

Types & Programming Languages

Exercises 4

These exercises are based on the material in Lectures 6, 7, 8 and 9.

- The typing rule for the equality test operator is

$$\frac{\Gamma \triangleright e : \text{int} \quad \Gamma \triangleright f : \text{int}}{\Gamma \triangleright e == f : \text{bool}}$$

It also makes sense to allow equality tests on boolean values, and we might try the general typing rule

$$\frac{\Gamma \triangleright e : T \quad \Gamma \triangleright f : T}{\Gamma \triangleright e == f : \text{bool}}$$

What problems are likely to be caused by this general form of equality test, if we have function types in the language? (Hint: think about how you would define reduction rules for the general equality test.)

- Work out the type of each of the following expressions, in the simply typed lambda calculus combined with SEL. (Using the notation `fn x:T => e` for $\lambda x : T.e$).

- `fn x:int => x+1`
- `fn x:int => (fn y:bool => (x == 1) & y)`
- `fn x:int->int => x(x(2))`
- `fn x:int->int => (fn y:int => xy)`

- Assuming that we have record types, what is the type of this expression?

`{ a = 1+2, b = 1==2, c = { x = 3 } }`

Show how this expression reduces to a value.

- Assume that we have sum types or variant types. A very useful idea is the *option type*. Given a type `T` we construct the type

`option T = < none:unit, some:T >`.

An expression of type `option T` is either the value `none()` or an expression `some(e)` where `e` is an expression of type `T`.

For example, `option int` has values `none()`, `some(1)`, `some(2)` etc. and we could also construct expressions such as `some(1+2)` and so on.

Option types can be used to represent the results of computations which might return an error condition instead of a useful value: for example, a square root function might be given the type

`sqrt : float -> option float`

to allow for the error case when the argument is negative.

- Give an example of a `case` expression which would be used to analyse an expression of type `option int`, and show how it reduces.
 - Write down the typing rule for `case` expressions associated with the type `option int`.
- We have seen how reference types allow us to describe the assignable variables of an imperative language: for example, the Java code

```
{ int x = 1;
  x = x + 1;
}
```

corresponds to

```
let val x = ref 1
in x := (!x) + 1
end
```

When we use objects in Java there is slightly more going on. A variable of some object type `MyClass` stores either `null` or a reference to an object of type `MyClass`. So a Java variable `MyClass x` corresponds to a value of type `option ref MyClass`, with the value `none()` representing `null`.

Assume that we have a Java function

```
MyClass f(int x).
```

A typical use of `f` might look like this:

```
MyClass x = f(1);
if (x == null)
  ... code 1 ...
else
  ... code 2 ...
```

Show how this style of coding would appear in our language.