



University  
of Glasgow

# *AspectJ for Runtime Behaviour Modification of Java Libraries*

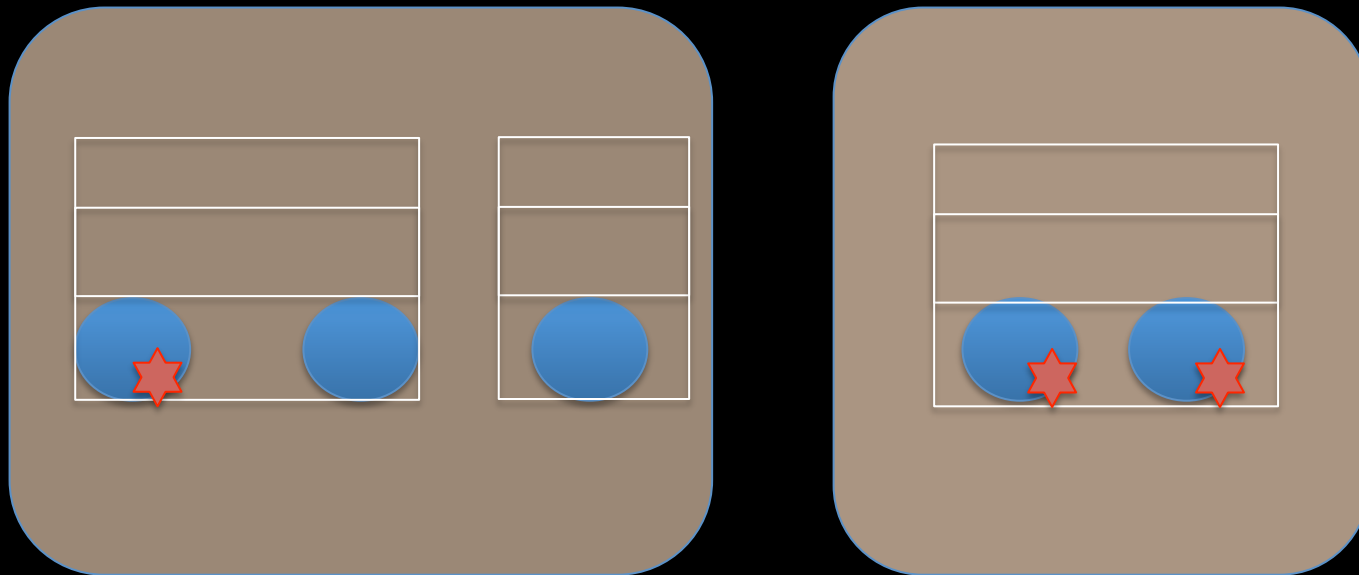
[Jeremy.Singer@glasgow.ac.uk](mailto:Jeremy.Singer@glasgow.ac.uk)

@jsinger\_compsci

1. What is **AspectJ**?
2. What are **approximate data structures**?
3. Using **AspectJ** to support **approx data structures**

1. What is **AspectJ**?

- Aspect-oriented programming enables clean modularization of cross-cutting concerns.



# Example Aspects

- Logging code
  - log each access to a resource
  - instrument each method entry point
- Resource management code
  - reference counting
  - implicit destruction

# AspectJ Syntax

- **Pointcut** – where to intercept execution
- **Advice** – what to compute

# Aspect Specification (1)

- **Pointcut** – where to intercept execution

```
public aspect HelloAspect {  
    pointcut greeting() :  
        HelloWorld.sayHello(..) ;  
}
```

## Aspect Specification (2)

- **Advice** – what to compute

```
public aspect HelloAspect {  
    after() returning: greeting() {  
        System.out.println("hello world");  
    }  
}
```



Is **AspectJ** a  
*domain-specific*  
*language?*

# Problems with **AspectJ**

- confusing syntax
  - what is  
`thisEnclosingJoinPointStaticPart?`
- evolving syntax
  - incremental changes to language, tracks Java releases
- poor and inconsistent documentation
  - ask David!

2. What are  
approximate data  
structures?

Conventionally, when you store a value at an indexed array location, you expect to retrieve that same value when you look up the same location later.

In approximate data structures, this 'contract' is probabilistic.

Approximate data structures are useful for dynamically adjusting application heap footprints based on resource availability.

Motivate – make `OutOfMemoryErrors`  
extinct

Possible application *genres* include:

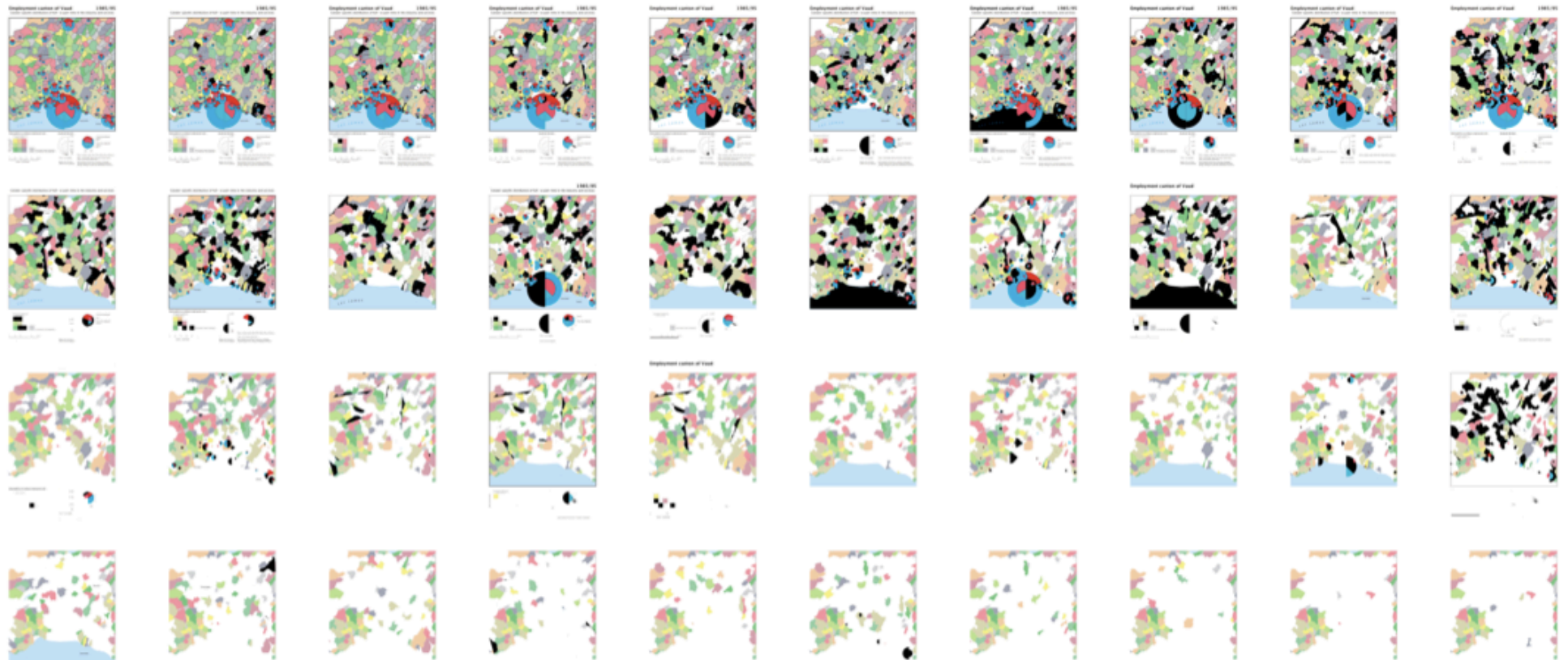
- big data processing
- long-running server applications

# Example (1)

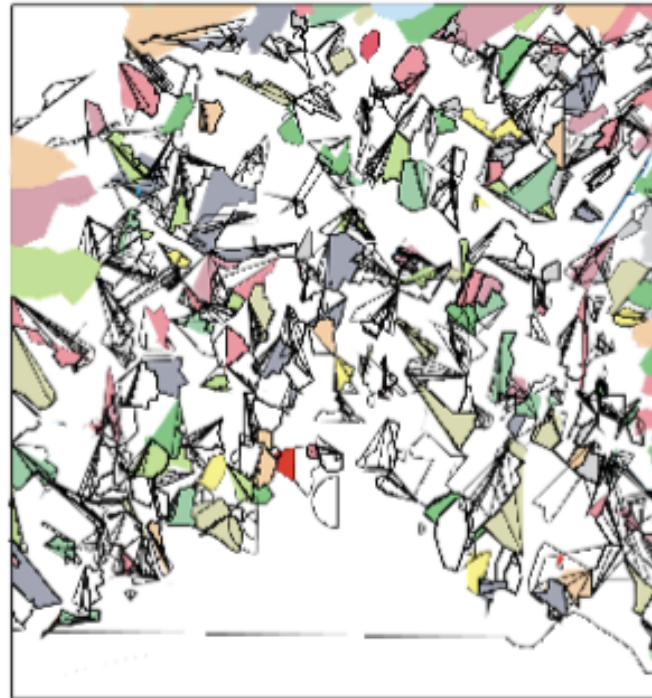
- the batik SVG processing tool



# Example (1)



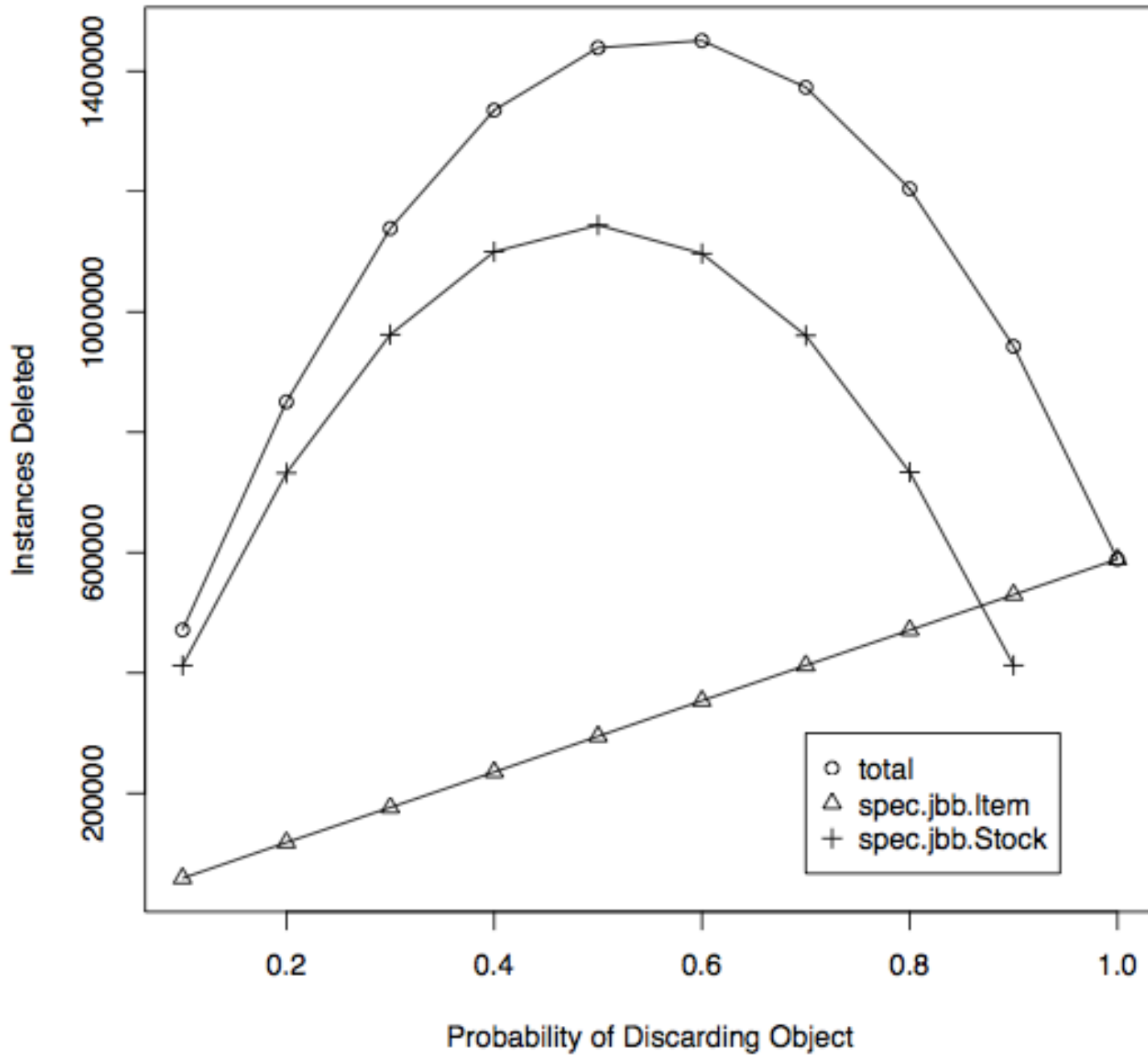
# Example (1)



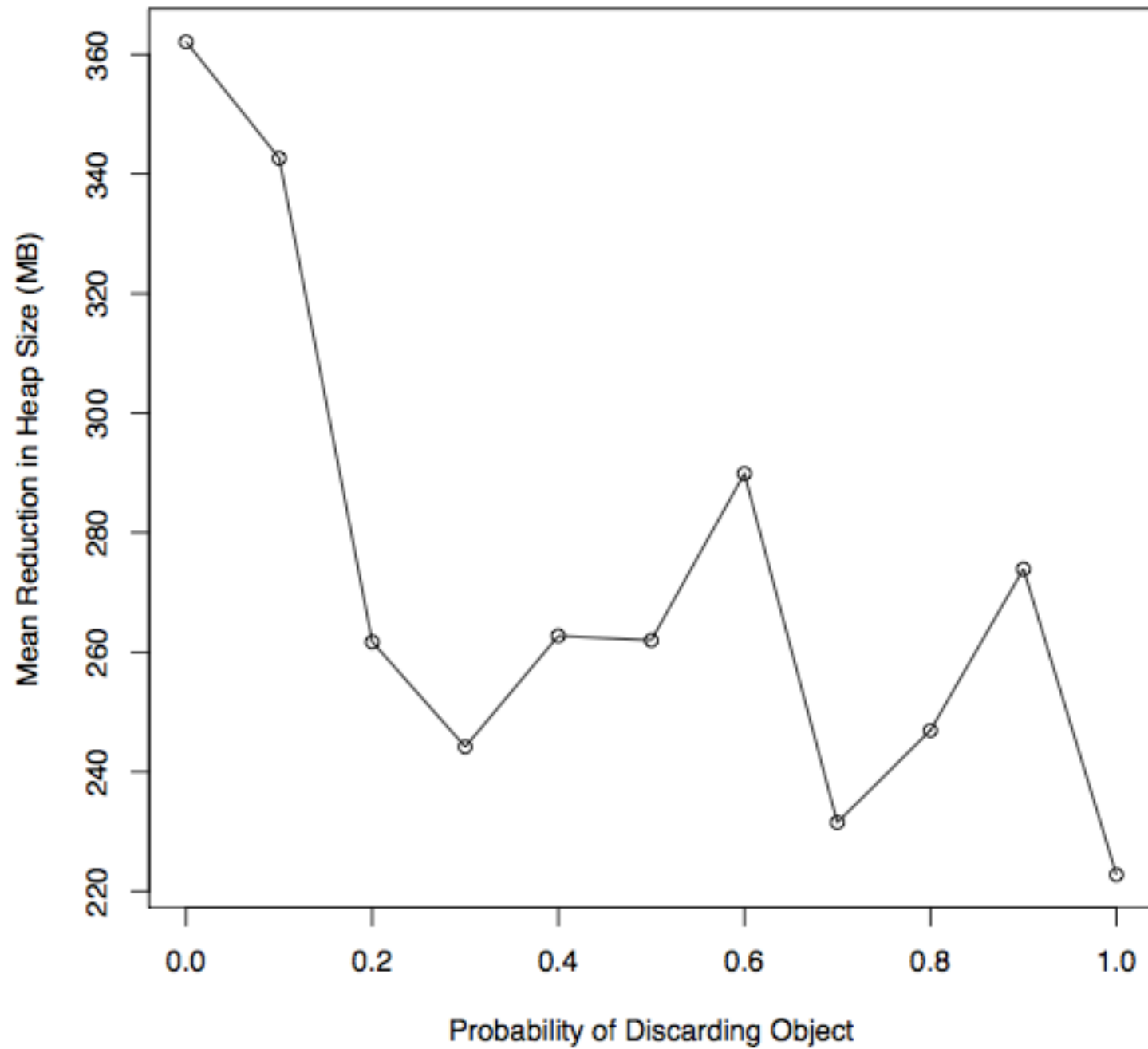
## Example (2)

- the PseudoJBB benchmark

### Deleted Objects by Type for the pjob Benchmark



Mean Reduction in Heap Size for pjbb2005



# Our motivation

- AnyScale applications adapt their computation (e.g. parallelism, precision) based on current platform resources. The application execution may migrate to a different platform as directed by cost/benefit calculations.
- [anyscale.org](https://anyscale.org)

3. Using **AspectJ** to  
support  
**approximate data  
structures**

# Pragmatics

- We instrument Java Collections classes to intercept data stores (e.g. `List.add`, `HashMap.put` and modify their semantics.
- probability of interception
- alternative action – drop or replace



```
1: procedure ADD TO A COLLECTION(collection, input)
2:   Increment insertionCounter for collection
3:   if collection is not empty then
4:     if insertionCounter > threshold then
5:       Set discard true with probability  $p_{drop}$ 
6:     end if
7:   end if
8:   if discard then
9:     for o in a permutation of collection do
10:      if type(o) == type(input) then
11:        Replace input with o
12:        Exit loop
13:      end if
14:    end for
15:   end if
16:   return proceed(collection, input)
17: end procedure
```

# Big Questions

1. Are approximate  
collections sensible?

2. How else can we forget data systematically?

3. Should the programmer, the compiler or the runtime be responsible for selecting data to forget?

4. Is AspectJ an appropriate language to support this concept?

