The Design of GUMSMP: a Multilevel Parallel Haskell Implementation

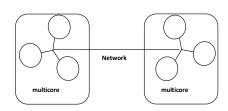
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Parallel Architectures

- Parallel architectures are increasingly multi-level e.g. clusters of multicores.
- A hybrid parallel programming model is often used to exploit parallelism across the cluster of multicores e.g. using MPI + OpenMP.
- Managing two abstractions is a burden for the programmer and increases the cost of porting to a new platform.
- The Main Goal: Providing efficient control of hierarchical architectures using GpH.



GpH(Glasgow Parallel Haskell)

- Semi-explicit parallel Haskell.
- Parallelism is expressed by two primitives added to the Haskell program: par and pseq.
- Example:

GpH(Glasgow Parallel Haskell)

- Evaluation strategies: polymorphic and higher order functions controlling parallelism.
- Potentially add extensions to refine placement e.g parBound.
- Tow main implementations :
 - GHC-SMP shared memory.
 - GHC-GUM distributed memory.

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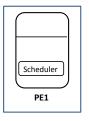
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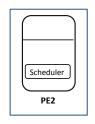
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 - Provides a scalable model.
 - Efficient exploitation of the the specifics of distributed and shared memory on different levels of the hierarchy.
 - Provides a single high-level programming model.

GUMSMP Design Overview

- Memory Management: the same virtual shared heap as GHC-GUM.
- Communication: the same mechanism implemented in GHC-GUM.
- Load Balancing: the combination of GHC-SMP and GHC-GUM mechanisms (using the hierarchy-aware policy).

- Searching for Local Work.
- Searching for Remote Work.

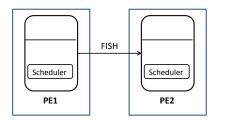






PE1 needs work

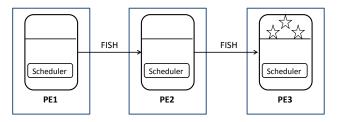
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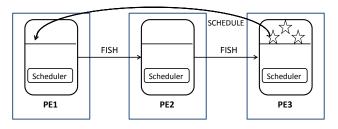
PE1 sends fish message

- Searching for Local Work.
- 2 Searching for Remote Work.



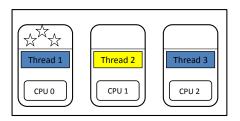
PE2 forwards the message

- Searching for Local Work.
- 2 Searching for Remote Work.



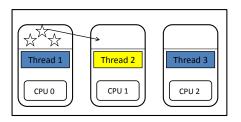
PE3 sends work to PE1

- Processor's Spark Pool is implemented as a bounded work-stealing queue.
- The owner can push and pop from one end of the queue without synchronization.
- Other threads can steal from the other end of the queue.



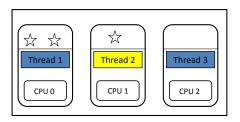
PE1 creates 'spark thread' to get work

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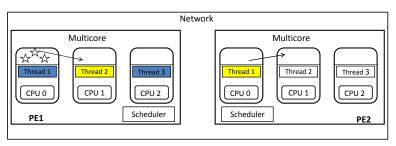
The 'spark thread' steals spark

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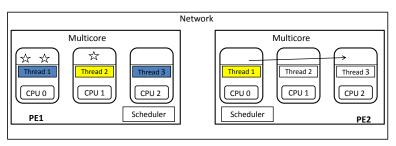
Evaluates the spark

- Work distribution of GUMSMP is hierarchy aware.
- It uses a work-stealing algorithm, through sending FISH message, on networks (inherited from GHC-GUM).
- Within a multicore it will search for a spark by directly accessing spark pools (inherited from GHC-SMP).



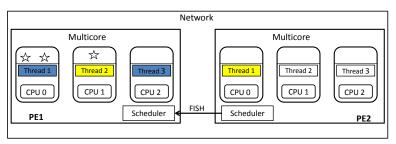
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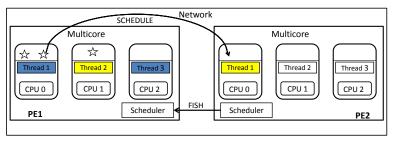
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PE2 sends Fish message

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PE1 replay with spark

GUMSMP Design Objectives

 Hierarchy aware load balancing Important to maintain even load distribution, but accept imbalances as the communication cost increases.

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- Hierarchy aware load balancing Important to maintain even load distribution, but accept imbalances as the communication cost increases.
- Mostly passive load distribution Essential to maintain passive load distribution, but switch to active in some cases e.g high-watermark.
- Effective latency hiding The system must be designed so that communication cost is not in the critical path of cooperating computations.

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- 2 Separate spark pool, dedicated to imported sparks.
 - (+) Keep related piece of work together.
 - (+) Useful in some situation e.g no idle processors any more.
 - (-) Requires additional stealing step.

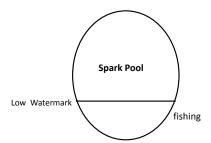
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- 2 Low-Watermark mechanism.



 $\label{process} \textbf{Work-offloading:} \ \ \text{How to process the received work-requesting message} \ ?$

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- Select spark from the processor with largest spark pool.
 - (-) Impose additional overheads.
- Random.

Ongoing Work

- performance within 7% of the original GHC-SMP implementation.
- Complete the implementation of the enhanced work distribution policy.
- Assess the quality of the enhanced work distribution policy on hierarchical architectures

• The shared memory component of the hybrid system shows

Conclusion

- The design of the new multi-level parallel Haskell implementation GUMSMPis presented.
- Designed for high-performance computation on multilevel architectures
 e.g. networks of multi-cores.
- The design focuses on flexible work distribution policies.
 - Hierarchy aware load balancing.
 - Mostly passive load distribution.
 - Effective latency hiding.
- The main benefits:
 - scalable model.
 - efficient exploitation of distributed and shared memory on different levels of the hierarchy.
 - single programming model.

Thank You..