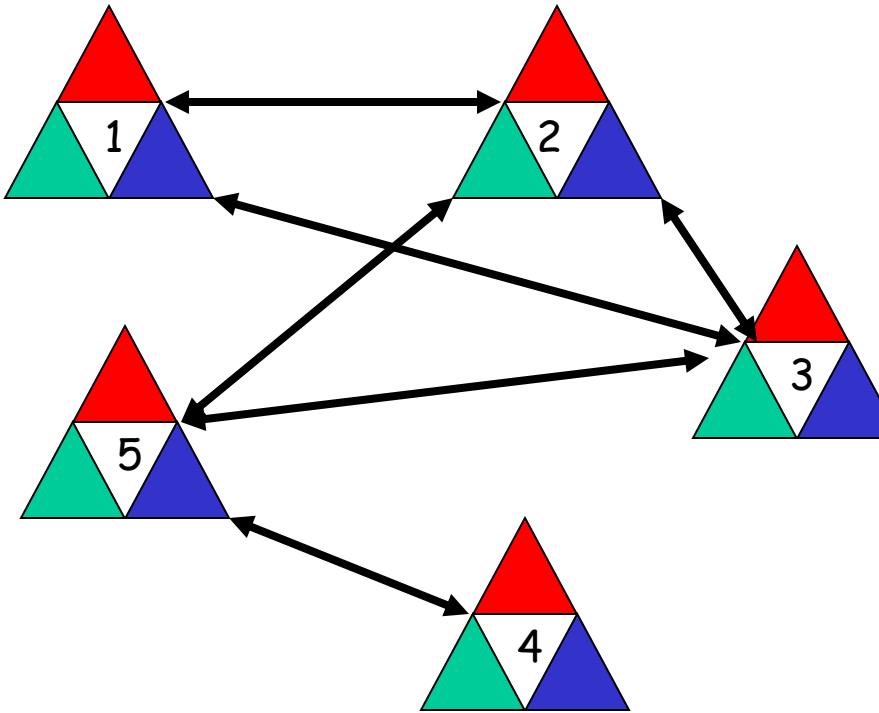


Chronological Backtracking (BT)

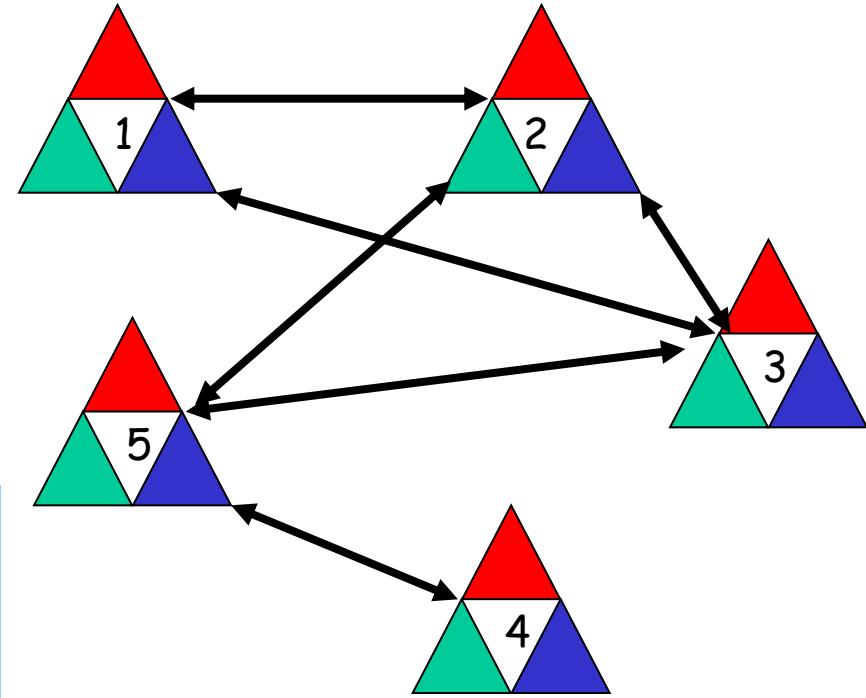
An example problem



Colour each of the 5 nodes, such that if they are adjacent, they take different colours

Representation (csp1)

- variables v_1, v_2, v_3, v_4 , and v_5
- domains d_1, d_2, d_3, d_4 , and d_5
- domains are the three colours {R,B,G}
- constraints
 - $v_1 \neq v_2$
 - $v_1 \neq v_3$
 - $v_2 \neq v_3$
 - $v_2 \neq v_5$
 - $v_3 \neq v_5$
 - $v_4 \neq v_5$



Assign a value to each variable, from its domain, such that the constraints are satisfied

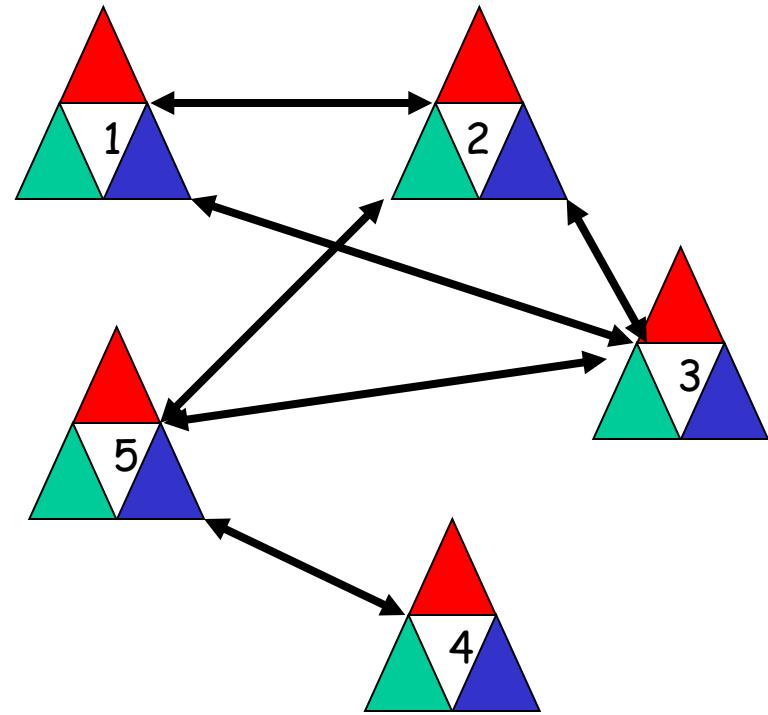
$$CSP = (V, D, C)$$

Chronological Backtracking (BT)

As a pair of mutually recursive functions

bt-label

```
bt-label(i,v,d,cd,n)
begin
  if i > n
    then return "solution"
  else begin
    consistent := false;
    for v[i] ∈ cd[i] while ¬consistent
      do begin
        consistent := true;
        for h := 1 to i-1 while consistent
          do consistent := check(v,i,h);
        if ¬consistent
          then cd[i] := cd[i] \ v[i];
        end
        if consistent
          then bt-label(i+1,v,d,cd,n)
        else bt-unlabel(i,v,d,cd,n)
      end
    end
```



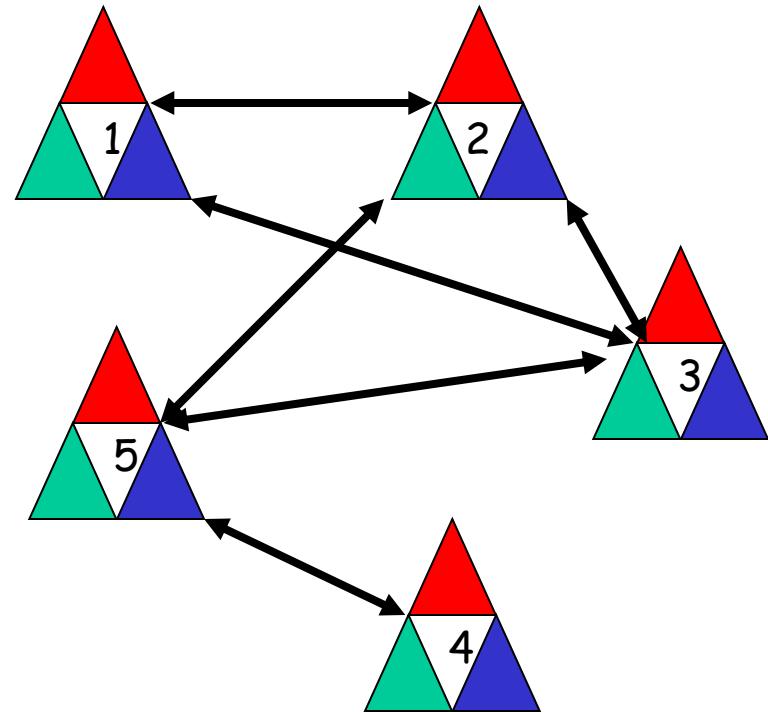
i the index of the current variable
h the index of a past variable (local)
v an array of variables to be instantiated
d an array of domains
cd an array of current domains
n the number of variables

constraints are binary.
consistent is boolean (local)
 $bt(v,d,n) = \text{for } i := 1 \text{ to } n \text{ } cd[i] := d[i];$
 $\text{return } bt\text{-label}(1,v,d,cd,n)$

Chronological Backtracking (BT)

bt-label

```
bt-label(i,v,d,cd,n)
begin
  if i > n
  then return "solution"
  else begin
    consistent := false;
    for v[i] ∈ cd[i] while ¬consistent
    do begin
      consistent := true;
      for h := 1 to i-1 while consistent
      do consistent := check(v,i,h);
      if ¬consistent
      then cd[i] := cd[i] \ v[i];
    end
    if consistent
    then bt-label(i+1,v,d,cd,n)
    else bt-unlabel(i,v,d,cd,n)
  end
```

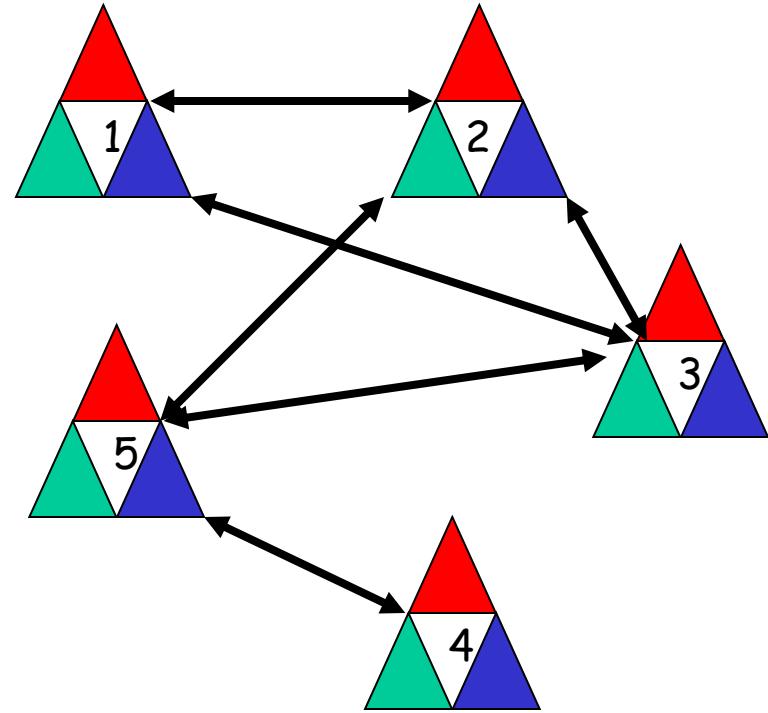


Find a consistent instantiation for $v[i]$

Chronological Backtracking (BT)

bt-label

```
bt-label(i,v,d,cd,n)
begin
  if i > n
  then return "solution"
  else begin
    consistent := false;
    for v[i] ∈ cd[i] while ¬consistent
    do begin
      consistent := true;
      for h := 1 to i-1 while consistent
      do consistent := check(v,i,h);
      if ¬consistent
        then cd[i] := cd[i] \ v[i];
        end
      if consistent
        then bt-label(i+1,v,d,cd,n)
        else bt-unlabel(i,v,d,cd,n)
    end
```

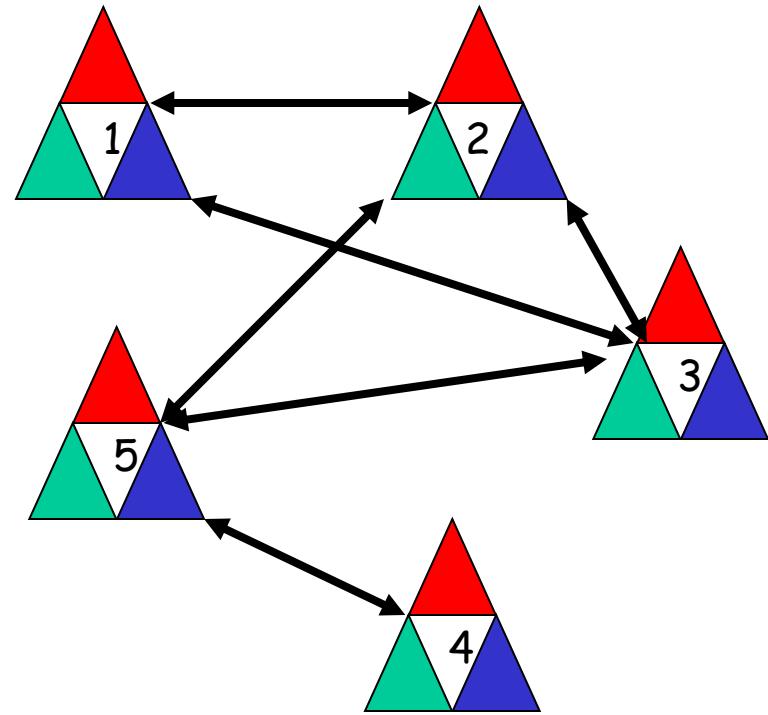


check backwards, from current to past

Chronological Backtracking (BT)

bt-label

```
bt-label(i,v,d,cd,n)
begin
  if i > n
    then return "solution"
  else begin
    consistent := false;
    for v[i] ∈ cd[i] while ¬consistent
      do begin
        consistent := true;
        for h := 1 to i-1 while consistent
          do consistent := check(v,i,h);
        if ¬consistent
          then cd[i] := cd[i] \ v[i];
        end
      if consistent
        then bt-label(i+1,v,d,cd,n)
        else bt-unlabel(i,v,d,cd,n)
  end
```

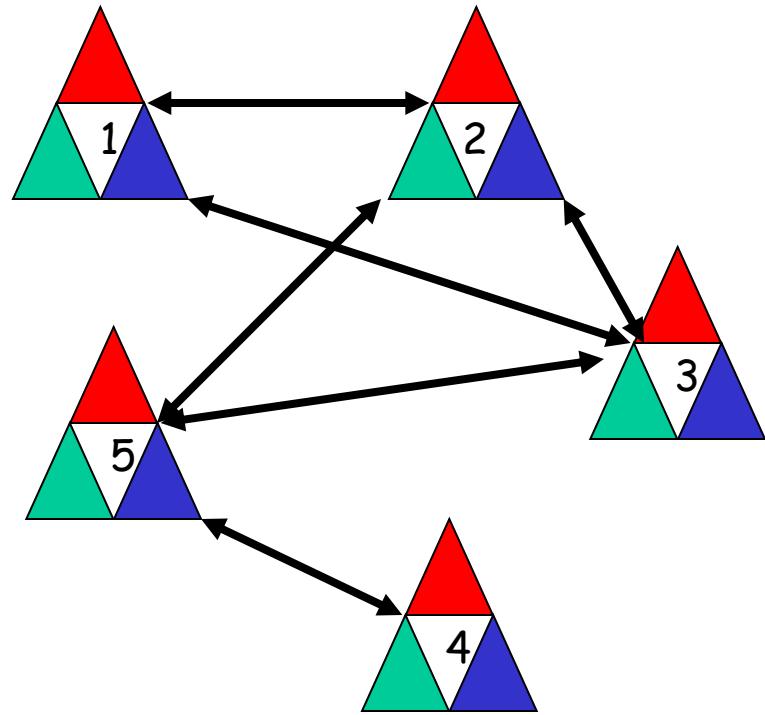


reurse

Chronological Backtracking (BT)

bt-unlabel

```
bt-unlabel(i,v,d,cd,n)
begin
  if i = 0
    then return "fail"
  else begin
    h := i - 1;
    cd[h] := cd[h] \ v[h];
    cd[i] := d[i];
    if cd[h] ≠ nil
      then bt-label(h,v,d,cd,n)
    else bt-unlabel(h,v,d,cd,n)
  end
end
```

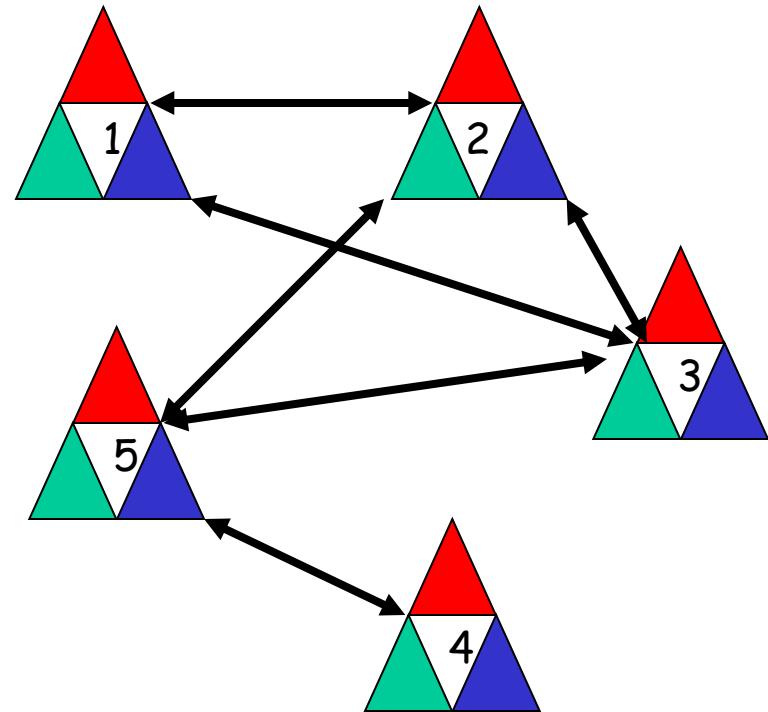


Chronological Backtracking (BT)

bt-unlabel

```
bt-unlabel(i,v,d,cd,n)
begin
  if i = 0
    then return "fail"
  else begin
    h := i - 1;
    cd[h] := cd[h] \ v[h];
    cd[i] := d[i];
    if cd[h] ≠ nil
      then bt-label(h,v,d,cd,n)
    else bt-unlabel(h,v,d,cd,n)
  end
end
```

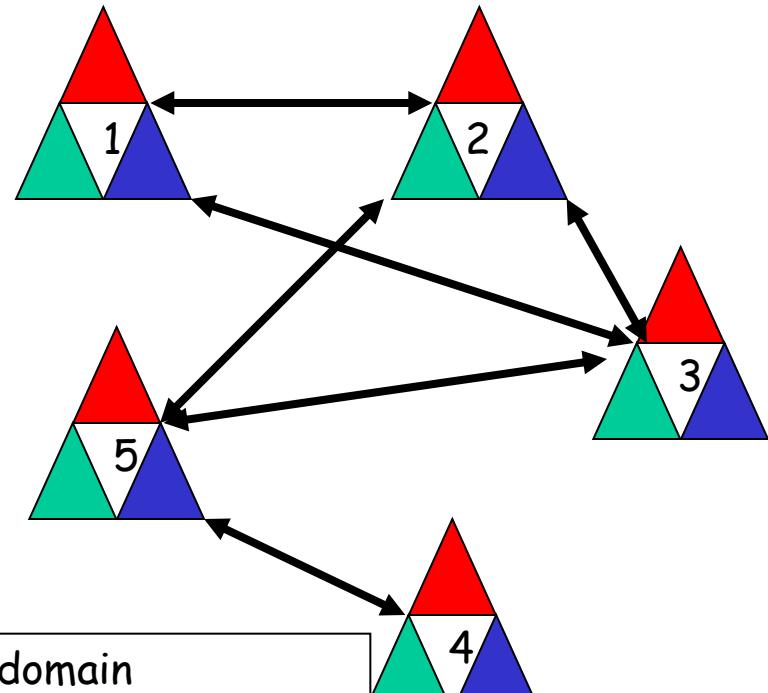
Past variable



Chronological Backtracking (BT)

bt-unlabel

```
bt-unlabel(i,v,d,cd,n)
begin
  if i = 0
    then return "fail"
  else begin
    h := i - 1;
    cd[h] := cd[h] \ v[h];
    cd[i] := d[i];
    if cd[h] ≠ nil
      then bt-label(h,v,d,cd,n)
    else bt-unlabel(h,v,d,cd,n)
  end
end
```

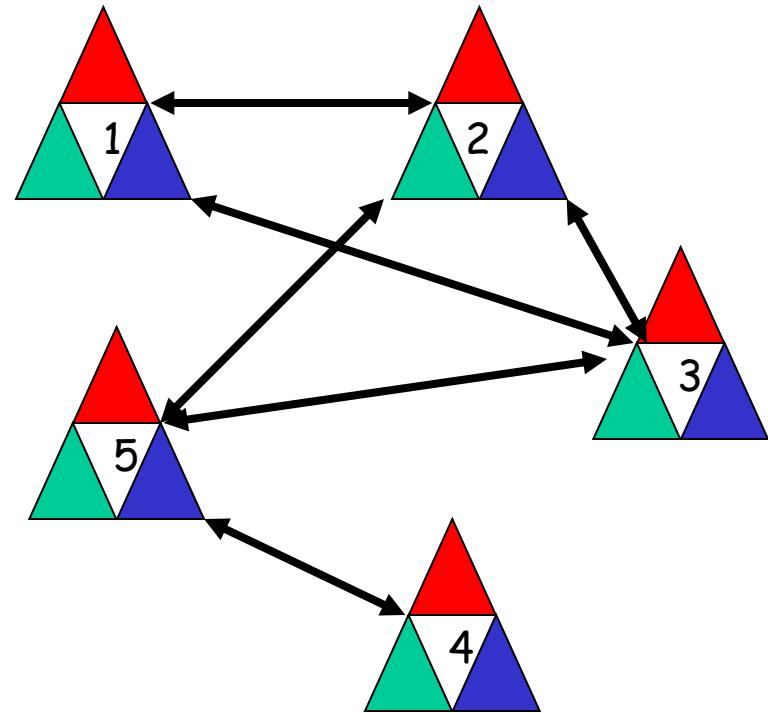


Chronological Backtracking (BT)

bt-unlabel

```
bt-unlabel(i,v,d,cd,n)
begin
  if i = 0
    then return "fail"
  else begin
    h := i - 1;
    cd[h] := cd[h] \ v[h];
    cd[i] := d[i];
    if cd[h] ≠ nil
      then bt-label(h,v,d,cd,n)
    else bt-unlabel(h,v,d,cd,n)
  end
end
```

recurse



Iterative implementation of BT

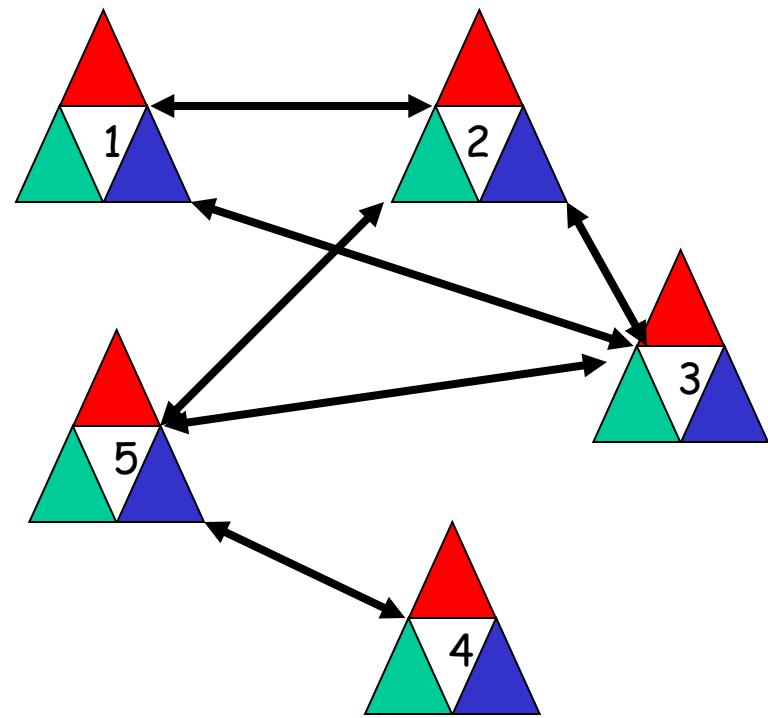
```
search(n,status)
begin
    consistent := true;
    status := "unknown";
    i := 1;
    while status = "unknown"
        do begin
            if consistent
                then i := label(i,consistent)
            else i := unlabel(i,consistent);
            if i > n
                then status = "solution"
            else if i = 0
                then status := "impossible"
            end
        end
```

```
bt-label(i,consistent)
begin
    consistent := false;
    for v[i] in cd[i] while ¬consistent
        do begin
            consistent := true;
            for h := 1 to i-1 while consistent
                do consistent := check(v,i,h);
            if ¬consistent
                then cd[i] := cd[i] \ v[i];
            end;
            if consistent then return i+1 else return I
        end
```

```
bt-unlabel(i,consistent)
begin
    h := i-1;
    cd[i] := d[i];
    cd[h] := cd[h] \ v[h];
    consistent := cd[h] ≠ nil
    return h;
end
```

This is more realistic. Why?

Chronological Backtracking (BT)



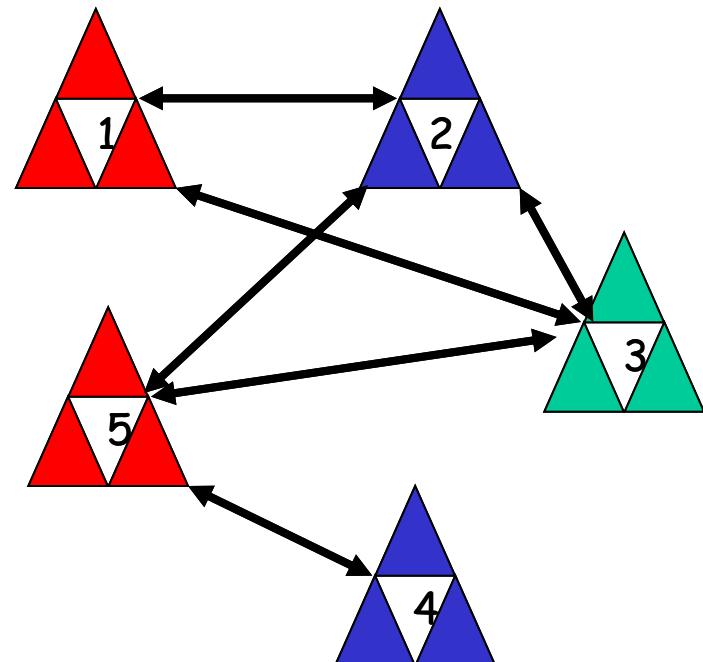
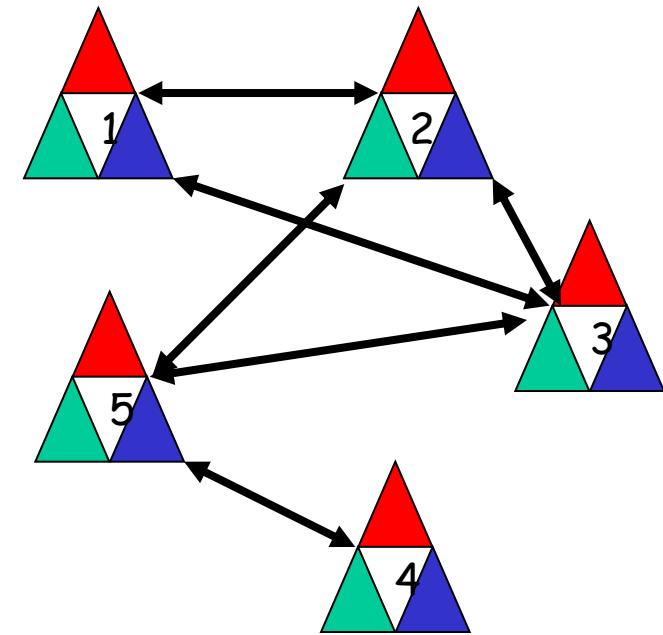
Three Different Views of Search

a trace
a tree
past, current, future

A Trace of BT (assume domain ordered {R,B,G})

- $V[1] := R$
- $V[2] := R$
 - $\text{check}(v[1], v[2])$ fails
- $V[2] := B$
 - $\text{check}(v[1], v[2])$ good
- $V[3] := R$
 - $\text{check}(v[1], v[3])$ fails
- $V[3] := B$
 - $\text{check}(v[1], v[3])$ good
 - $\text{check}(v[2], v[3])$ fails
- $V[3] := G$
 - $\text{check}(v[1], v[3])$ good
 - $\text{check}(v[2], v[3])$ good
- $V[4] := R$
- $V[5] := R$
 - $\text{check}(v[2], v[5])$ good
 - $\text{check}(v[3], v[5])$ good
 - $\text{check}(v[4], v[5])$ fails
- $V[5] := B$
 - $\text{check}(v[2], v[5])$ fails
- $V[5] := G$
 - $\text{check}(v[2], v[5])$ good
 - $\text{check}(v[3], v[5])$ fails

- backtrack!
- $V[4] := B$
- $V[5] := R$
 - $\text{check}(v[2], v[5])$ good
 - $\text{check}(v[3], v[5])$ good
 - $\text{check}(v[4], v[5])$ good
- solution found



16 checks and 12 nodes

A Tree Trace of BT (assume domain ordered {R,B,G})



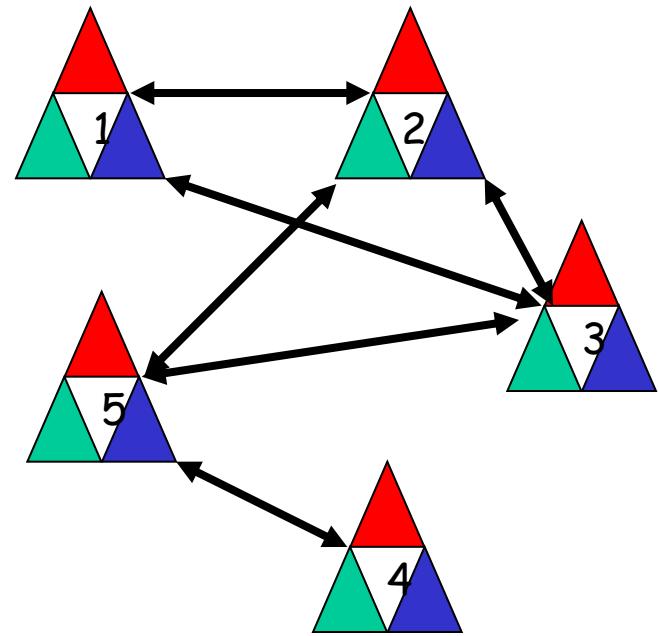
v1

v2

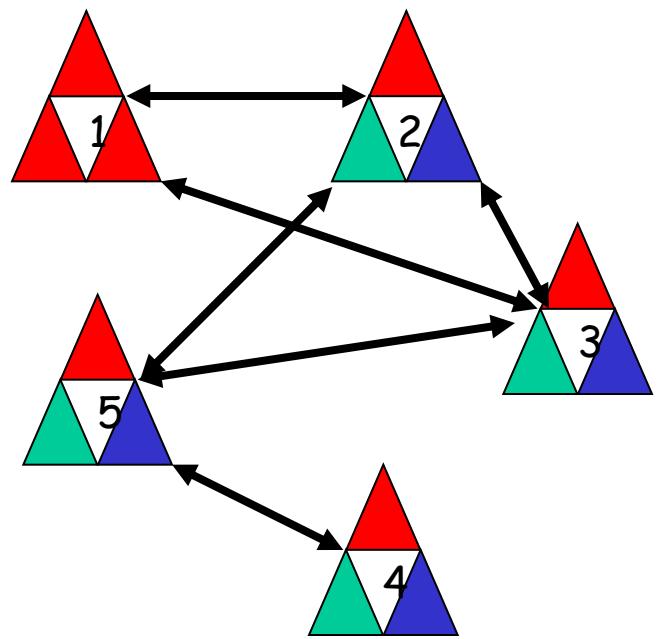
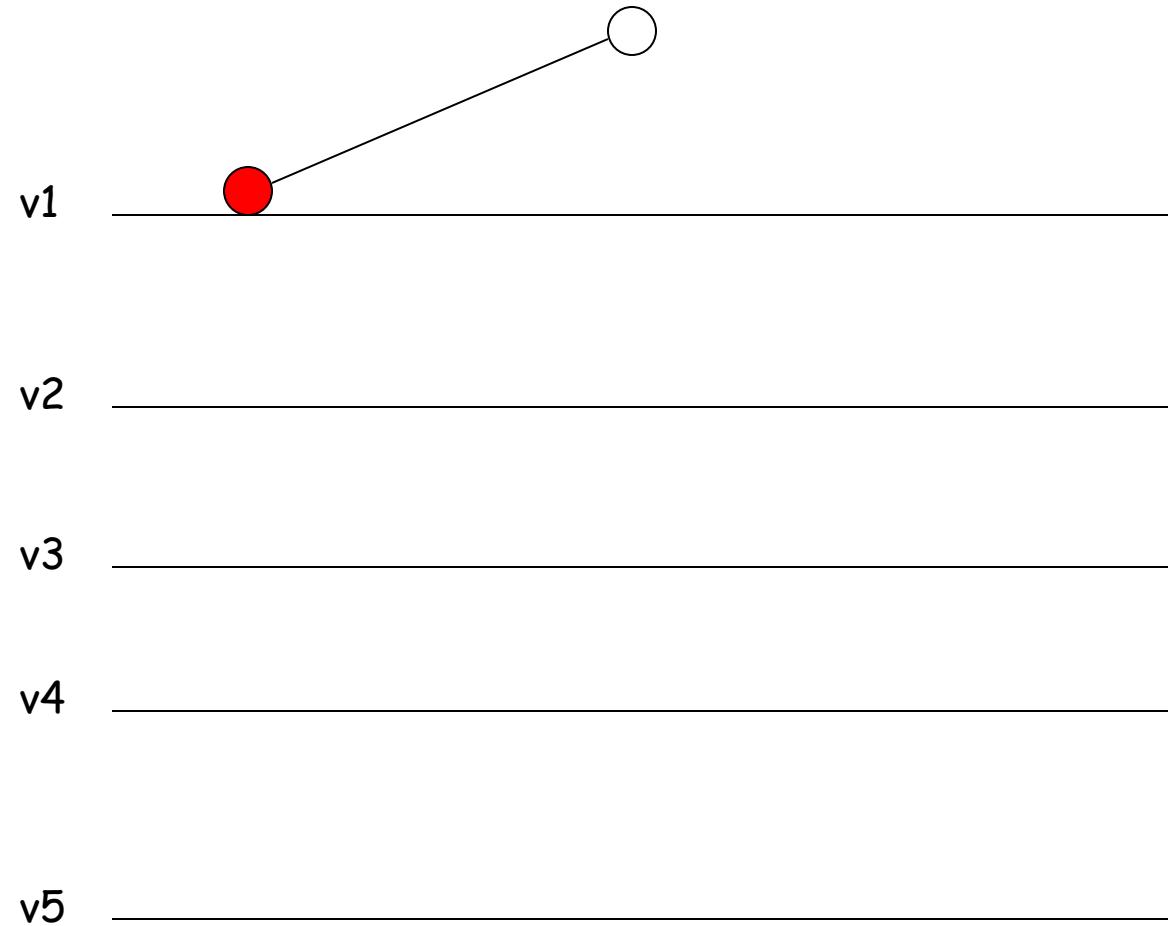
v3

v4

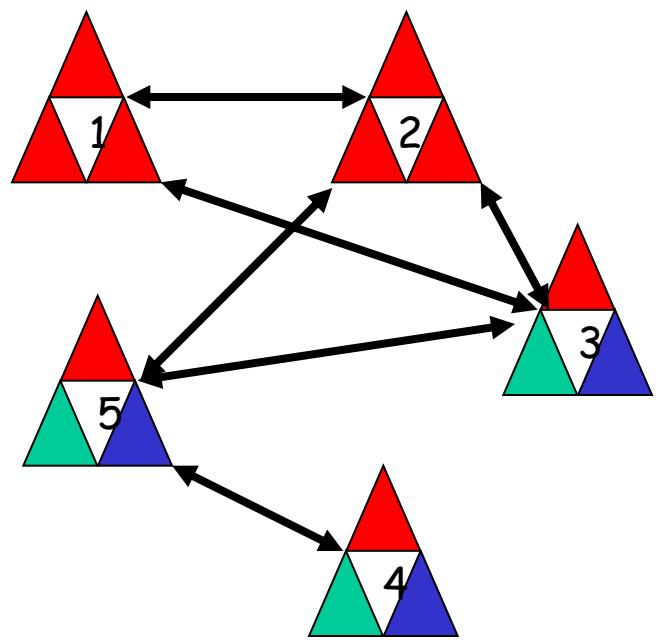
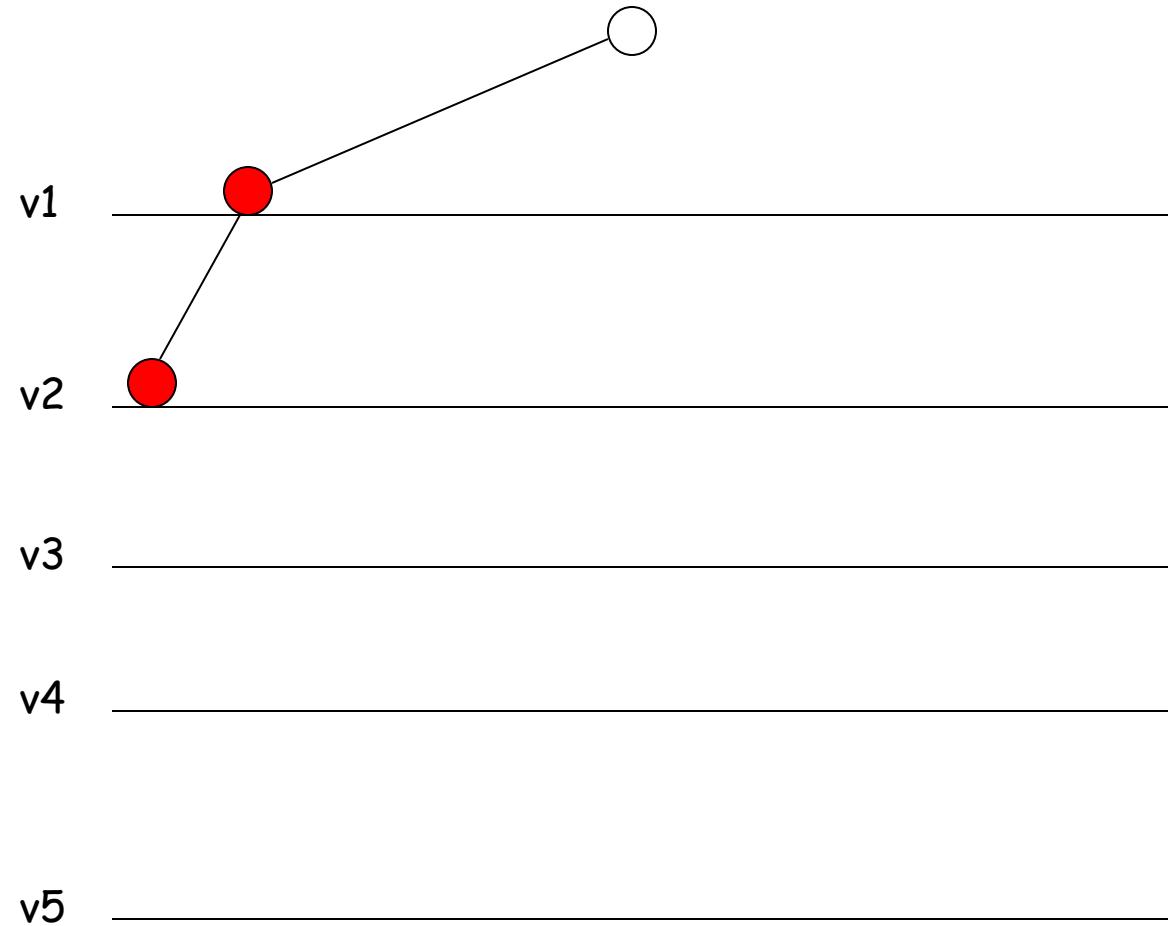
v5



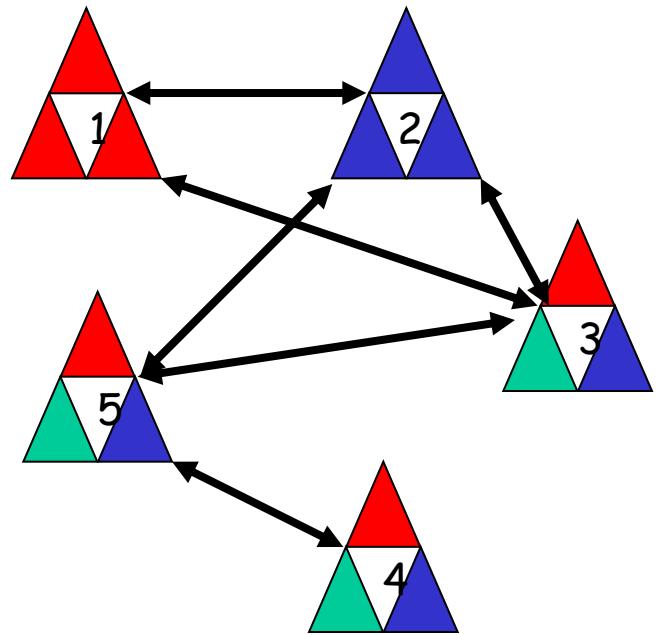
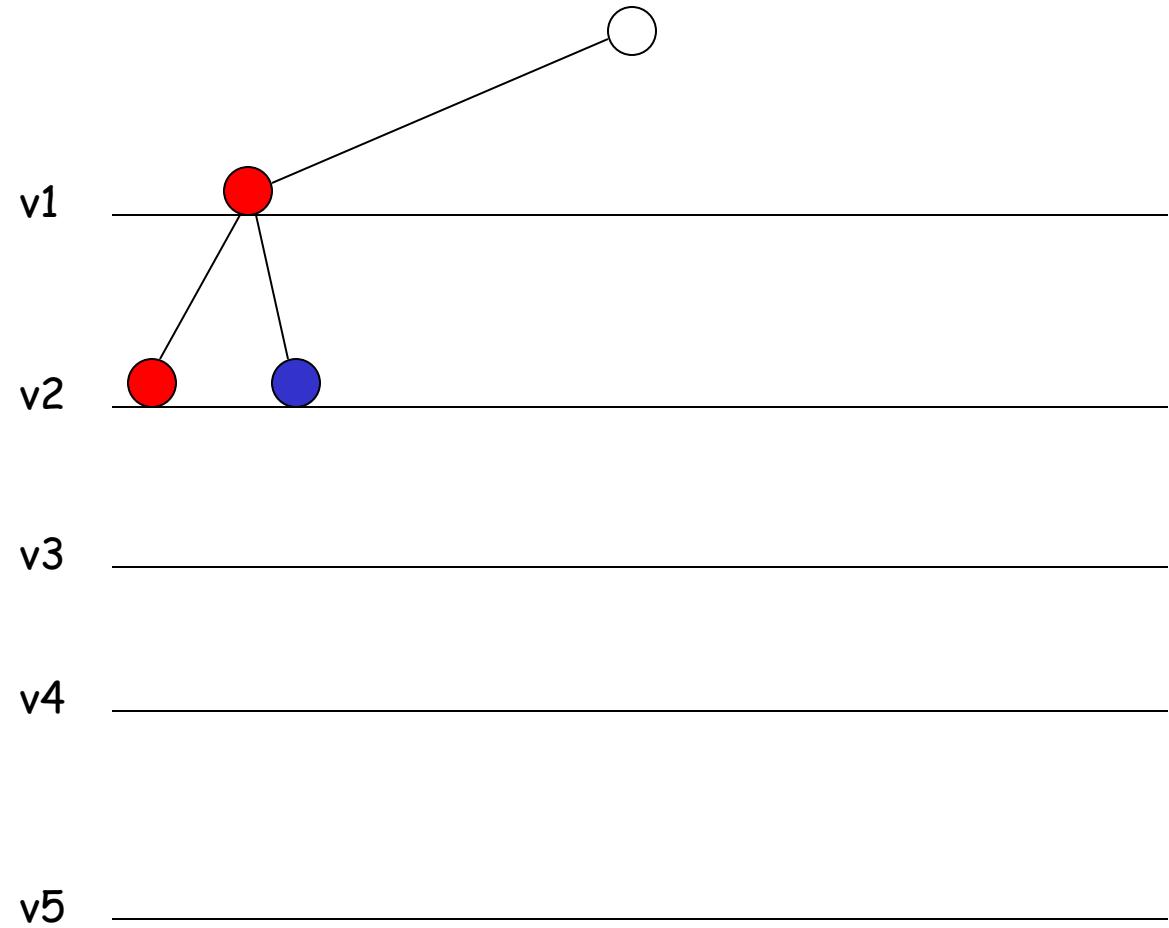
A Tree Trace of BT (assume domain ordered {R,B,G})



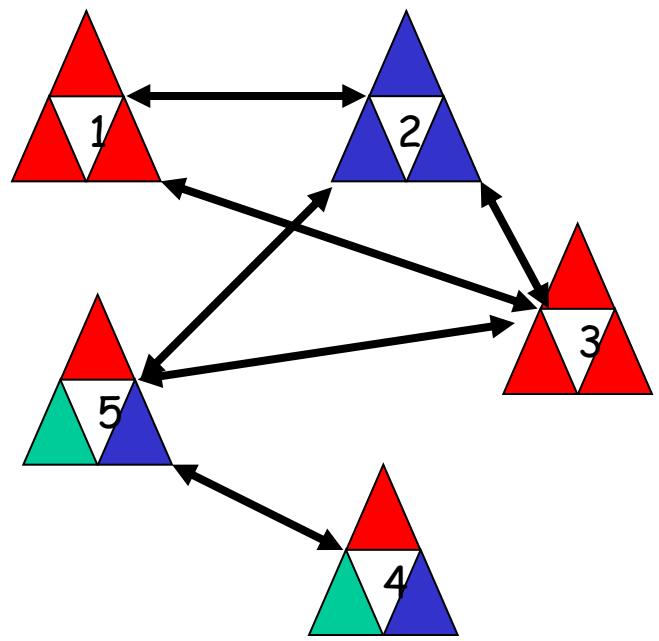
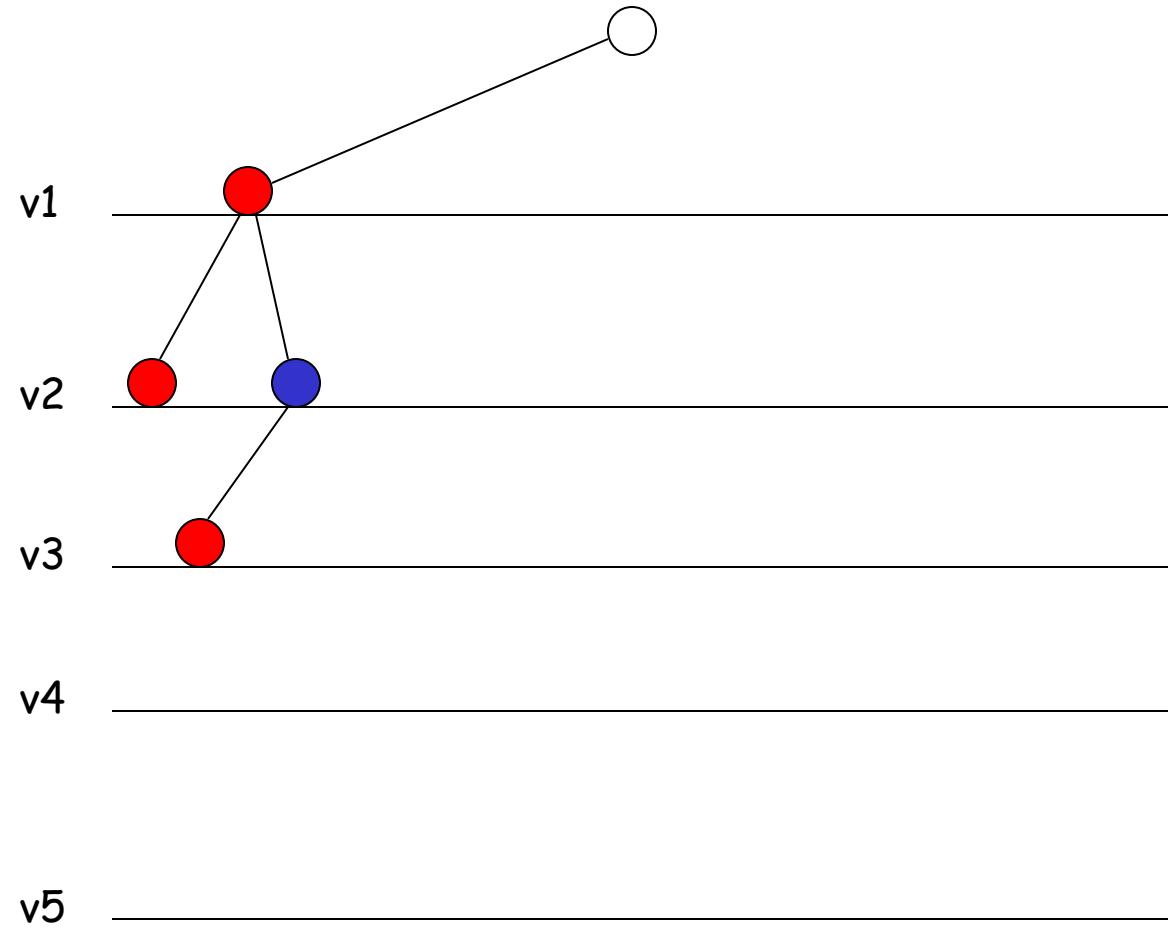
A Tree Trace of BT (assume domain ordered {R,B,G})



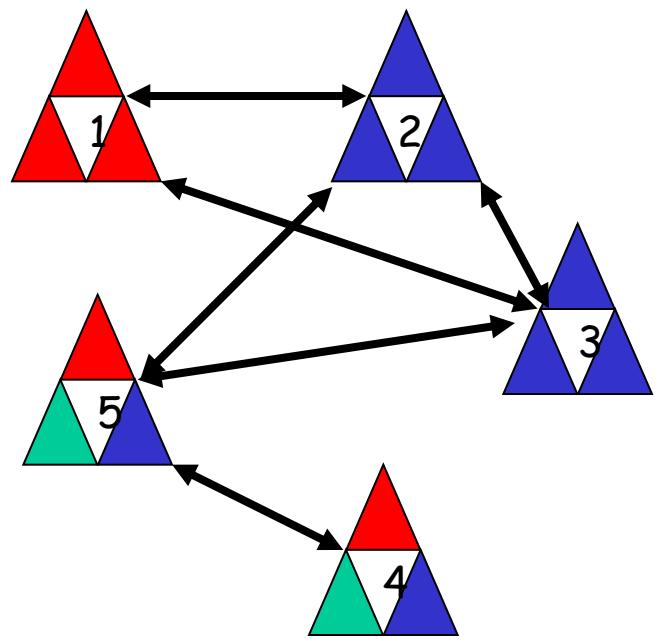
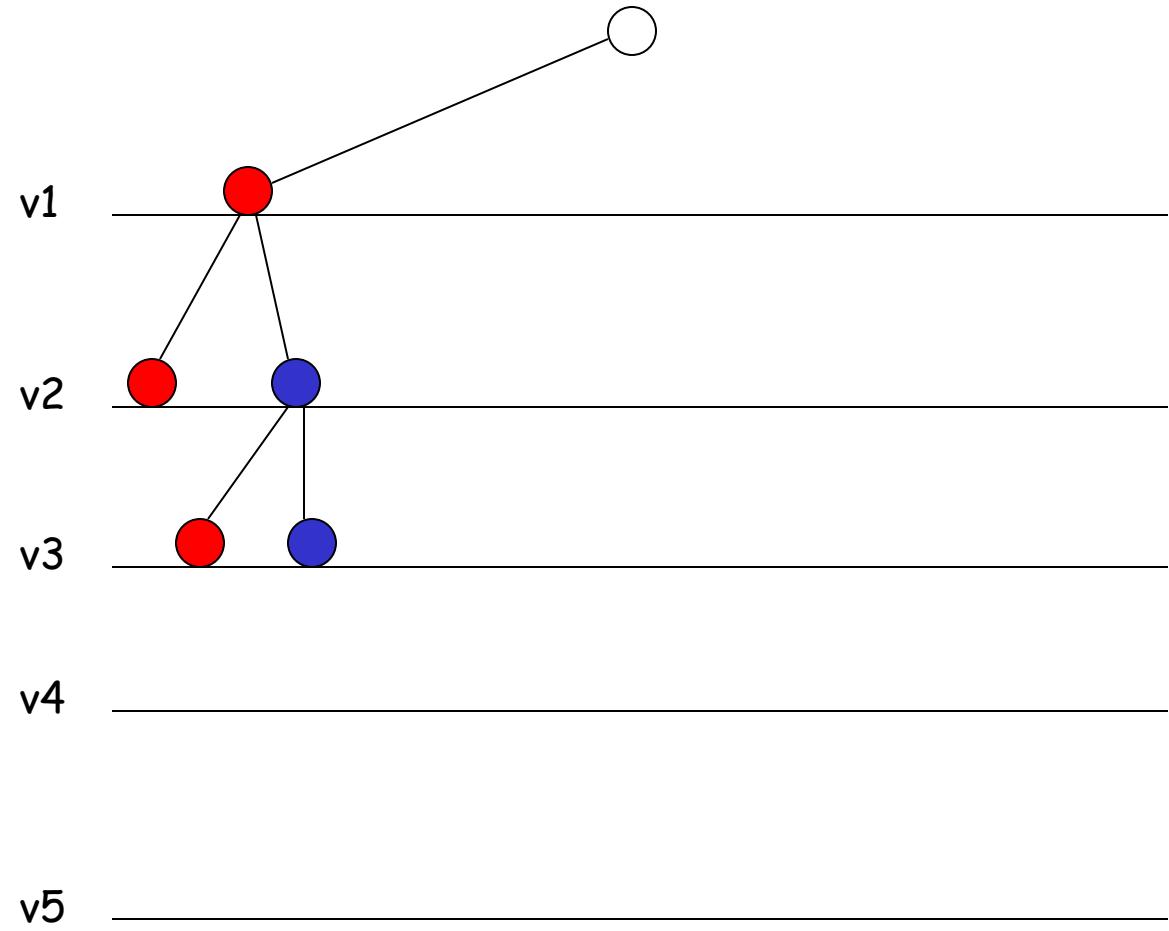
A Tree Trace of BT (assume domain ordered {R,B,G})



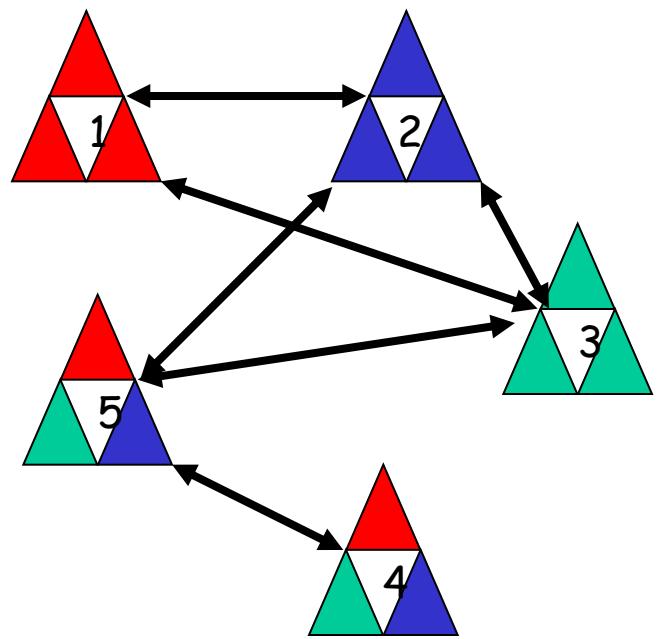
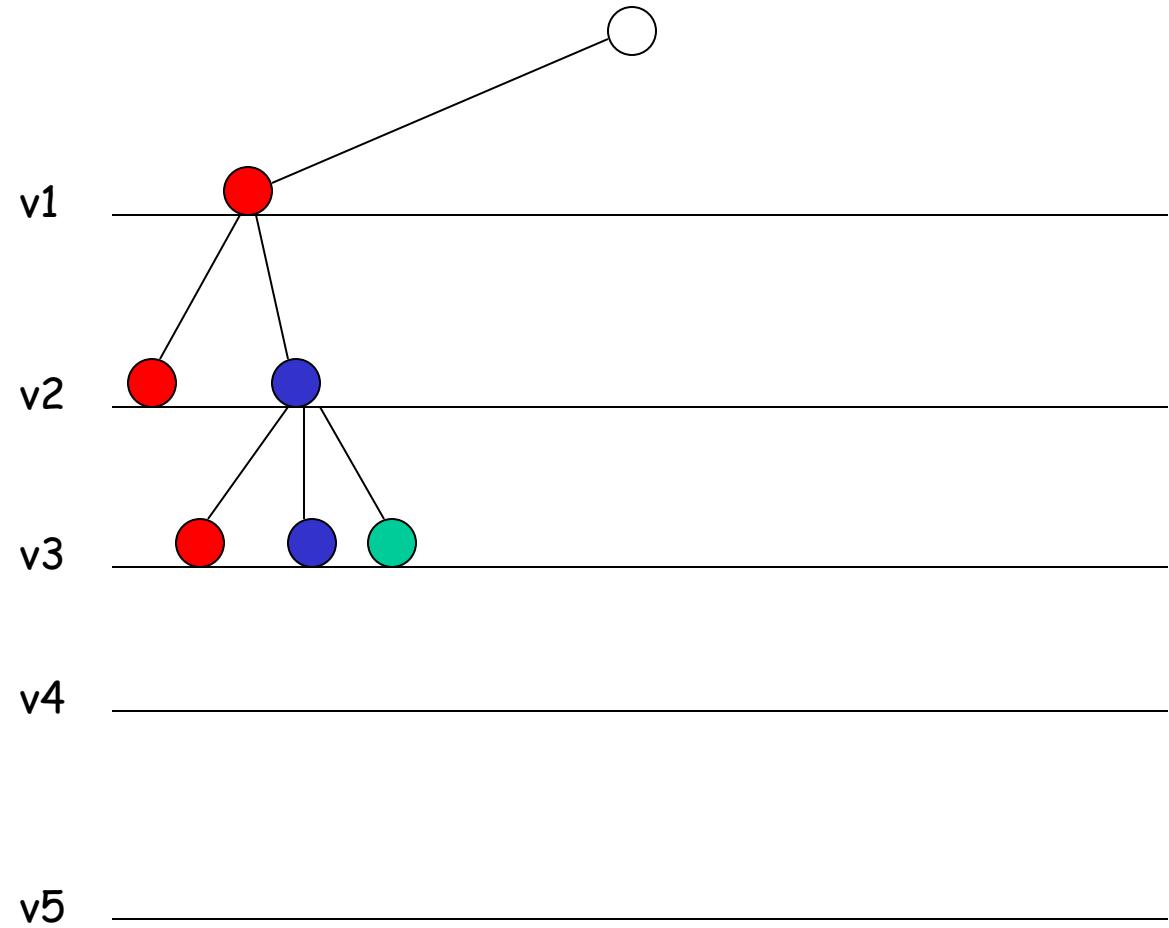
A Tree Trace of BT (assume domain ordered {R,B,G})



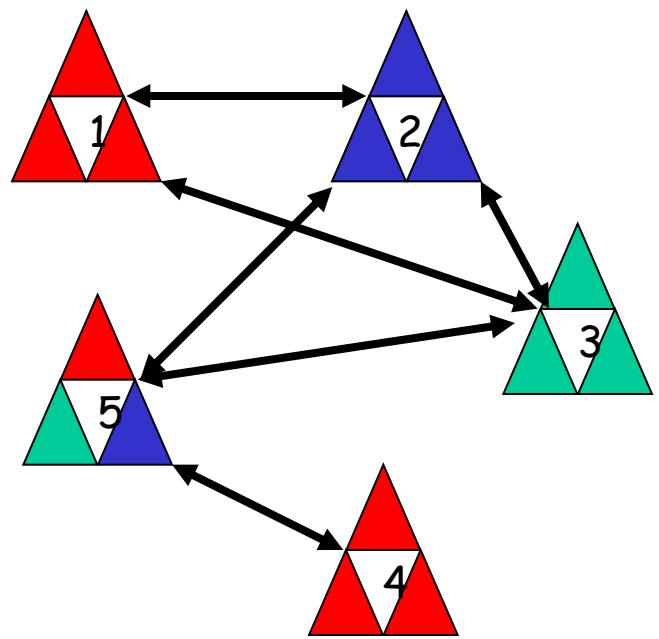
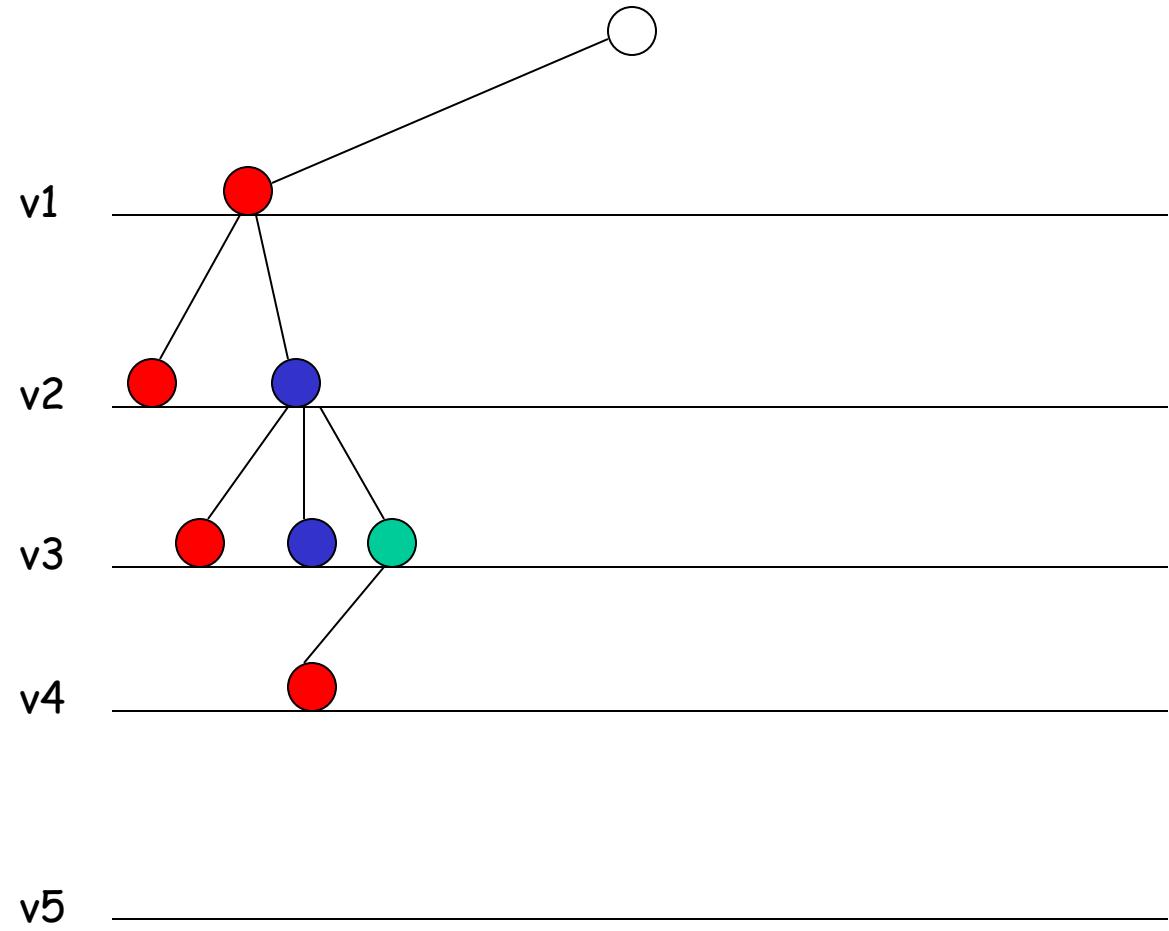
A Tree Trace of BT (assume domain ordered {R,B,G})



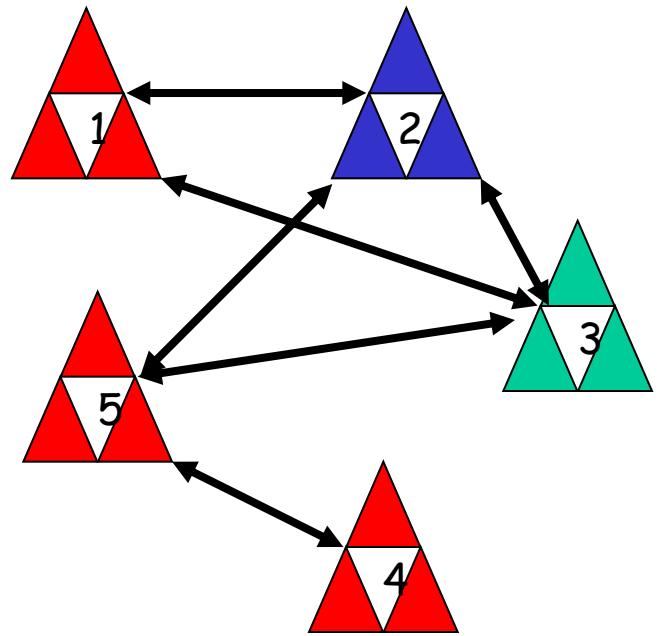
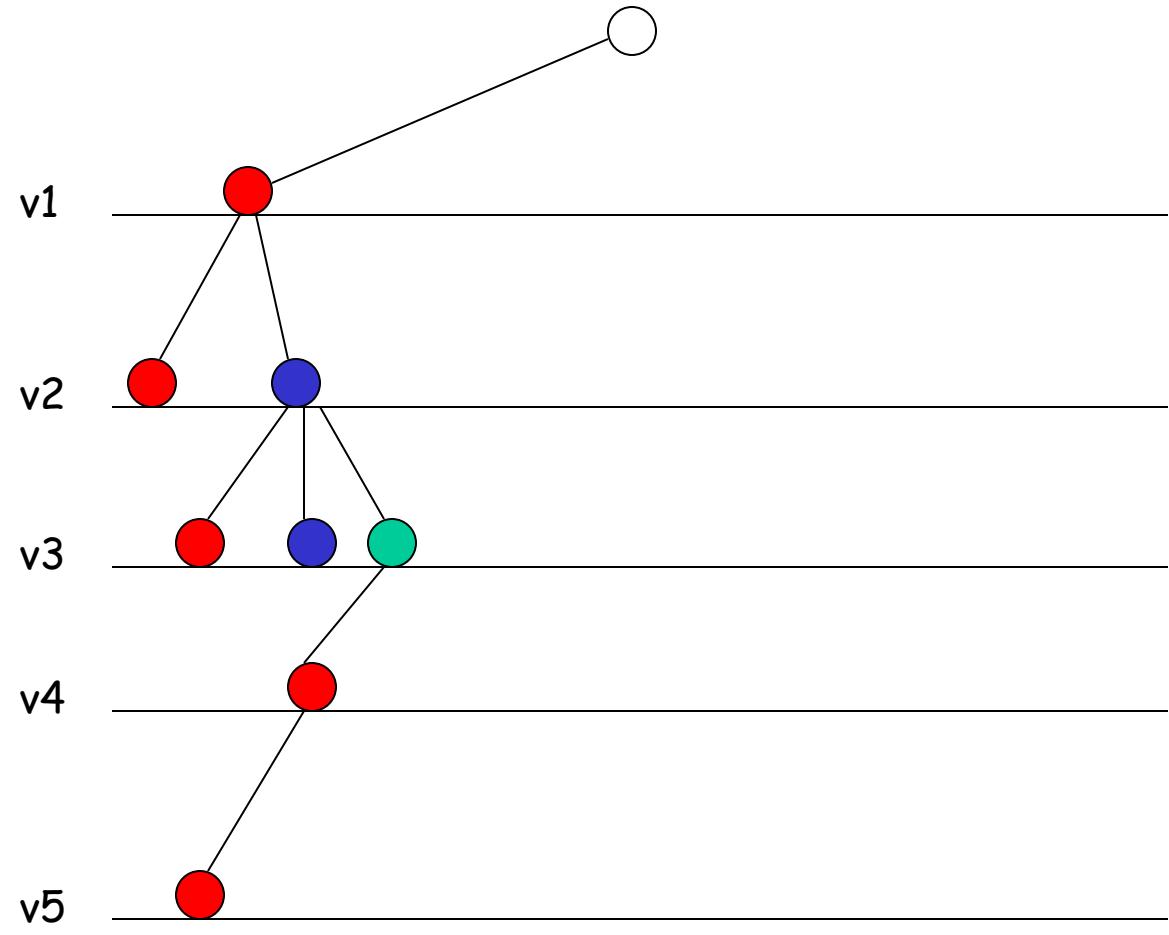
A Tree Trace of BT (assume domain ordered {R,B,G})



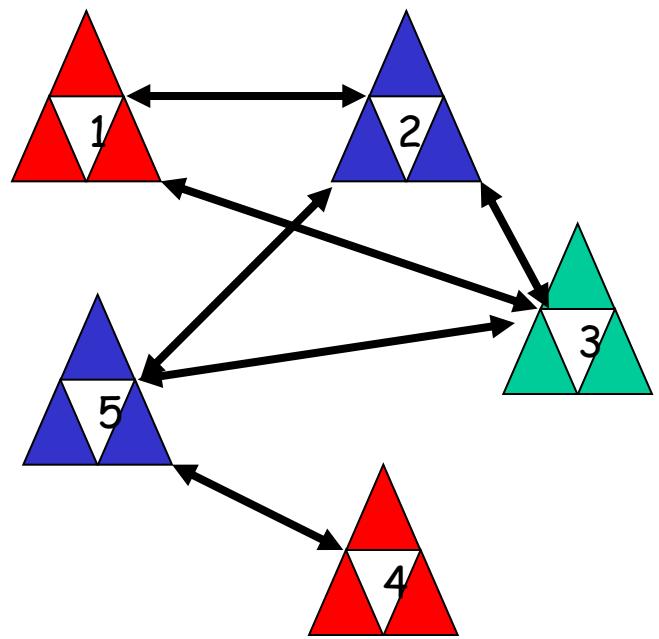
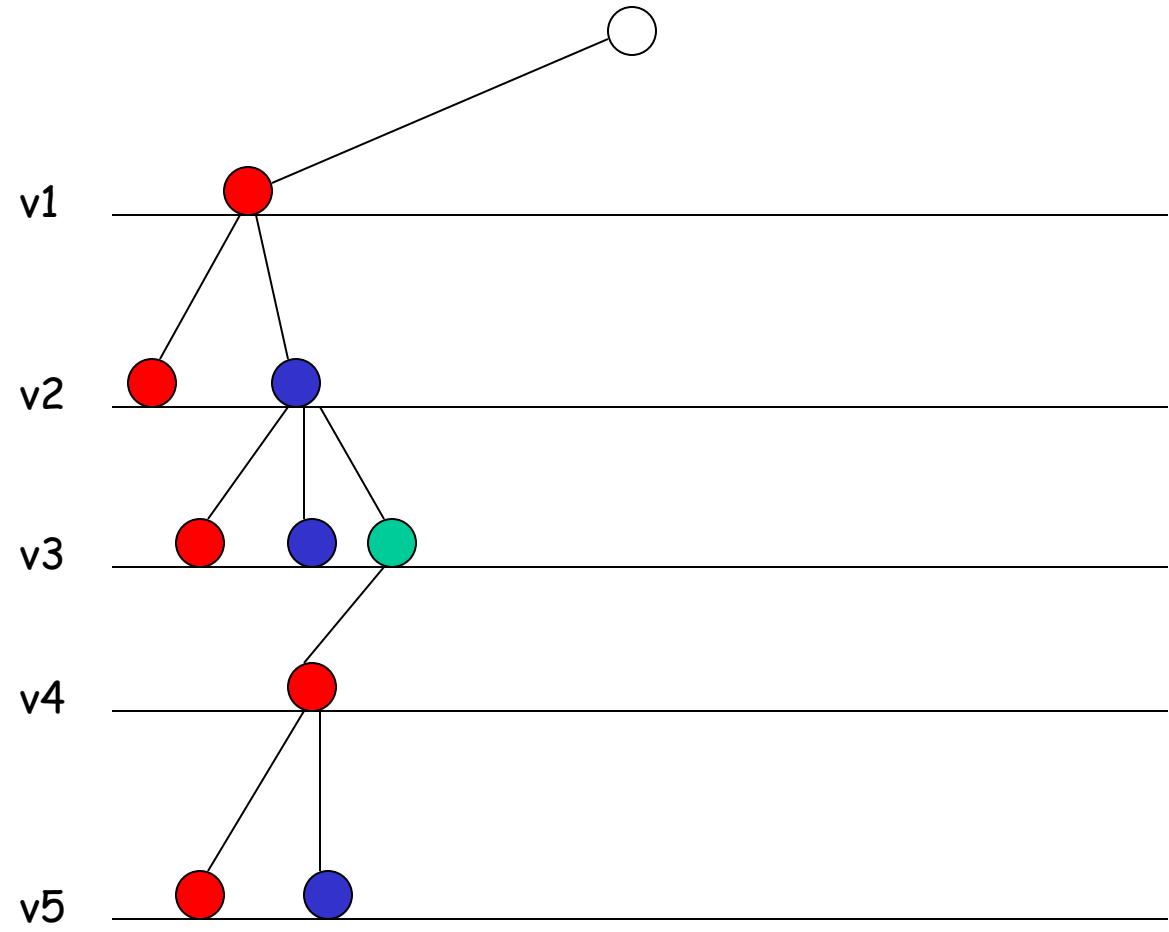
A Tree Trace of BT (assume domain ordered {R,B,G})



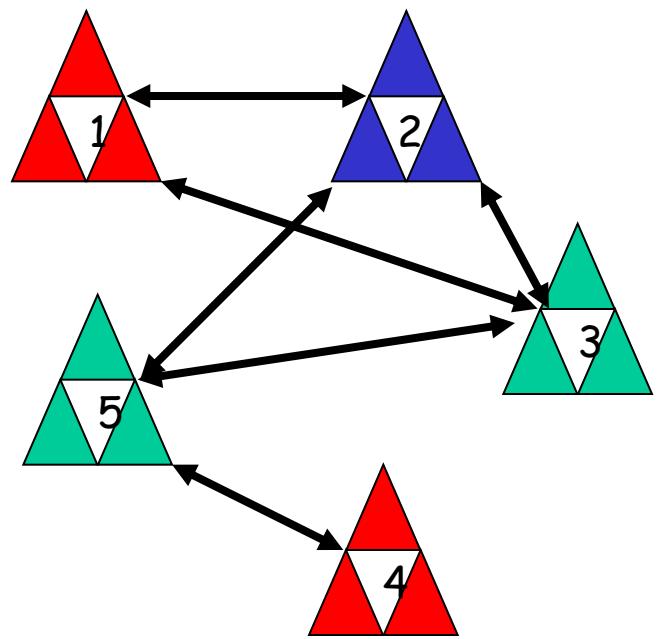
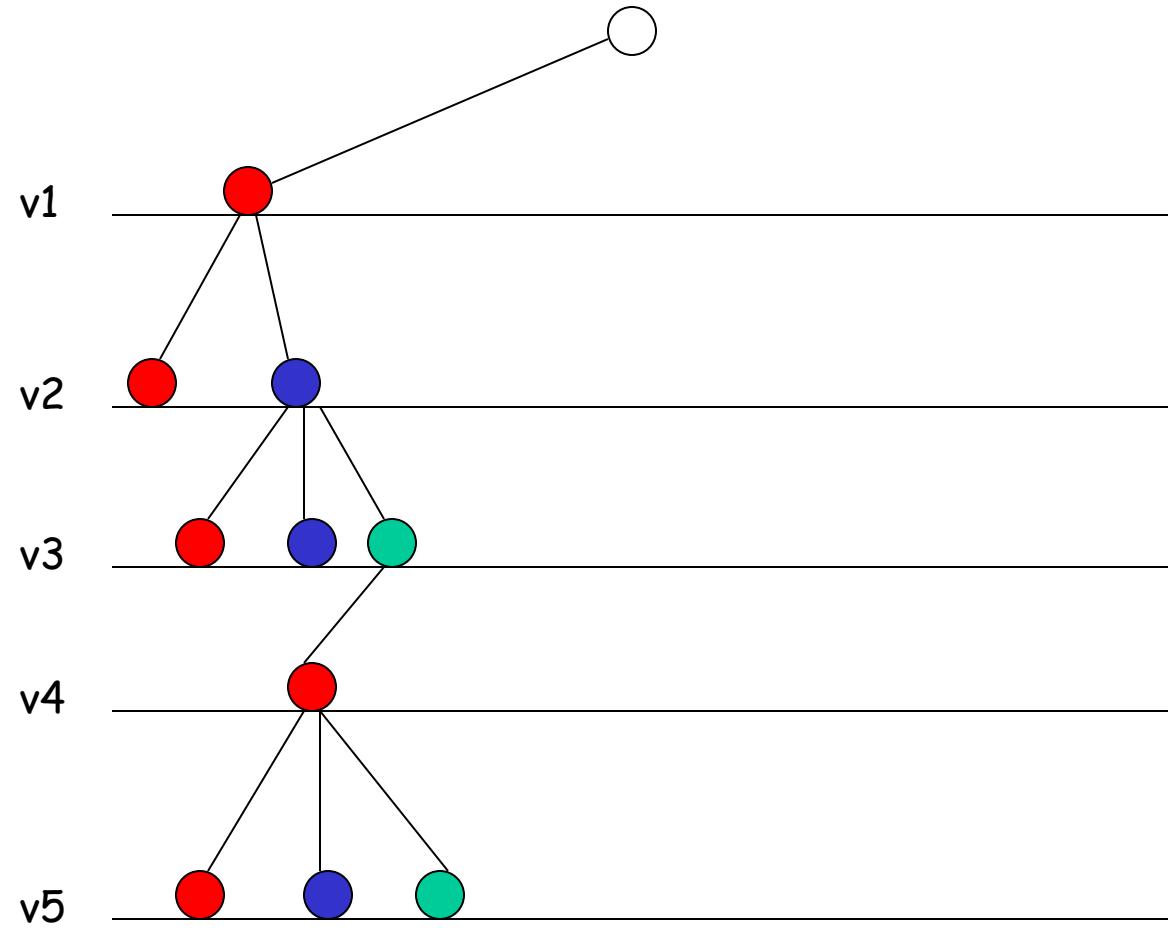
A Tree Trace of BT (assume domain ordered {R,B,G})



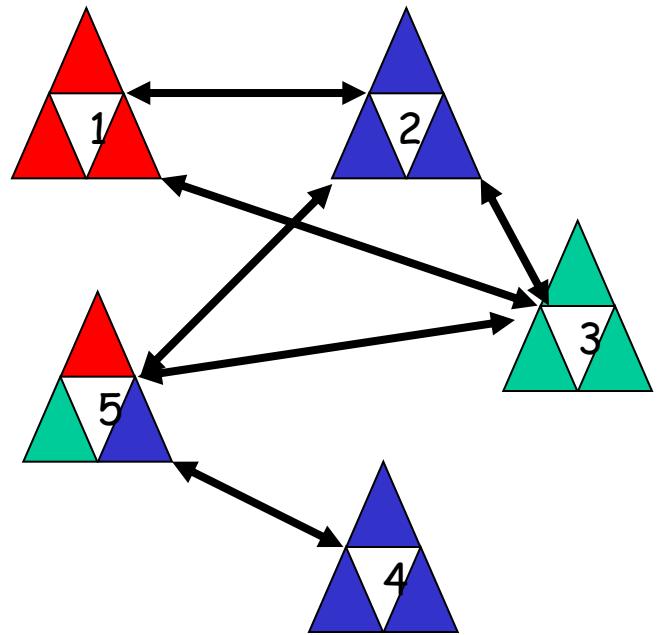
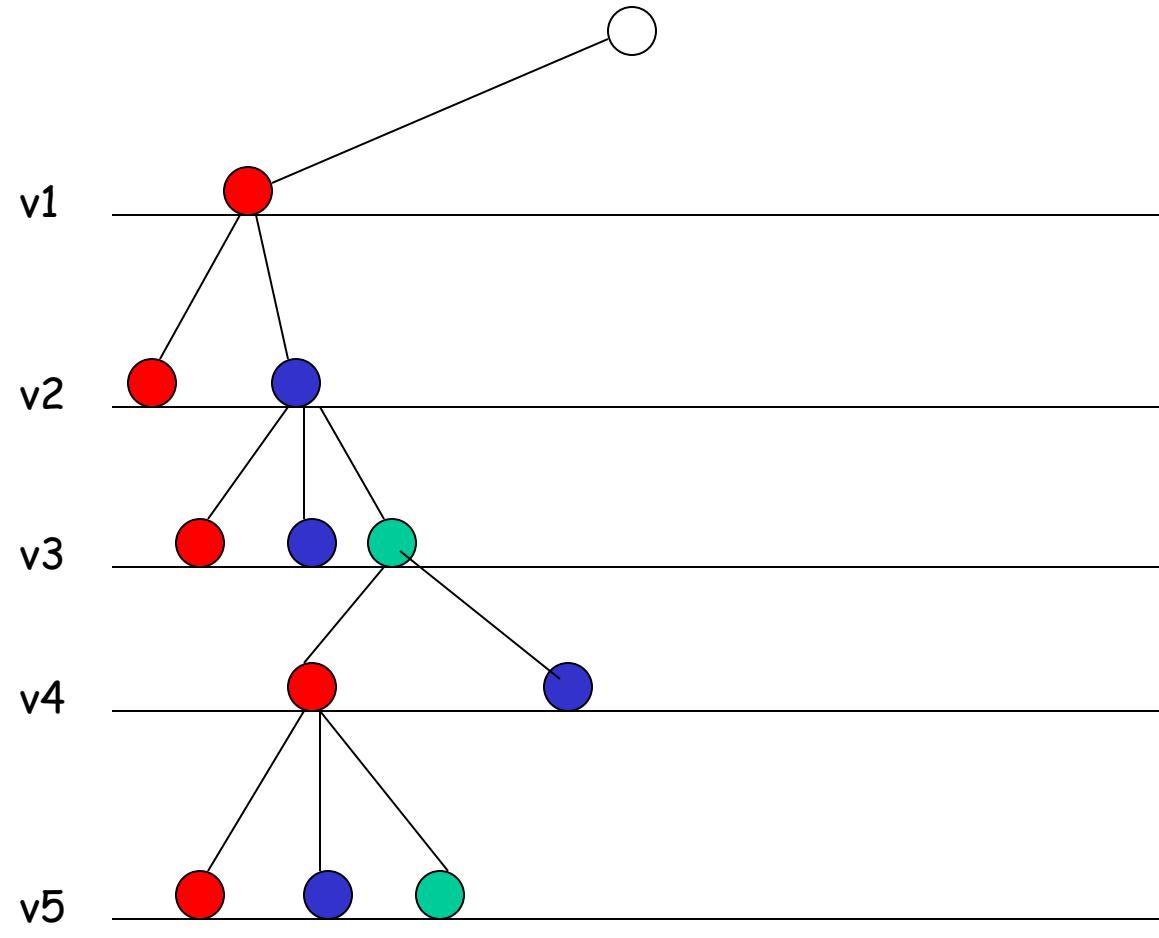
A Tree Trace of BT (assume domain ordered {R,B,G})



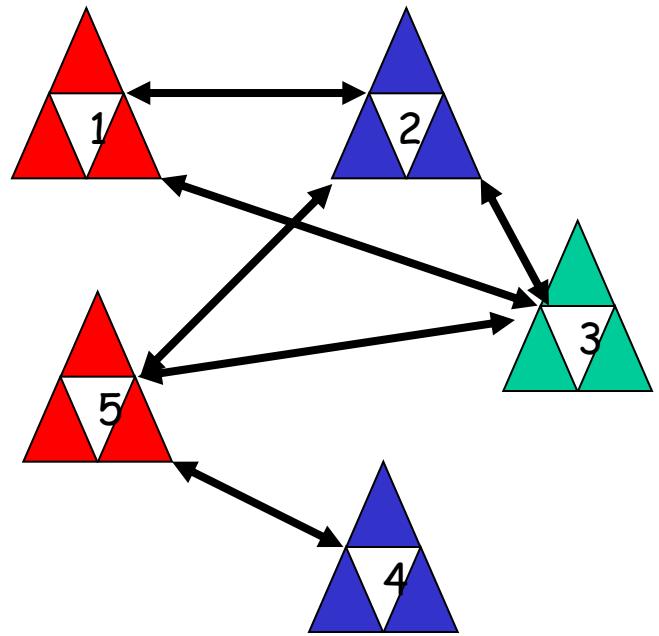
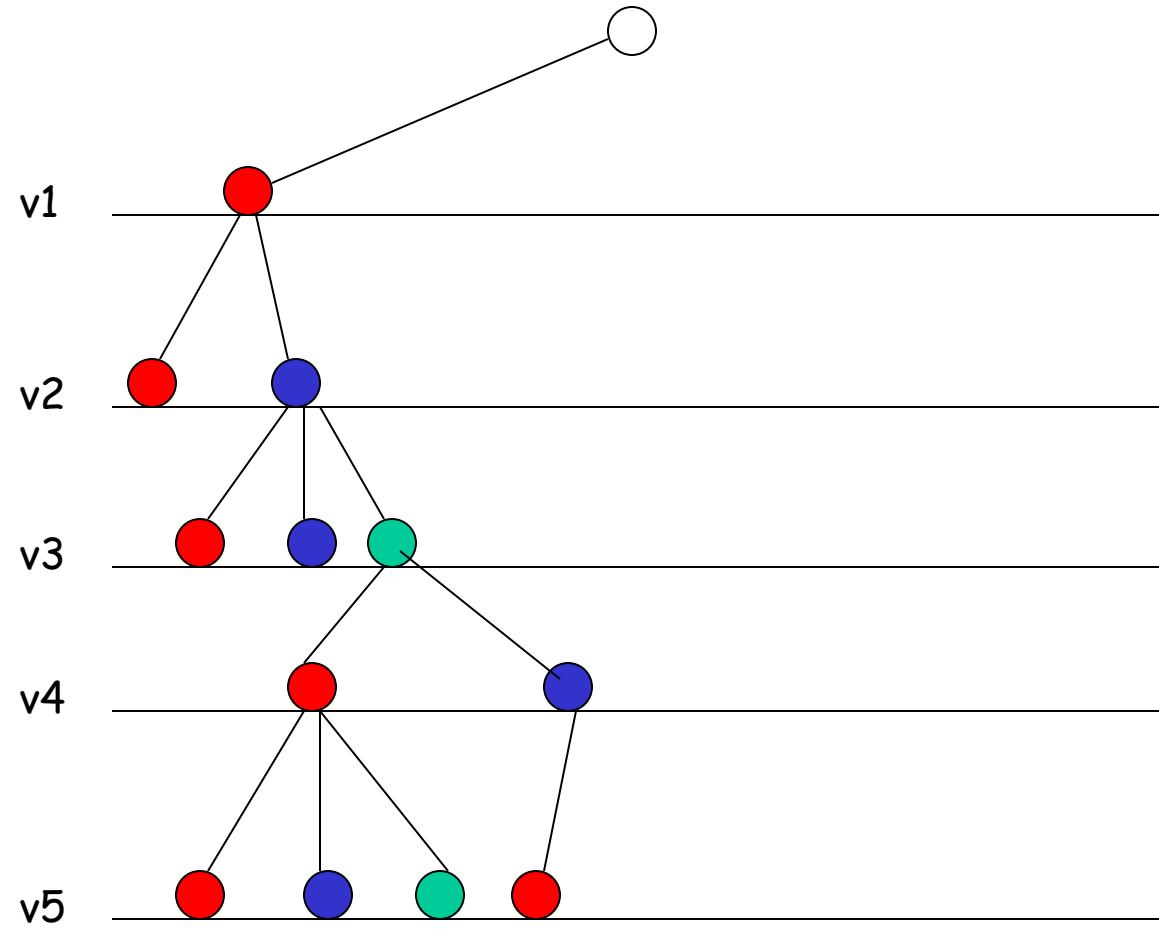
A Tree Trace of BT (assume domain ordered {R,B,G})



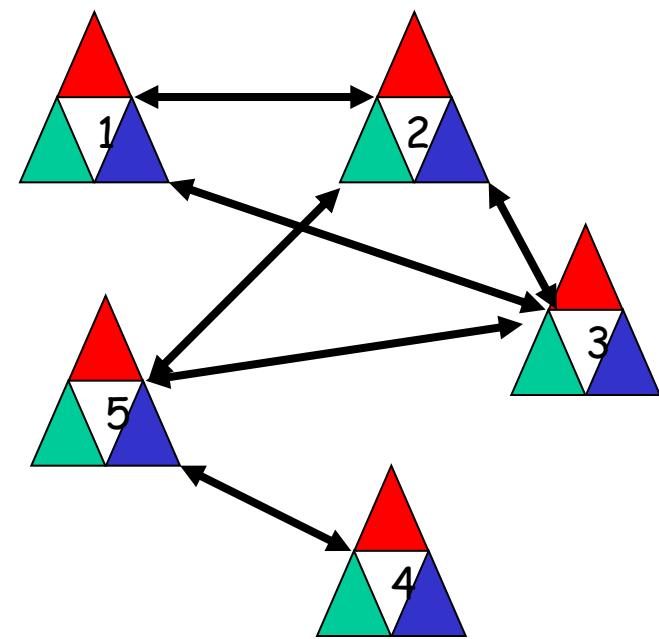
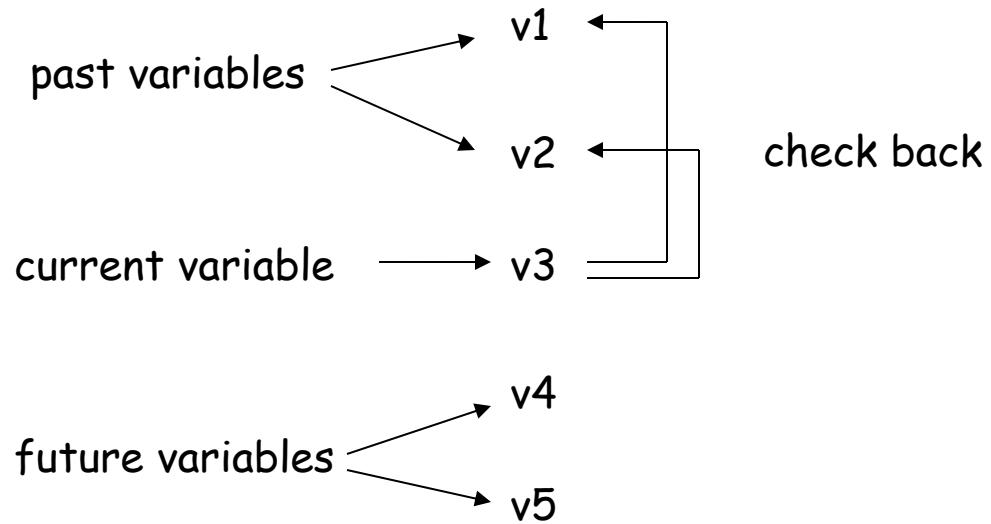
A Tree Trace of BT (assume domain ordered {R,B,G})



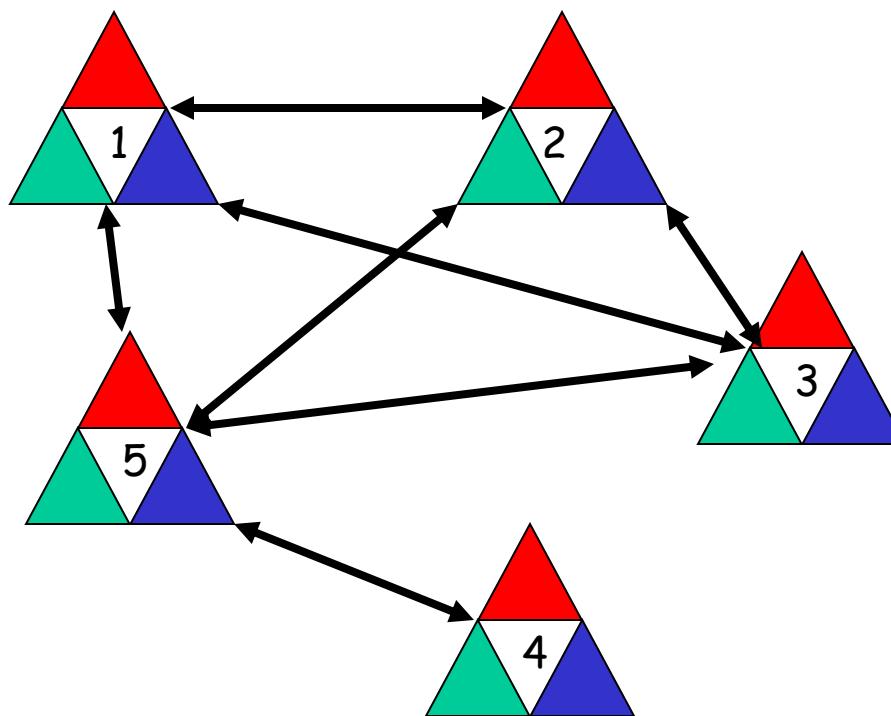
A Tree Trace of BT (assume domain ordered {R,B,G})



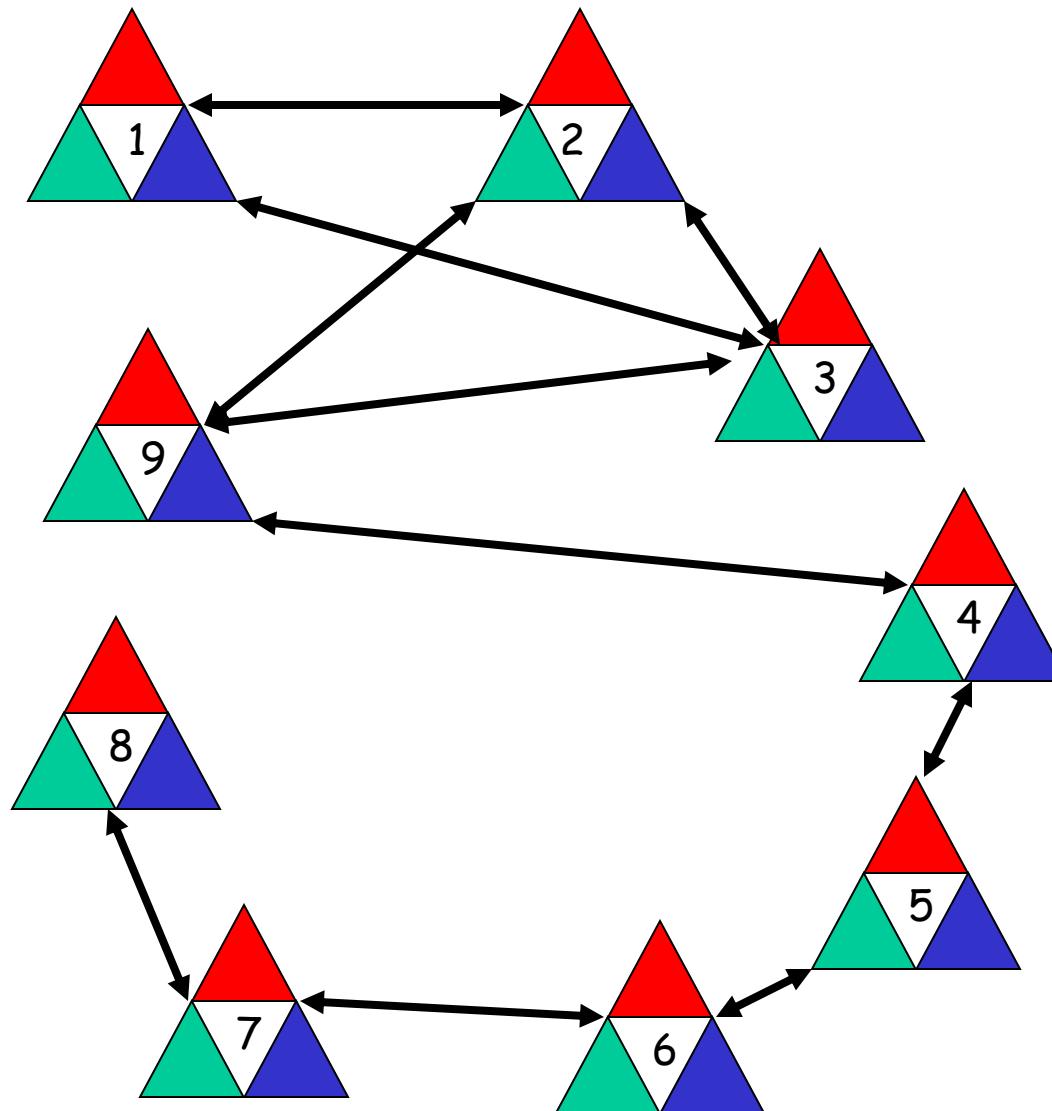
Another view



Can you solve this (csp2)?



Thrashing? (csp3c)

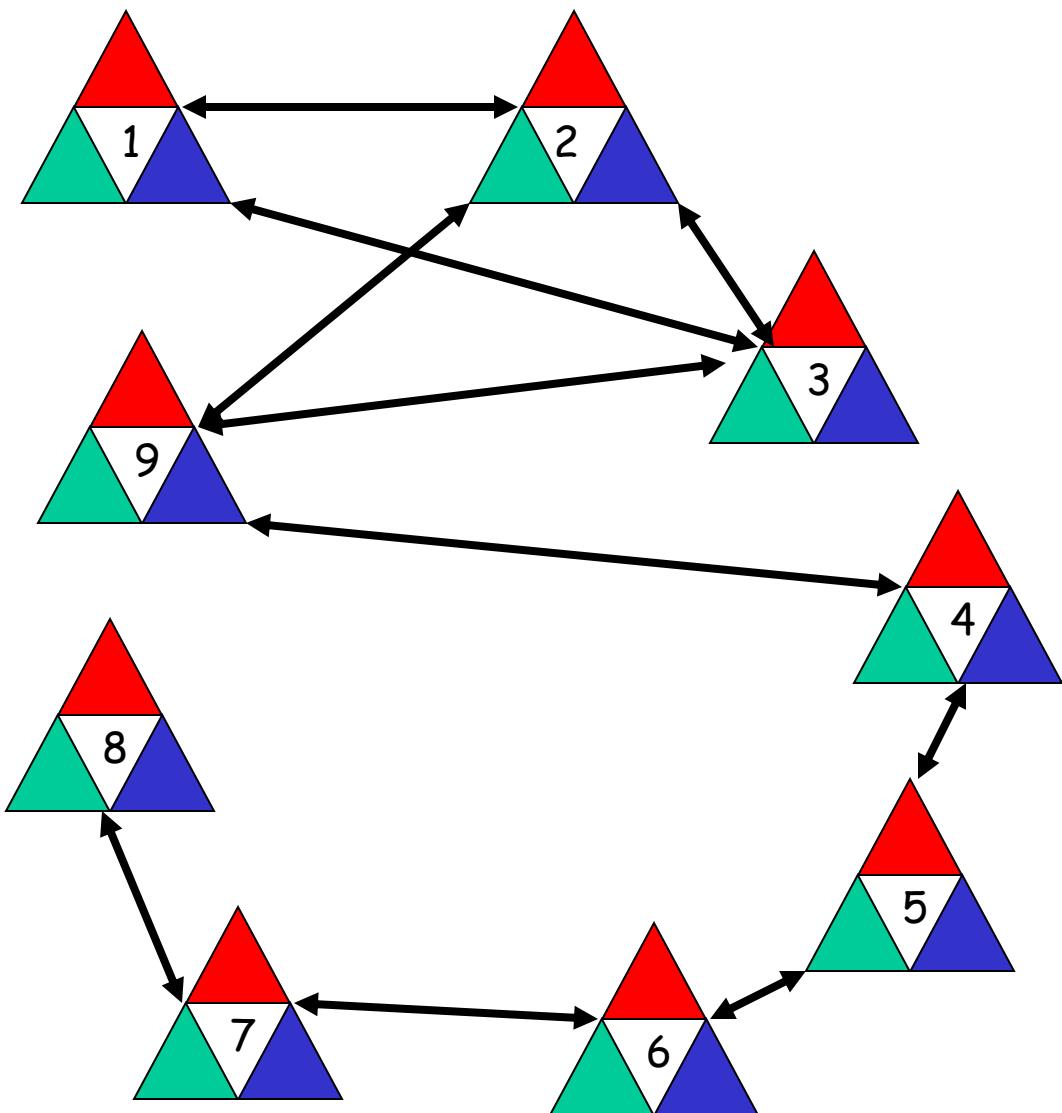


Thrashing? (csp3c)

```
V1 = R  
v2 = B  
v3 = G  
v4 = R  
v5 = B  
v6 = R  
v7 = B  
v8 = R  
v9 = conflict
```

The cause of the conflict with v9
is v4, but what will bt do?

Find out how it goes with
bt3 and csp4



questions

- Why measure checks and nodes?
 - What is a node anyway?
 - Who cares about checks?
- Why instantiate variables in lex order?
 - Would a different order make a difference?
- How many nodes could there be?
- How many checks could there be?
- Are all csp's binary?
- How can we represent constraints?
- Is it reasonable to separate V from D ?
 - Why not have variables with domains
- What is a *constraint graph*?
 - How can (V, D, C) be a graph?
- What is BT "thinking about"? i.e. why is it so dumb?
- What could we do to make BT smarter?
- Is BT doing any inferencing/propagation?
- What's the 1st reference to BT?
 - Golomb & Baumert JACM 12, 1965?
 - The Minotaur?