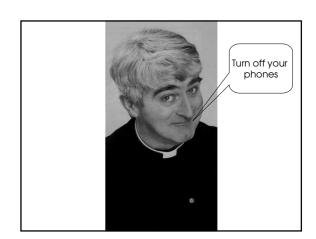
AF2



Primes, gcd, some examples, reading

- p>1 is prime if the only positive factors are 1 and p if p is not prime it is composite

The Fundamental Theorem of Arithmetic

Every positive integer can be expressed as a unique product of primes



$$n = \prod_{i=1}^{\kappa} p_i^{e_i}$$

How many prime numbers are there? Is there a largest prime number?

- assume we know all the prime numbers
 - · let p be the largest prime
 - · multiply all the primes together
 - · add 1
 - · call this n
- n is not divisible by 2, we get remainder 1
 n is not divisible by 3, we get remainder 1
 n is not divisible by 5, we get remainder 1

- n is not divisible by p, we get remainder 1
- But, by the FTA we know n has a prime factorisation
- Therefore n must have a prime divisor > p
- Therefore there is no greatest prime

Due to Euclid

PRIMES

If n is composite n has a prime divisor $\leq \sqrt{n}$

$$composite(n) \rightarrow n = a.b$$

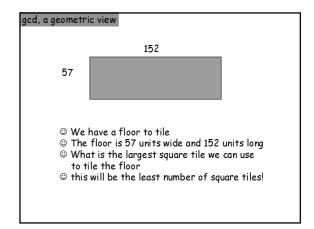
$$a \le \sqrt{n} \lor b \le \sqrt{n}$$

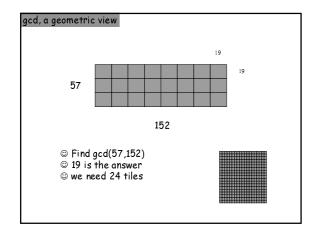
$$\neg (a > \sqrt{n} \land b > \sqrt{n})$$

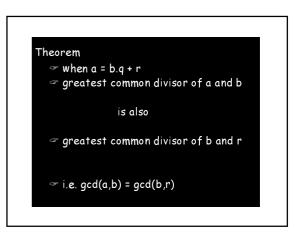
Therefore, the divisor α or b is either prime or due to the $% \left\{ 1,2,...,n\right\}$ fundamental theorem of arithmetic, can be expressed as a product of primes

∴ n has a divisor $\leq \sqrt{n}$

for RSA encryption "The most efficient factorisation methods (1999) require billions of years to factor 400 digit integers. Consequently ... when 200 digit primes ... messages encrypted ... cannot be found in a reasonable time." "When new factorisation techniques are found ... it will be necessary to use larger primes to ensure secrecy/security of messages."







A proof follows The proof is based on what we already know about division RTP: gcd(a,b) = gcd(b,r) where a = b.q + r

Sketch of the proof
Show that common divisors of a and b are also common divisors of b and r

(1) show that if some d|a and d|b then d|r
(2) show that if some d|b and d|r then d|a
(3) conclude that a common divisor of a and b is also a common divisor of b and r
(4) consequently gcd(a,b) = gcd(b,r)

RTP: gcd(a,b) = gcd(b,r) where a = b.q + r

(1) show that if some d|a and d|b then d|r

a = b.q + r ∴a - b.q = r

We have already proved that

- if d|b then d|b.q if d|a and d|b.q then d|(a b.q)

since d|a and d|b.q it follows that d|(a - b.q) since (a - b.q) = r it follows that d|r

Consequently a common divisor of a and b also divides r

RTP: gcd(a,b) = gcd(b,r) where a = b.g + r

(2) show that if some d|b and d|r then d|a

a = b.q + r

We have already proved that

- if d|b then d|b.qif d|b.q and d|r then d|(b.q + r)

since d|b,q and d|r it follows that d|(b,q + r) since (b,q + r) = a it follows that d|a

Consequently a common divisor of b and r also divides a

RTP: gcd(a,b) = gcd(b,r) where a = b.q + r

(3) conclude that a common divisor of a and b is also a common divisor of b and r

From (1) and (2) we have proved that

- 🖈 a common divisor of a and b is also a divisor of r
- 🕏 consequently a common divisor of a and b is also a common divisor of b and r

RTP: gcd(a,b) = gcd(b,r) where a = b.q + r

(4) consequently gcd(a,b) = gcd(b,r)

From (3) we can conclude that the greatest common divisor of a and b is also a greatest common divisor of b and r

This suggests an algorithm

- given a and bif a = 0 then b is the gcd
- if b = 0 then a is the gcd
- otherwise
 - a ÷b = q remainder r
 - we have established that we can substitute
 - gcd(a,b) with gcd(b,r)

 - now compute gcd(b,r)
 note, r is decreasing!

- · We have shown that when a = b.q + r
- any divisor of a and b is also a divisor of b and r
 consequently gcd(a,b) = gcd(b,r)
 We state that gcd(0,b) = b and gcd(a,0) = a

- · therefore, there is nothing to do if a or b is zero
- Due to our theorem
- we can reduce gcd(a,b) to a simpler problem gcd(b,r)
- · at each iteration we replace
 - b with r where $0 \le r < b$
 - · at each iteration b decreases
 - eventually this must terminate with b = 0
 - · and as we said, we are then done
 - · and a is the answer!

```
gcd(a,b)

• gcd(a,b)

• if b = 0

• then a

• else gcd(b,a mod b)

• (0) gcd(a,b)

• (1) if b = 0 goto 6

• (2) let r := a mod b

• (3) a := b

• (4) b := r

• (5) goto (1)

• (6) return(a)

gcd(a,b)

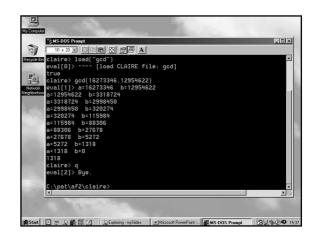
while b ≠ 0

do begin

r := a mod b;

a := b;

b := r;
end;
```



Numbers P Can you do arithmetic in different bases? P I expect so (see pages 169-177) Can you do addition, multiplication, subtraction, division P in different bases? (I expect so) Do you know the algorithm for this? P Again, I expect so.