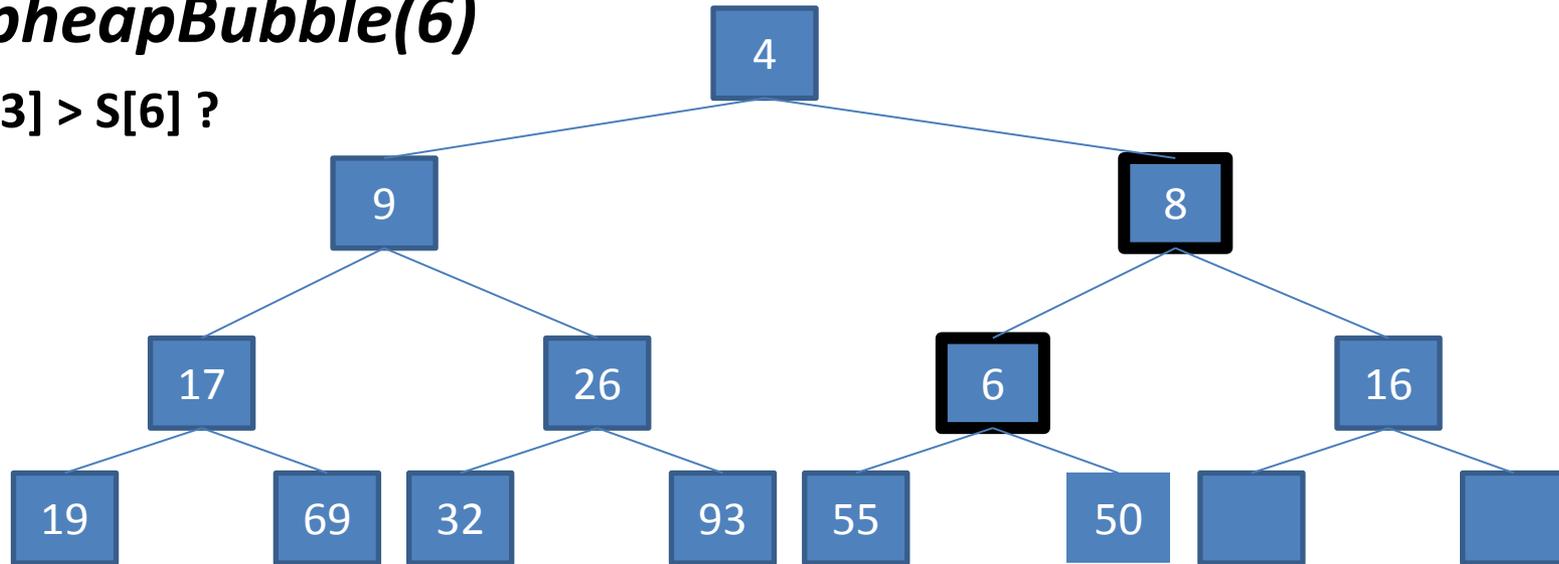
**	4	9	8	17	26	6	16	19	69	32	93	55	50		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

upheapBubble(6)

$S[3] > S[6]$?



S	**	4	9	8	17	26	6	16	19	69	32	93	55	50		
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

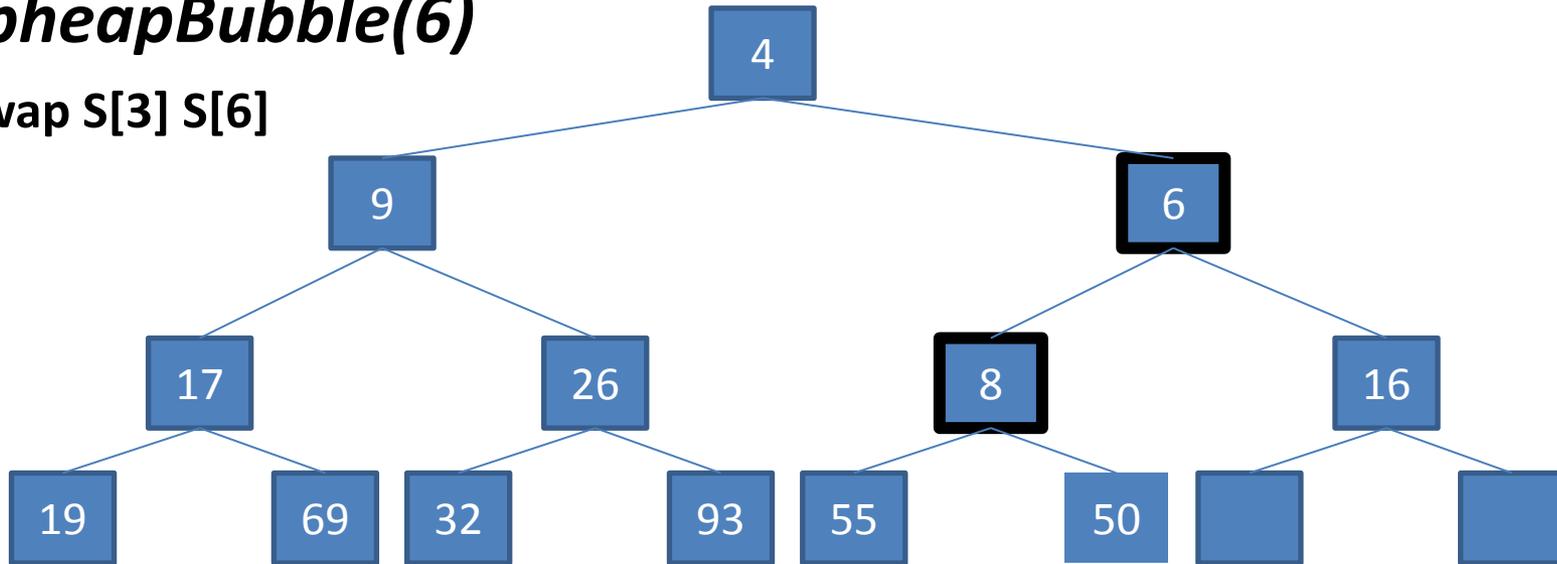
last: 13

capacity: 15



upheapBubble(6)

swap S[3] S[6]



S

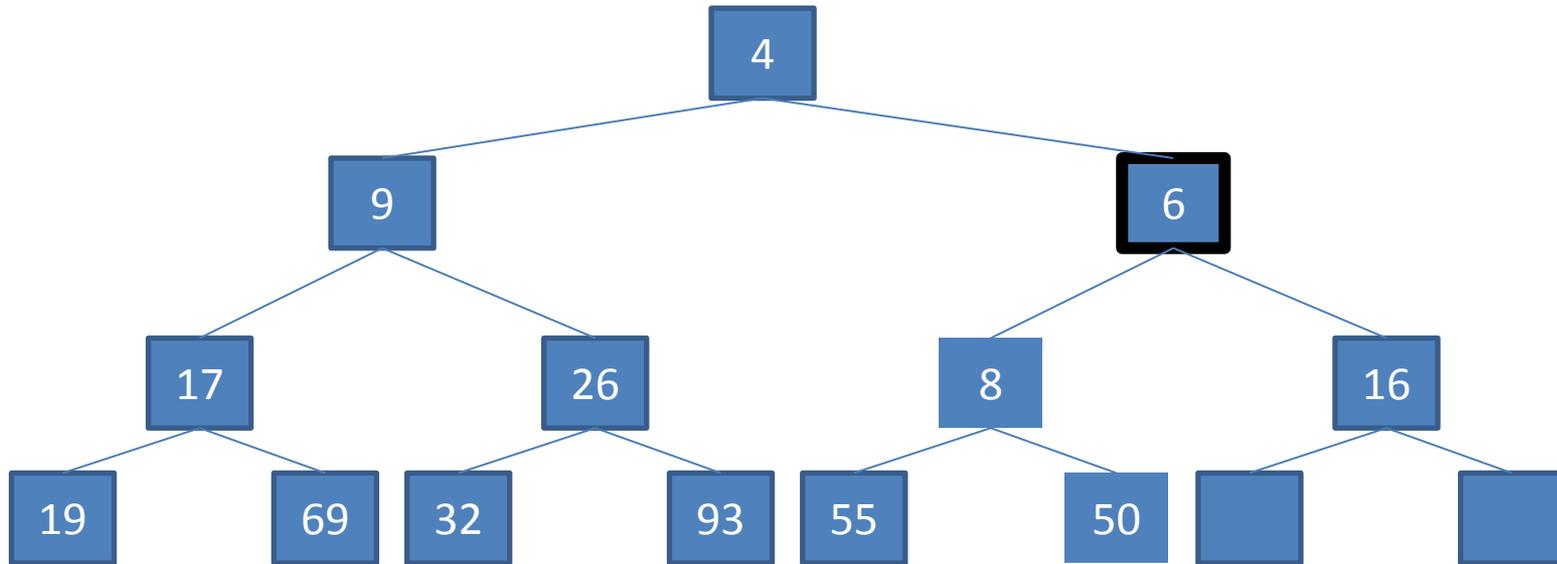
**	4	9	6	17	26	8	16	19	69	32	93	55	50		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15



An implementation of a Heap data structure



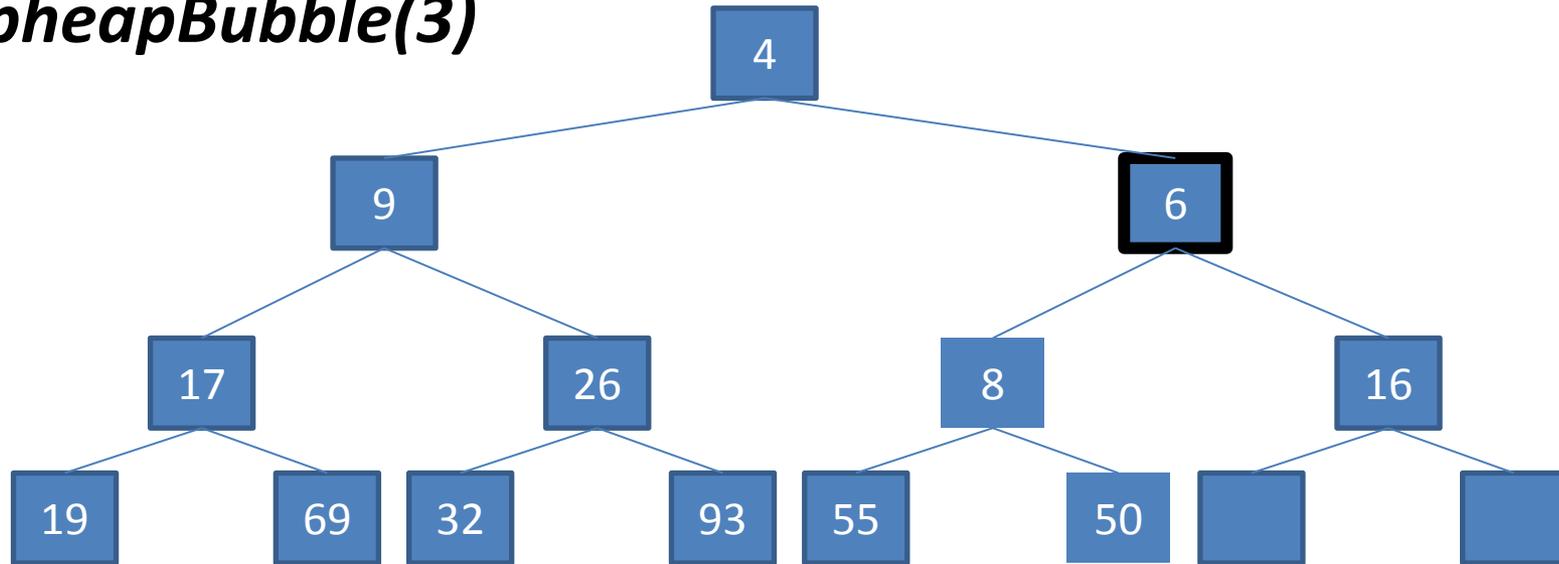
S

**	4	9	6	17	26	8	16	19	69	32	93	55	50		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

upheapBubble(3)



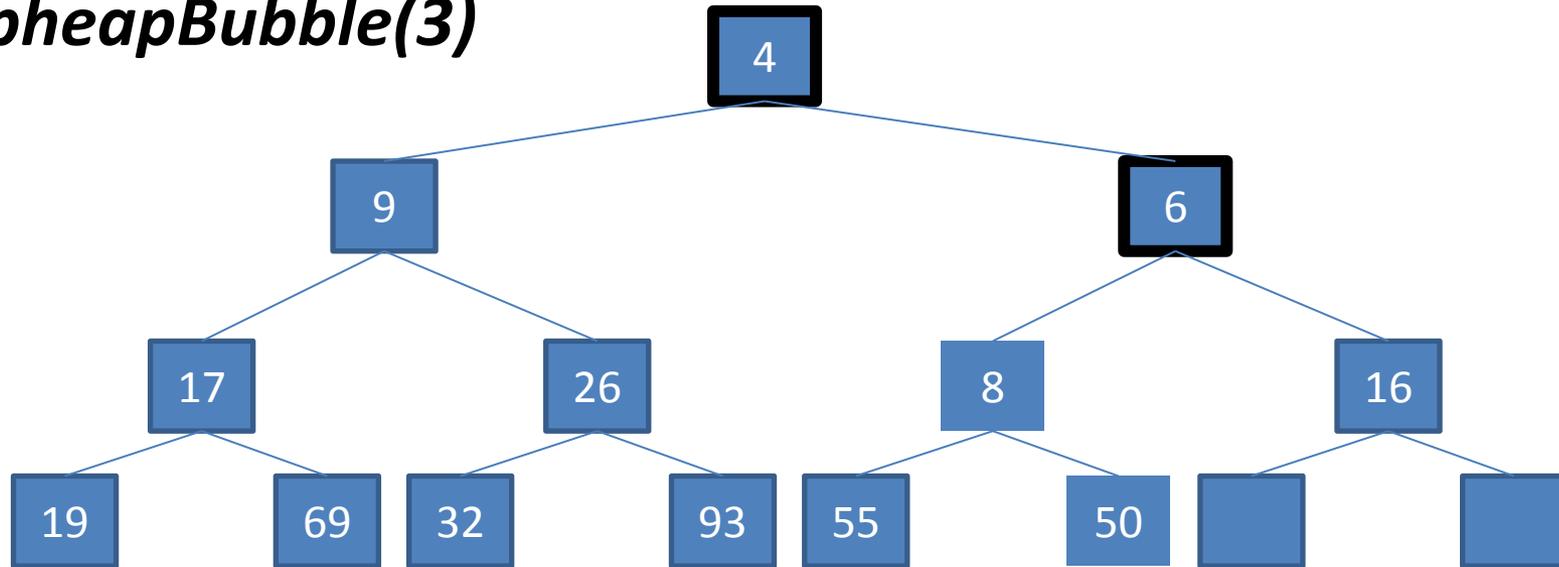
S

**	4	9	6	17	26	8	16	19	69	32	93	55	50		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

upheapBubble(3)



S

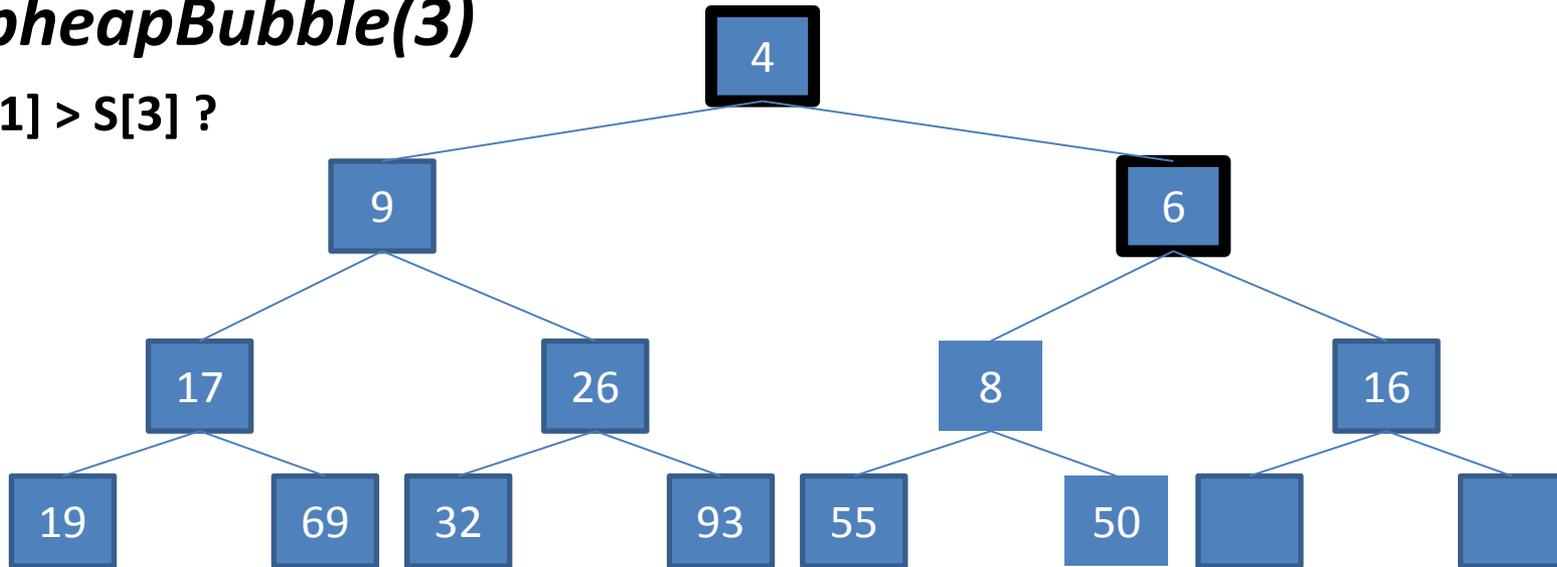
**	4	9	6	17	26	8	16	19	69	32	93	55	50		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

upheapBubble(3)

$S[1] > S[3]$?



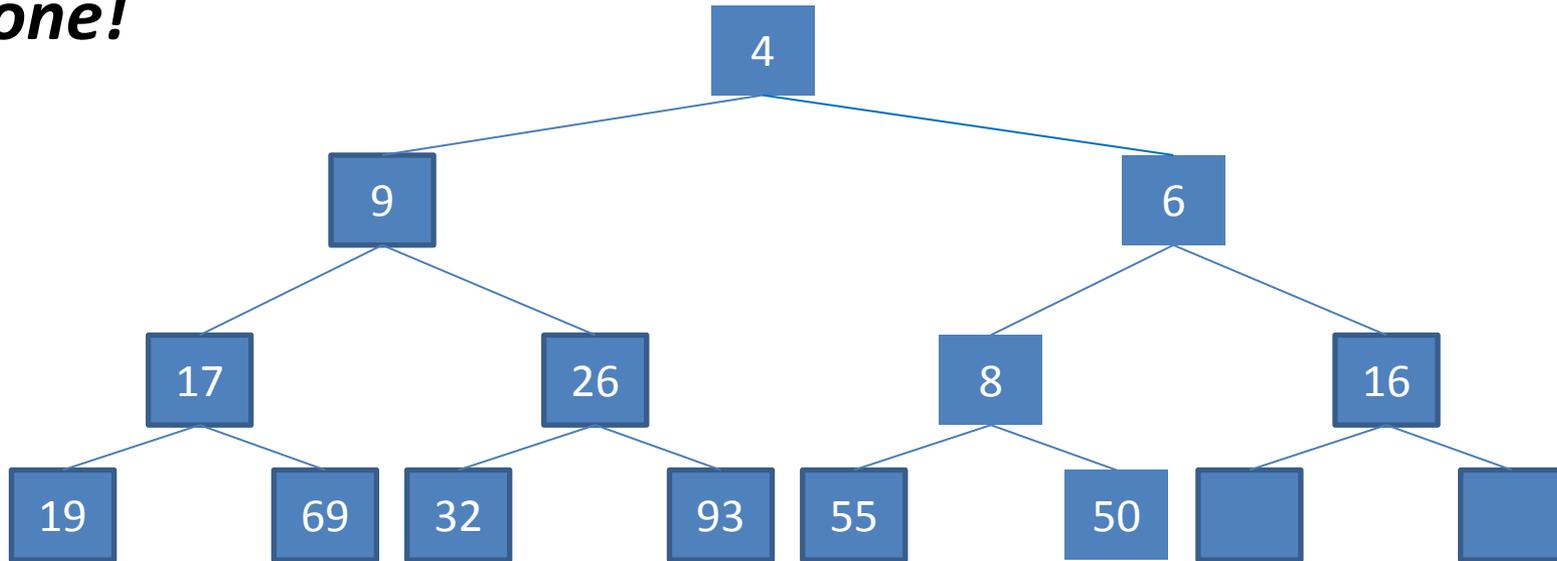
S	**	4	9	6	17	26	8	16	19	69	32	93	55	50		
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15



Done!



S

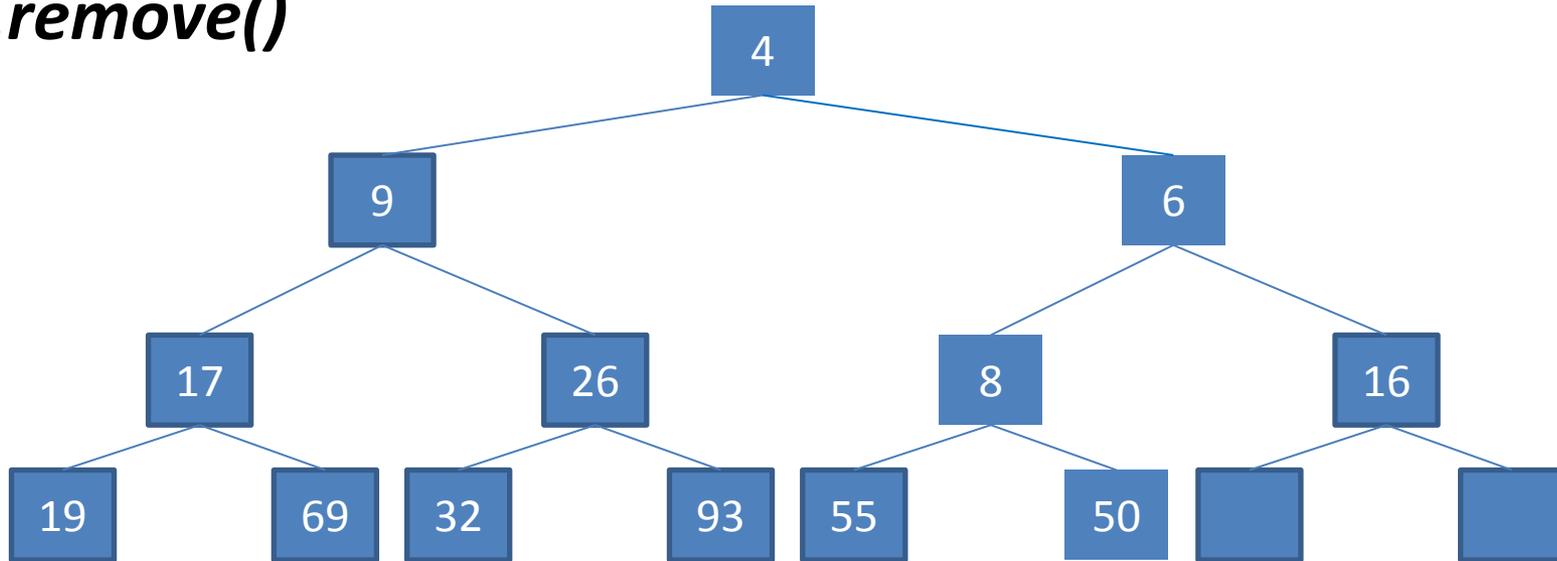
**	4	9	6	17	26	8	16	19	69	32	93	55	50		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

Removal from the heap H

H.remove()



S

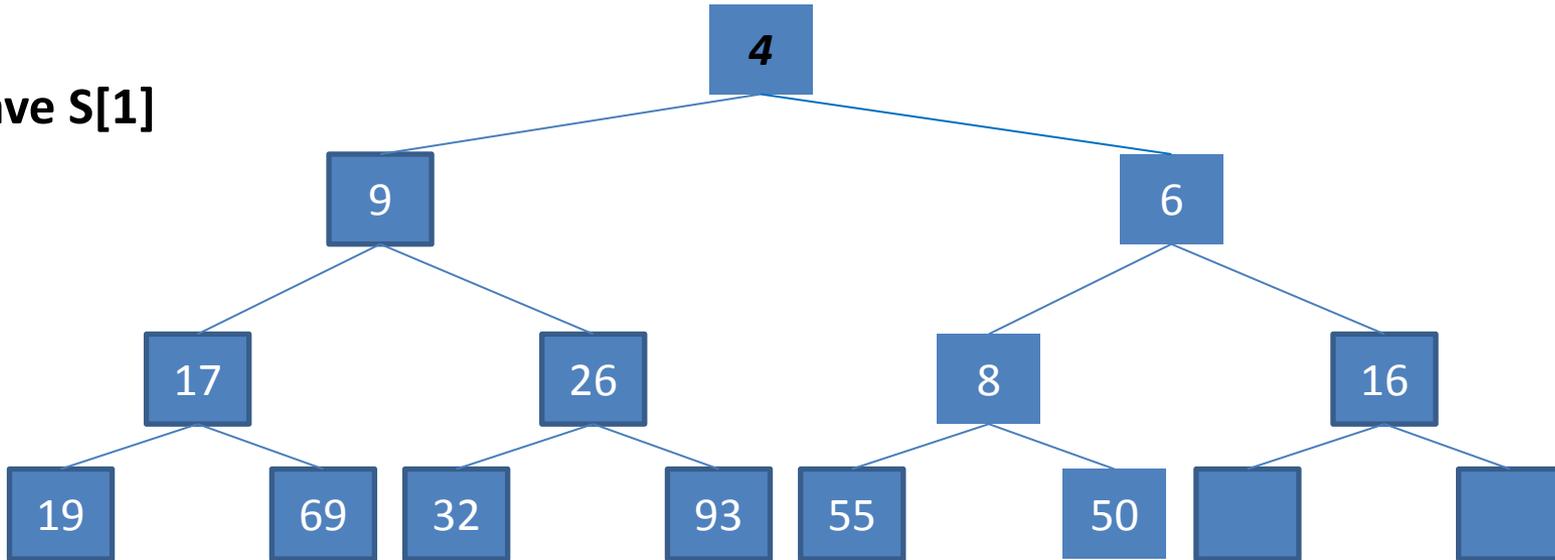
**	4	9	6	17	26	8	16	19	69	32	93	55	50		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

4

Save S[1]



S

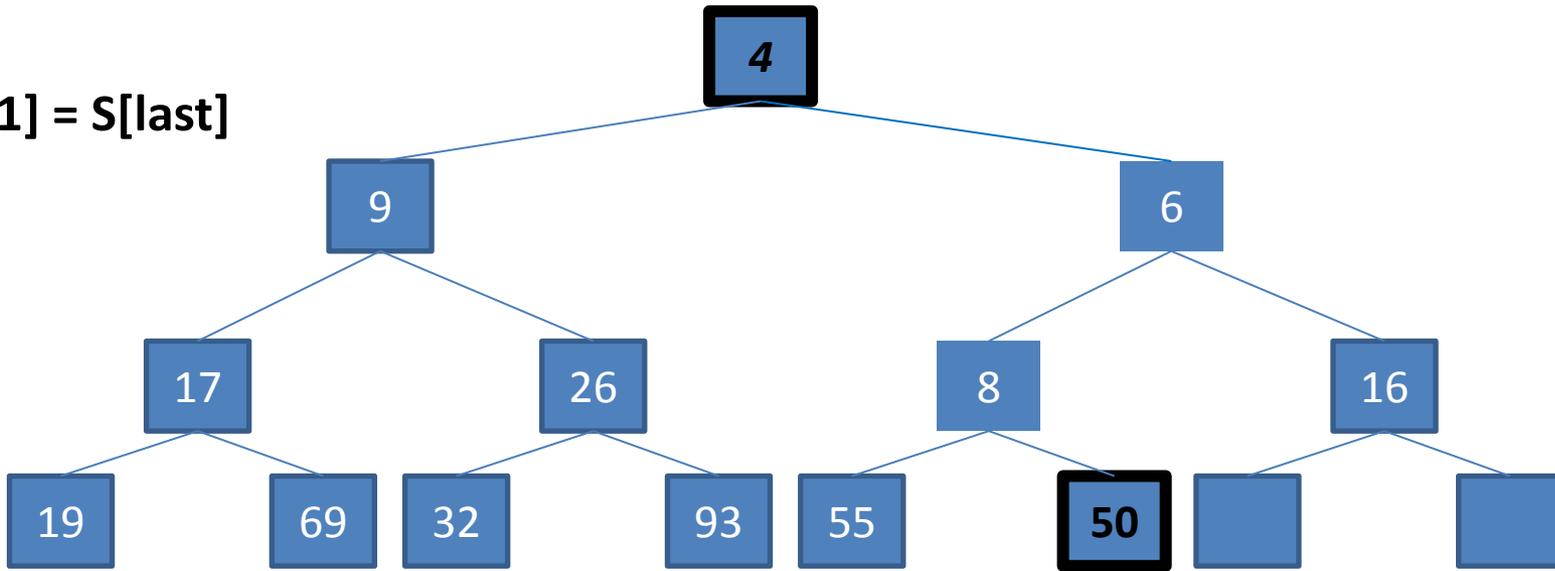
**	4	9	6	17	26	8	16	19	69	32	93	55	50		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

4

$S[1] = S[\text{last}]$



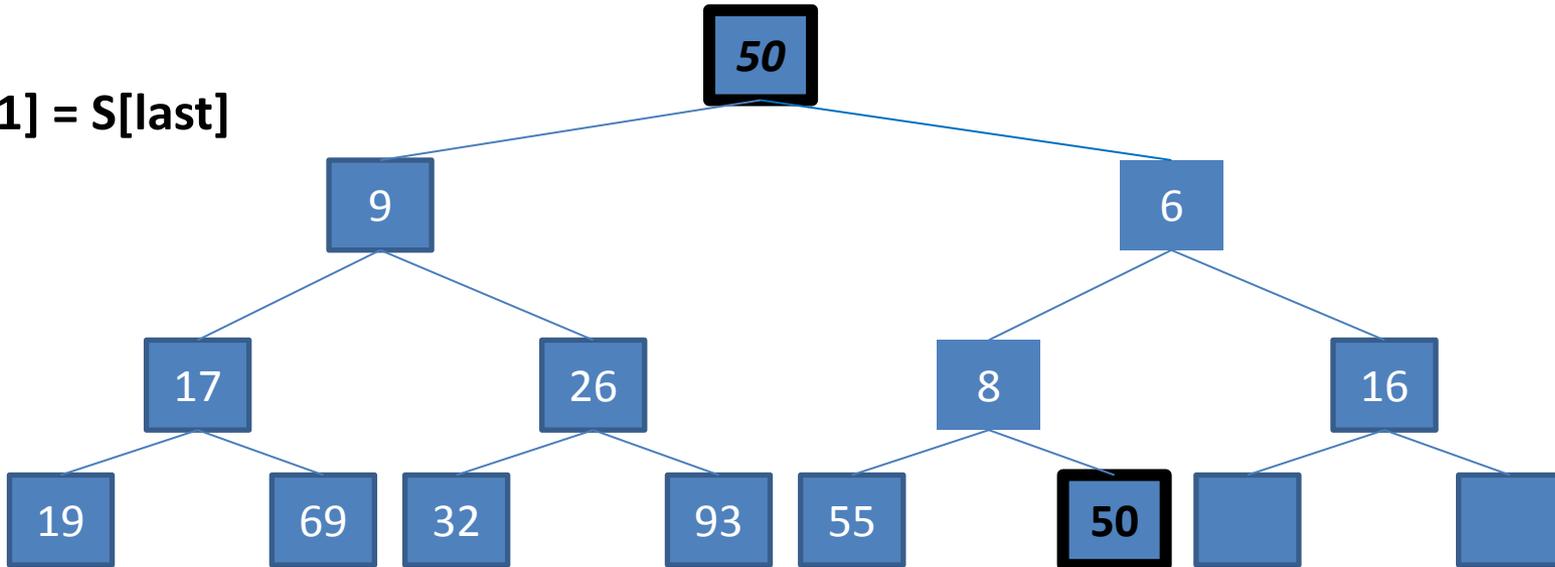
S	**	4	9	6	17	26	8	16	19	69	32	93	55	50		
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

4

$S[1] = S[\text{last}]$



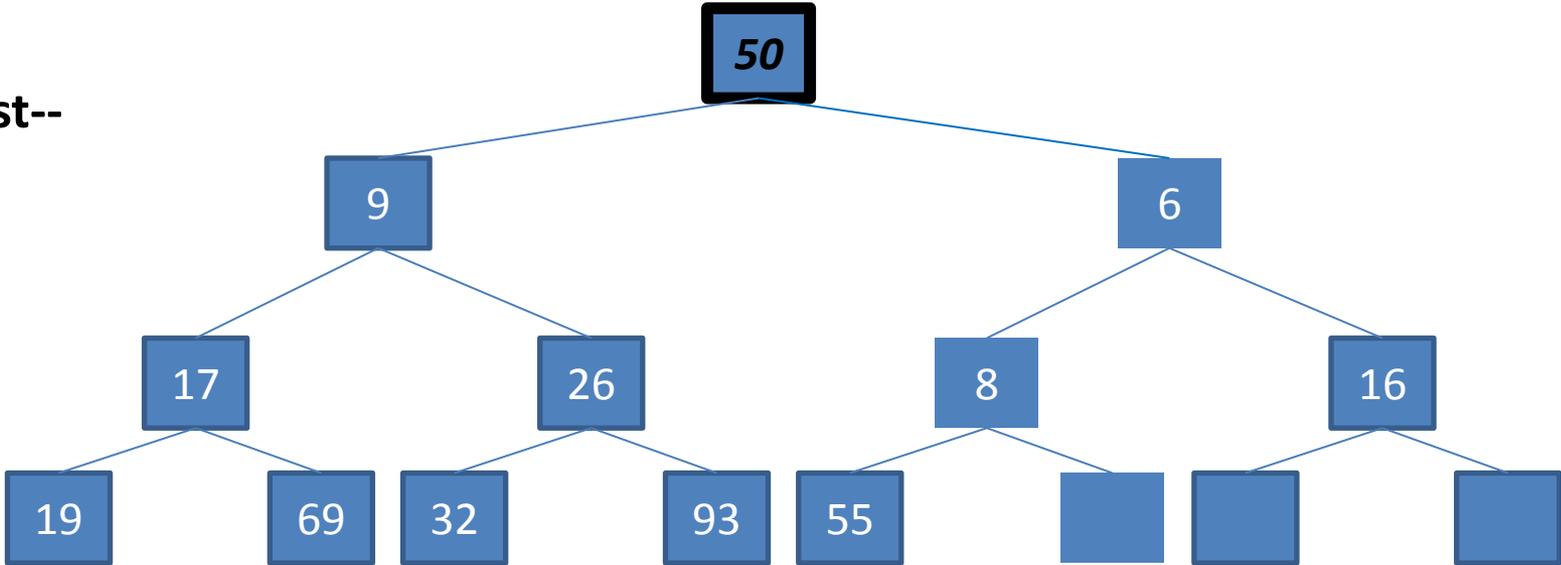
S	**	50	9	6	17	26	8	16	19	69	32	93	55	50		
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 13

capacity: 15

4

last--



S

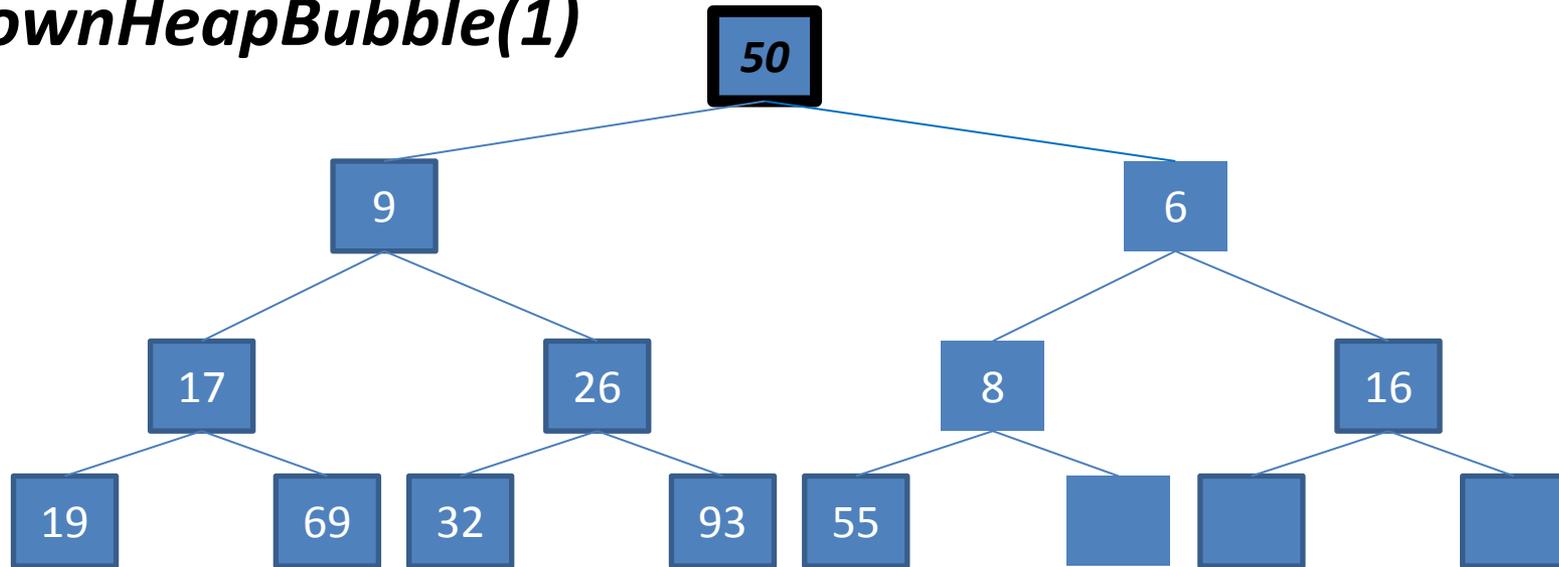
**	50	9	6	17	26	8	16	19	69	32	93	55			
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 12

capacity: 15

4

downHeapBubble(1)



S

**	50	9	6	17	26	8	16	19	69	32	93	55			
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

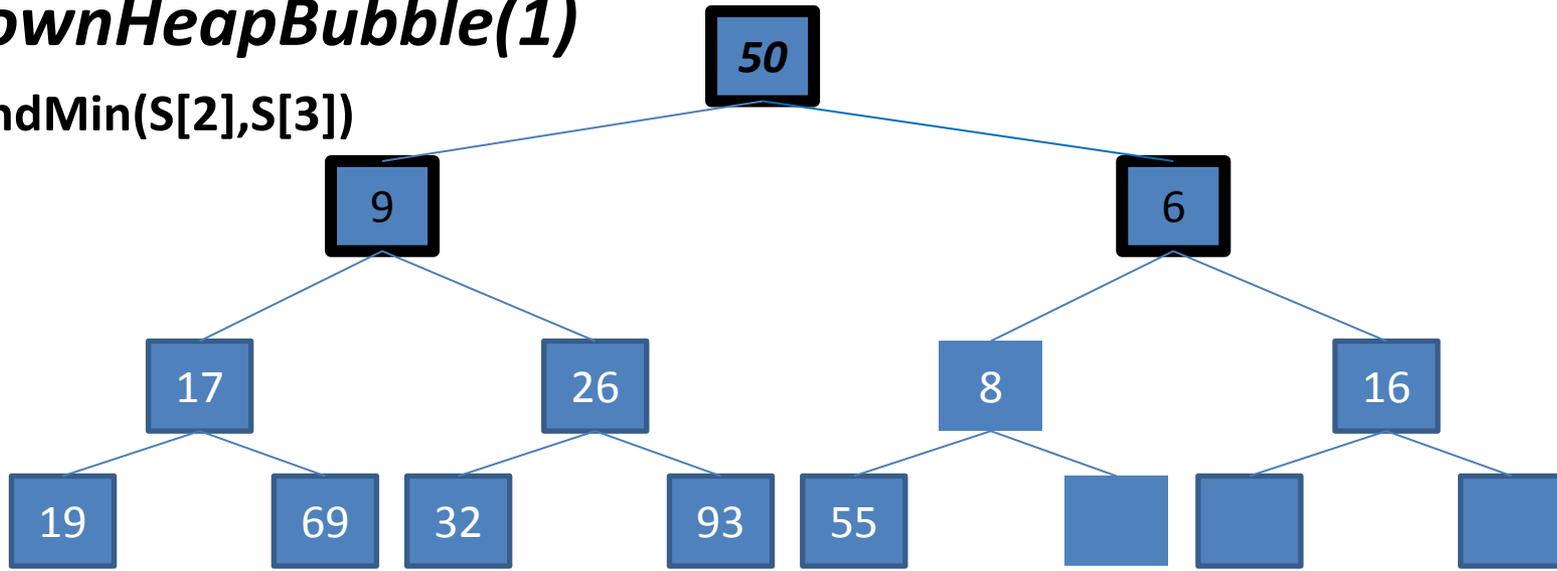
last: 12

capacity: 15

4

downHeapBubble(1)

findMin(S[2],S[3])



	**	50	9	6	17	26	8	16	19	69	32	93	55			
<i>S</i>	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 12

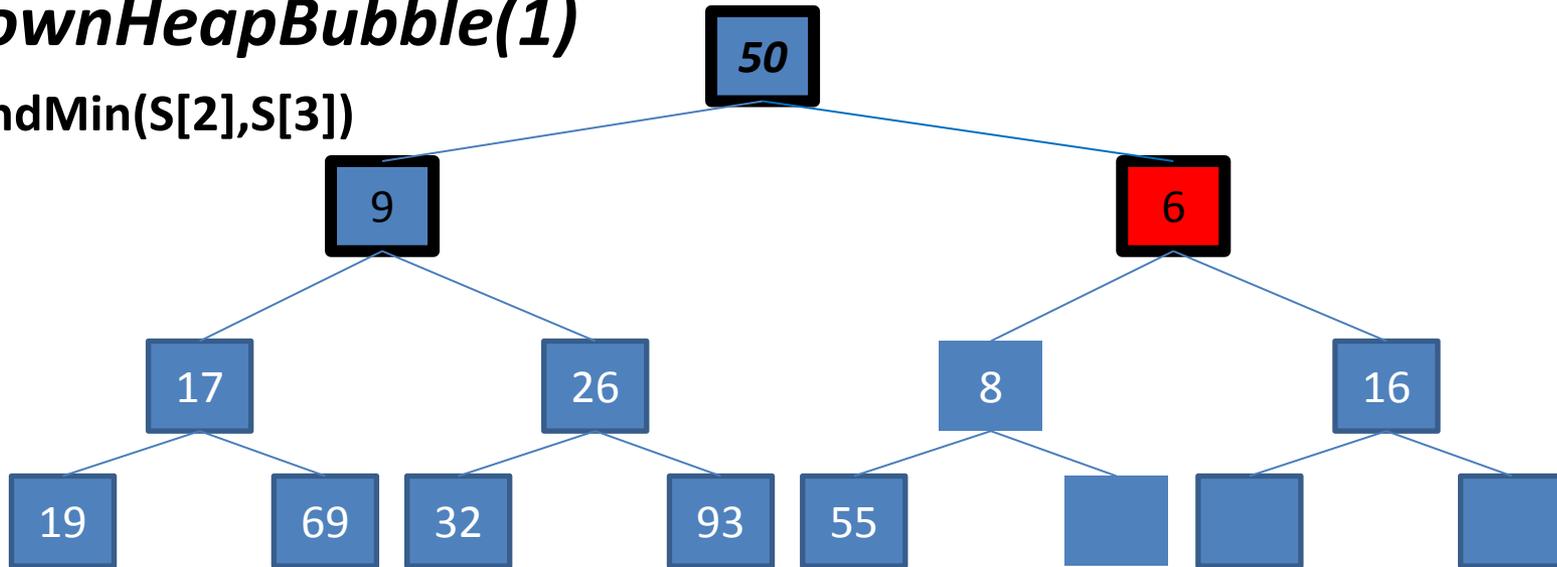
capacity: 15



4

downHeapBubble(1)

findMin(S[2],S[3])



	**	50	9	6	17	26	8	16	19	69	32	93	55			
S	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 12

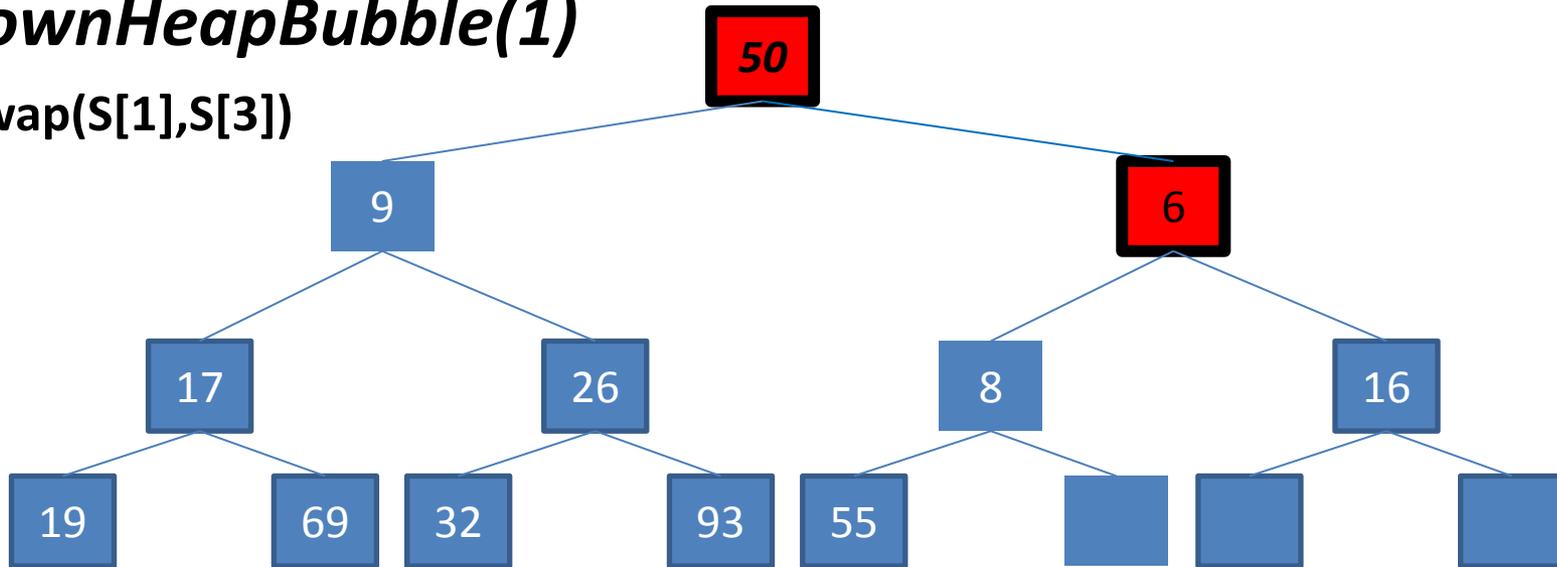
capacity: 15



4

downHeapBubble(1)

swap(S[1],S[3])



	**	50	9	6	17	26	8	16	19	69	32	93	55			
S	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 12

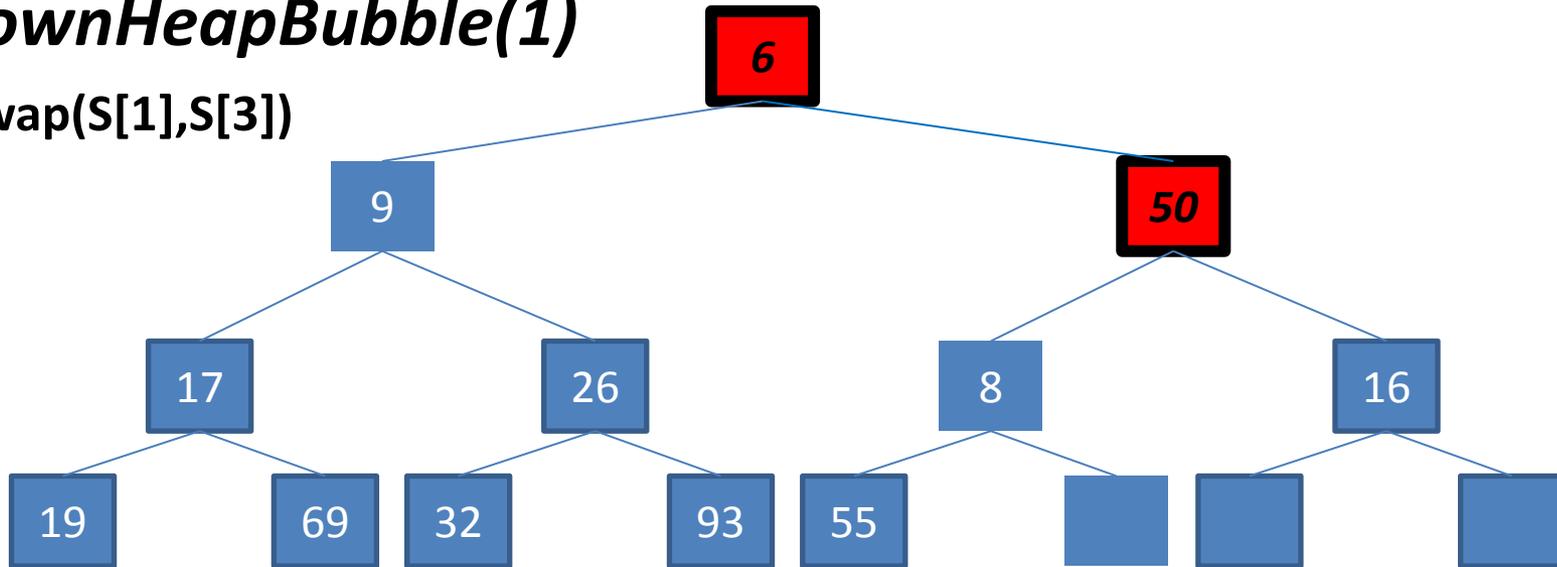
capacity: 15



4

downHeapBubble(1)

swap(S[1],S[3])



	**	6	9	50	17	26	8	16	19	69	32	93	55			
S	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

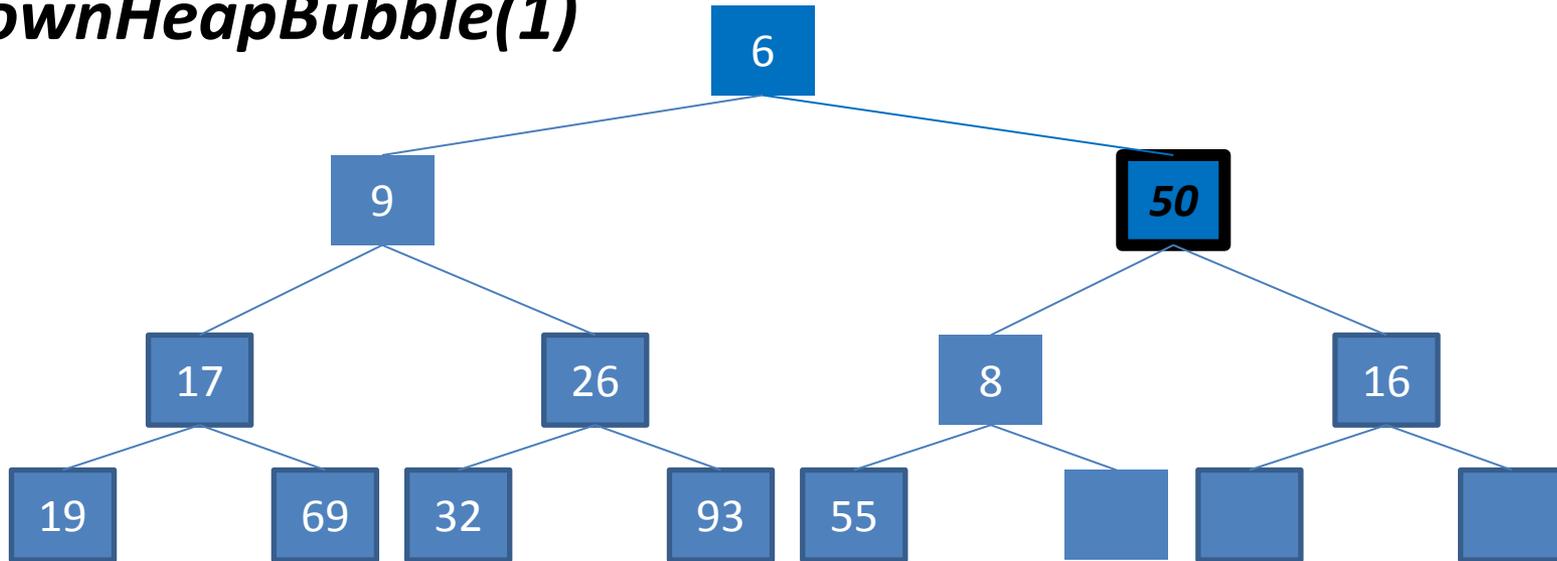
last: 12

capacity: 15



4

downHeapBubble(1)



S

**	6	9	50	17	26	8	16	19	69	32	93	55			
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

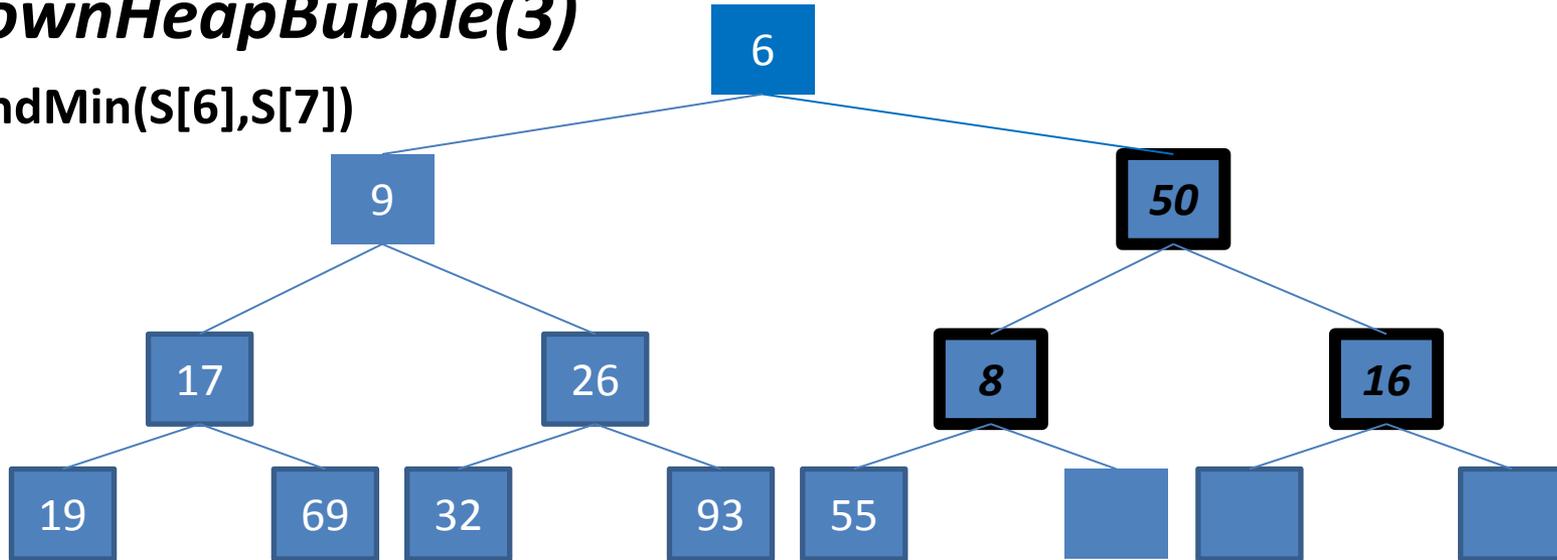
last: 12

capacity: 15

4

downHeapBubble(3)

findMin(S[6],S[7])



S	**	6	9	50	17	26	8	16	19	69	32	93	55			
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



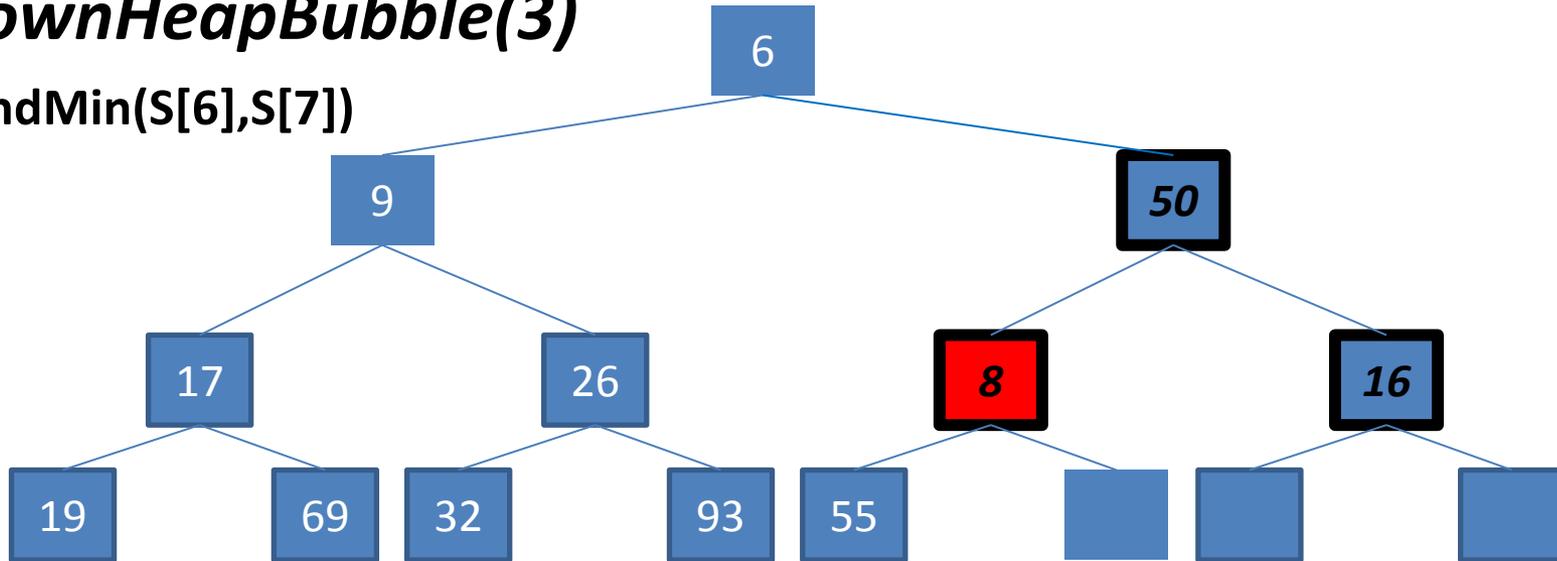
last: 12

capacity: 15

4

downHeapBubble(3)

findMin(S[6],S[7])



S	**	6	9	50	17	26	8	16	19	69	32	93	55			
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 12

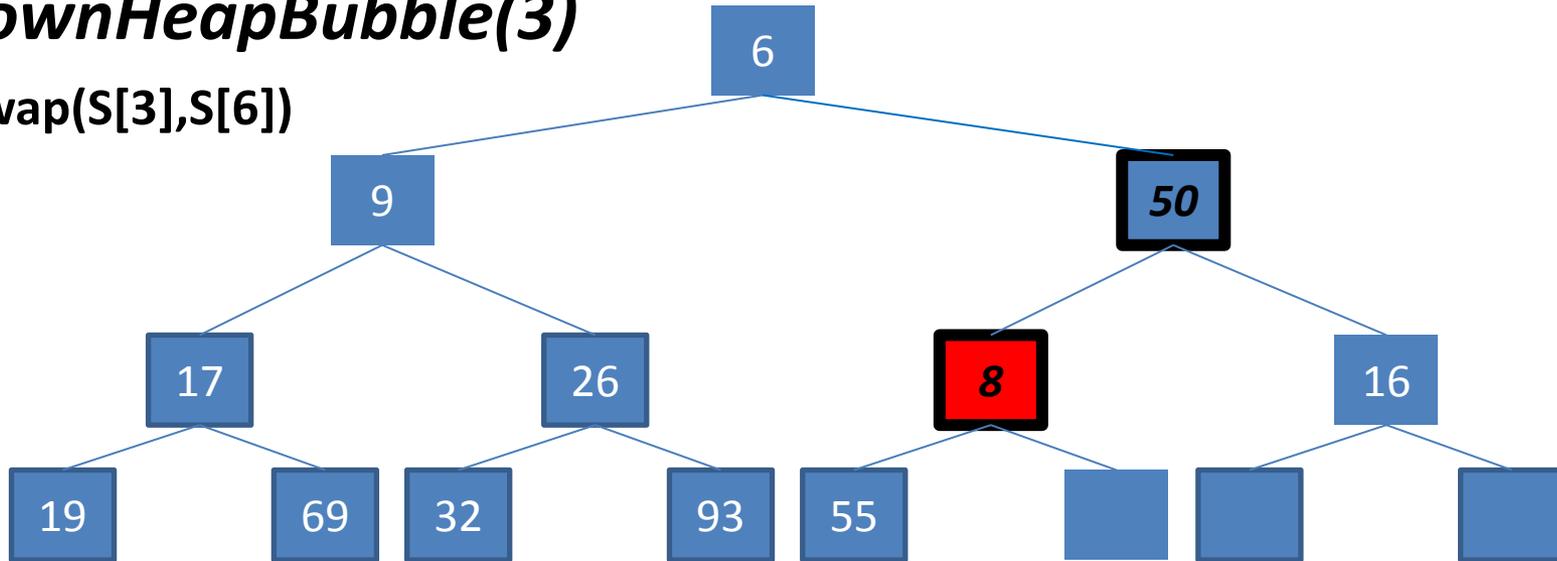
capacity: 15



4

downHeapBubble(3)

swap(S[3],S[6])



S	**	6	9	50	17	26	8	16	19	69	32	93	55			
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



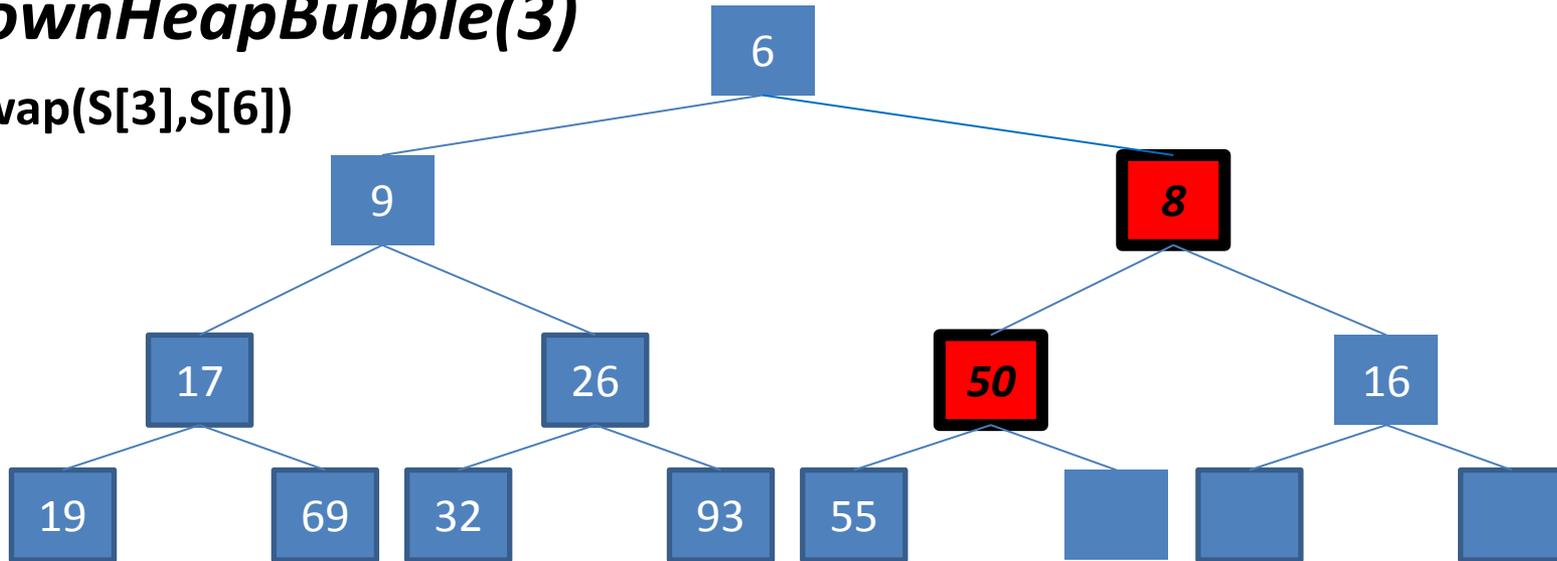
last: 12

capacity: 15

4

downHeapBubble(3)

swap(S[3],S[6])



S

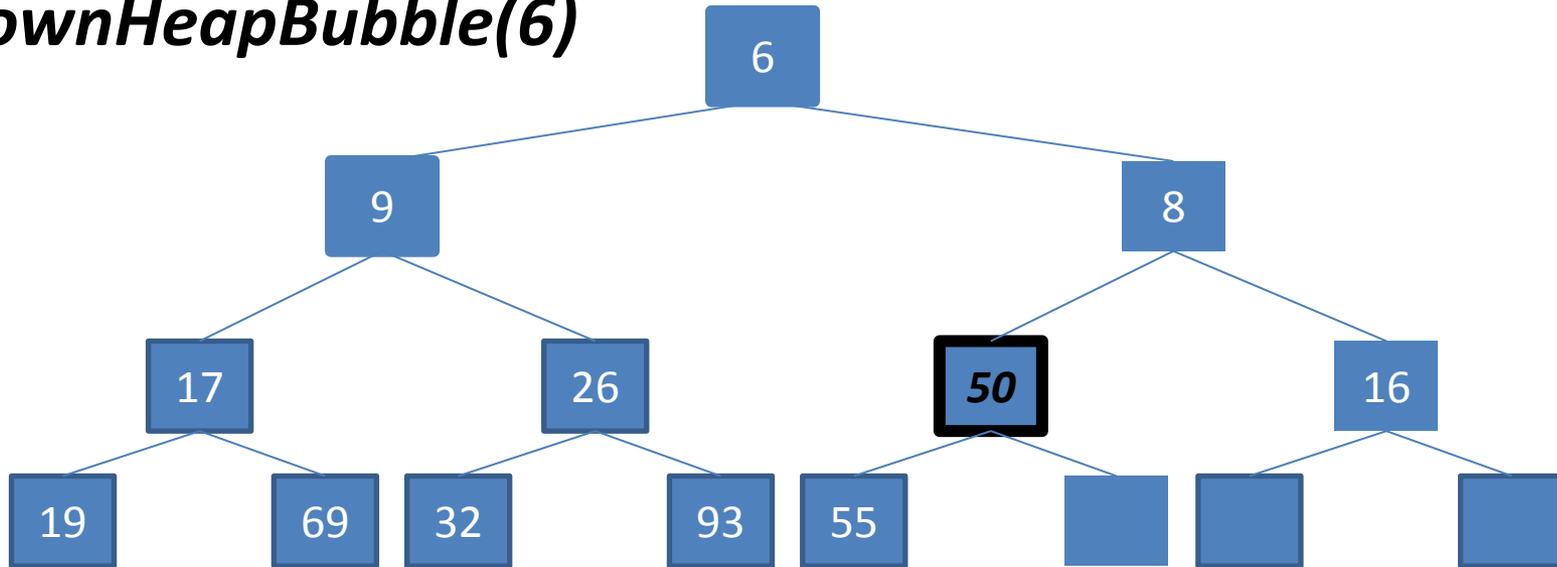
**	6	9	8	17	26	50	16	19	69	32	93	55			
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 12

capacity: 15



4

downHeapBubble(6)

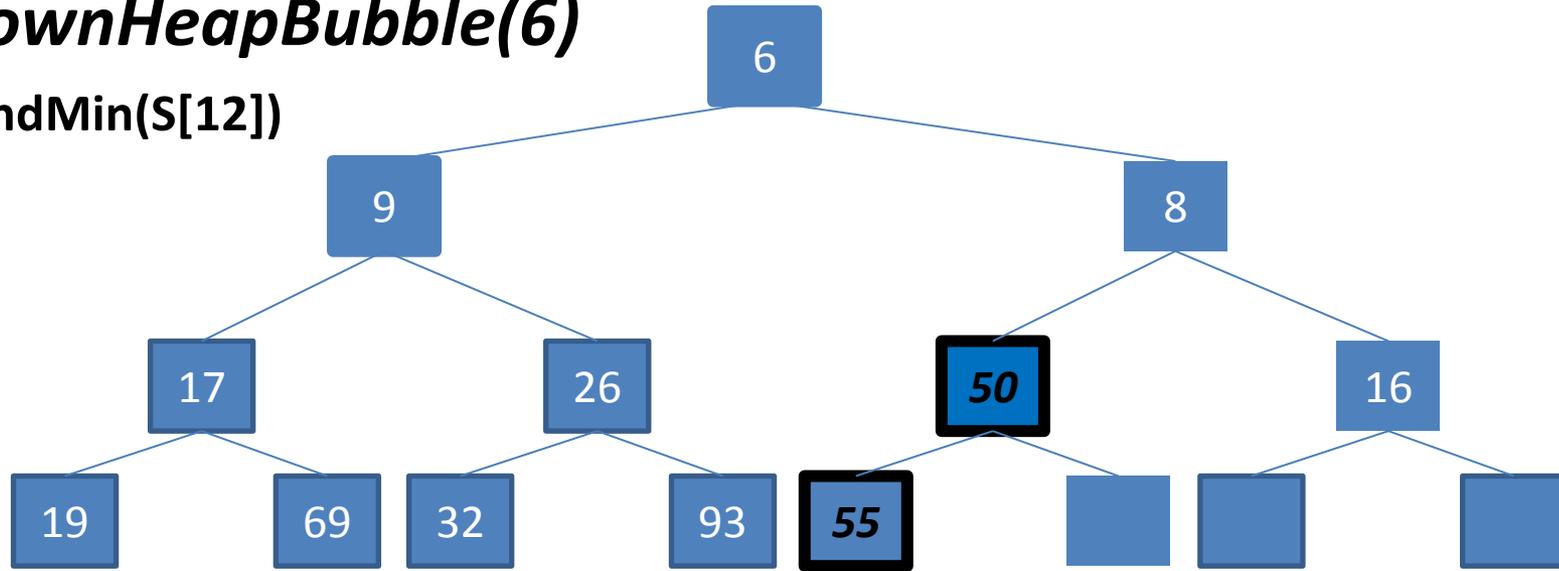
S	**	6	9	8	17	26	50	16	19	69	32	93	55			
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

*last: 12**capacity: 15*

4

downHeapBubble(6)

findMin(S[12])



S

**	6	9	8	17	26	50	16	19	69	32	93	55			
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



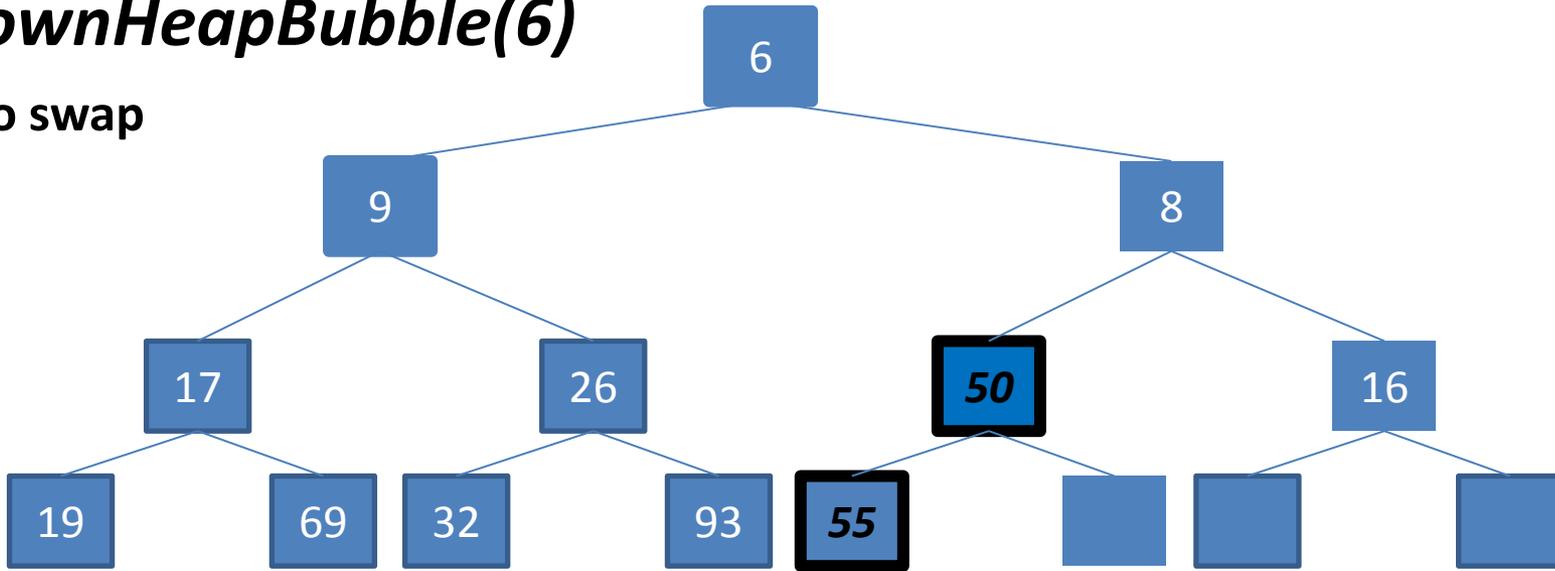
last: 12

capacity: 15

4

downHeapBubble(6)

No swap



S

**	6	9	8	17	26	50	16	19	69	32	93	55			
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



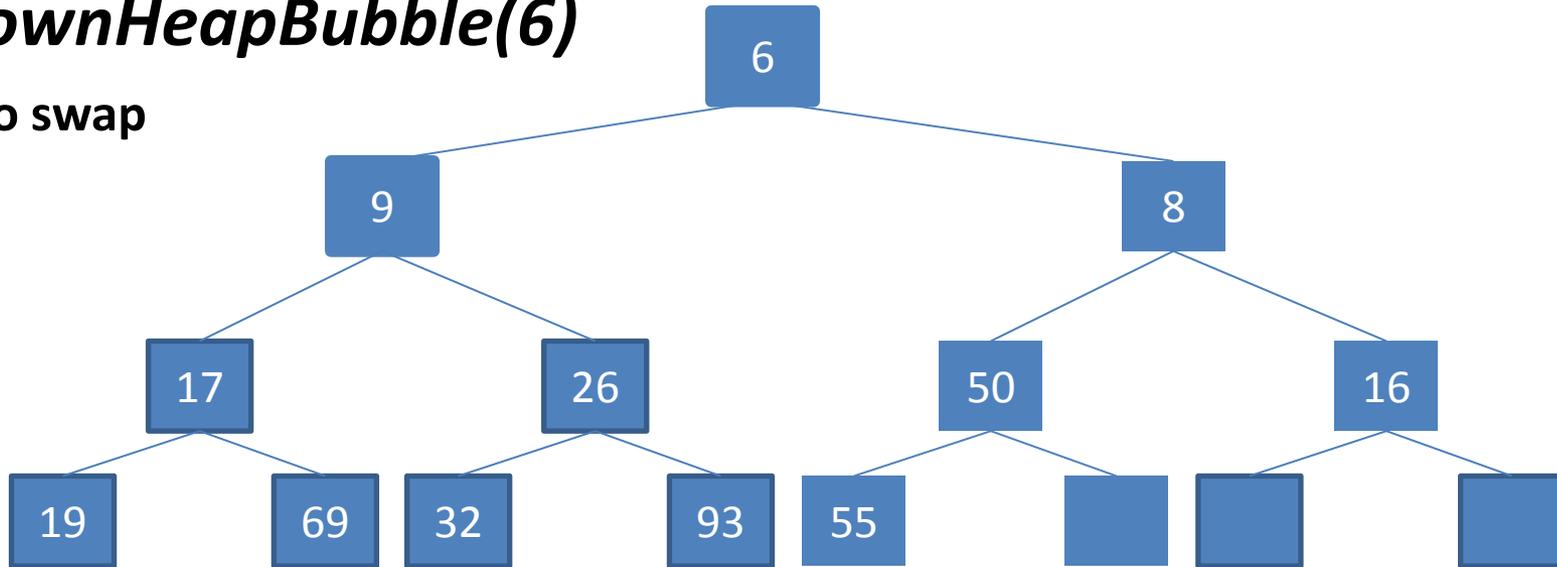
last: 12

capacity: 15

4

downHeapBubble(6)

No swap



S

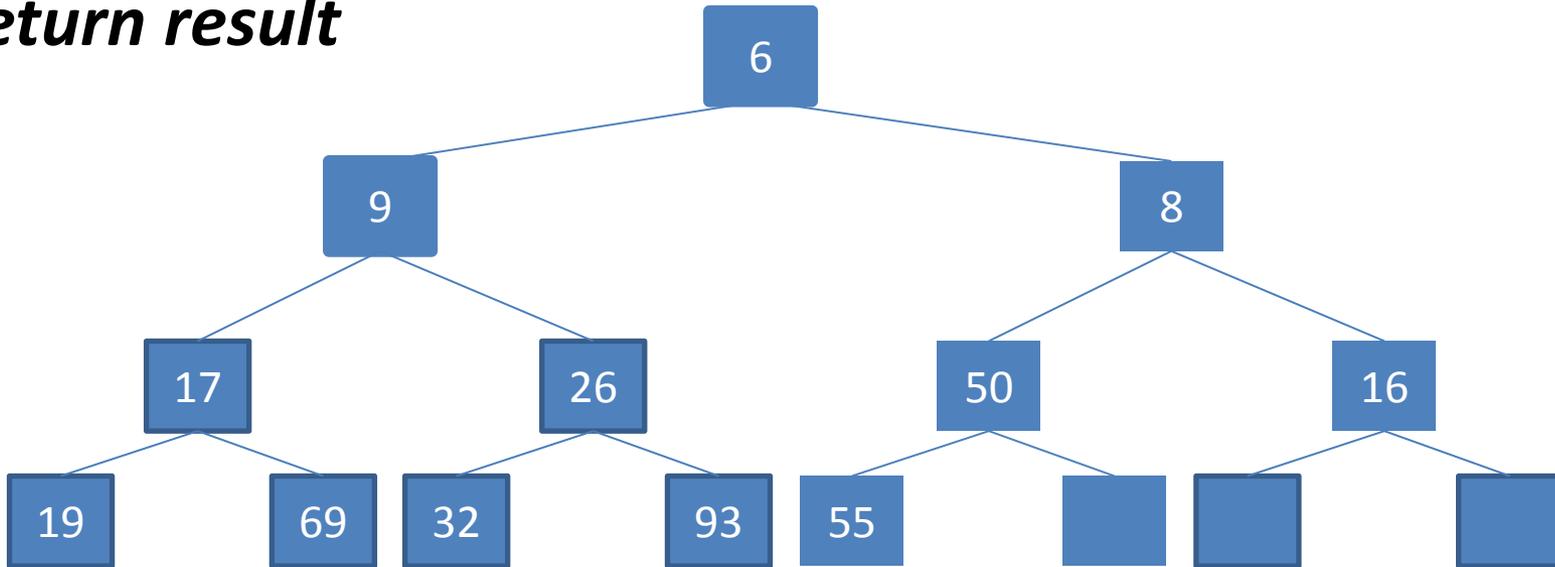
**	6	9	8	17	26	50	16	19	69	32	93	55			
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 12

capacity: 15



Return result



S

**	6	9	8	17	26	50	16	19	69	32	93	55			
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

last: 12

capacity: 15

Method	Time
size	$O(1)$
isEmpty	$O(1)$
insert	$O(\log(n))$
removeMin	$O(\log(n))$
min	$O(1)$

Have a look at priority queue as given in Java distribution

http://docs.oracle.com/javase/1.5.0/docs/api/java/util/PriorityQueue.html

PriorityQueue (Java 2 Platfo... x

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Overview Package **Class** Use Tree Deprecated Index Help

PREV CLASS NEXT CLASS

SUMMARY: NESTED | FIELD | [CONSTR](#) | [METHOD](#)

[FRAMES](#) [NO FRAMES](#) [All Classes](#)

DETAIL: FIELD | [CONSTR](#) | [METHOD](#)

Java™ 2 Platform
Standard Ed. 5.0

java.util

Class PriorityQueue<E>

[java.lang.Object](#)

- [java.util.AbstractCollection<E>](#)
- [java.util.AbstractQueue<E>](#)
- [java.util.PriorityQueue<E>](#)

Type Parameters:

E - the type of elements held in this collection

All Implemented Interfaces:

[Serializable](#), [Iterable<E>](#), [Collection<E>](#), [Queue<E>](#)

```
public class PriorityQueue<E>
extends AbstractQueue<E>
implements Serializable
```

An unbounded priority [queue](#) based on a priority heap. This queue orders elements according to an order specified at construction time, which is specified either according to their *natural order* (see [Comparable](#)), or according to a [Comparator](#), depending on which constructor is used. A priority queue does not permit null elements. A priority queue relying on natural ordering also does not permit insertion of non-comparable objects (doing so may result in [ClassCastException](#)).

The *head* of this queue is the *least* element with respect to the specified ordering. If multiple elements are tied for least value, the head is one of those elements -- ties are broken arbitrarily. The queue retrieval operations `poll`, `remove`, `peek`, and `element` access the element at the head of the queue.

A priority queue is unbounded, but has an internal *capacity* governing the size of an array used to store the elements on the queue. It is always at least as large as the queue size. As elements are added to a priority queue, its capacity grows automatically. The details of the growth policy are not specified.

18:03

... based on a priority heap

http://docs.oracle.com/javase/1.5.0/docs/api/java/util/PriorityQueue.html

PriorityQueue (Java 2 Platfo... x

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Overview Package **Class** Use Tree Deprecated Index Help

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Java™ 2 Platform
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java.util

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18:03

http://docs.oracle.com/javase/1.5.0/docs/api/java/util/PriorityQueue.html

PriorityQueue (Java 2 Platform)

Search java priority queue

PriorityQueue (SortedSet<E> extends E> o)
Creates a PriorityQueue containing the elements in the specified collection.

Method Summary

boolean	add(E o) Adds the specified element to this queue.
void	clear() Removes all elements from the priority queue.
Comparator<E> super E>	comparator() Returns the comparator used to order this collection, or null if this collection is sorted according to its elements natural ordering (using Comparable).
Iterator<E>	iterator() Returns an iterator over the elements in this queue.
boolean	offer(E o) Inserts the specified element into this priority queue.
E	peek() Retrieves, but does not remove, the head of this queue, returning null if this queue is empty.
E	poll() Retrieves and removes the head of this queue, or null if this queue is empty.
boolean	remove(Object o) Removes a single instance of the specified element from this queue, if it is present.
int	size() Returns the number of elements in this collection.

Methods inherited from class java.util.[AbstractQueue](#)

[addAll](#), [element](#), [remove](#)

18:04

Different method names ... add rather than insert



Your mission, should you choose to accept it ...

Exercise 4 (*assessed*)

- Implement the Heap
- Use it for sorting



FIN